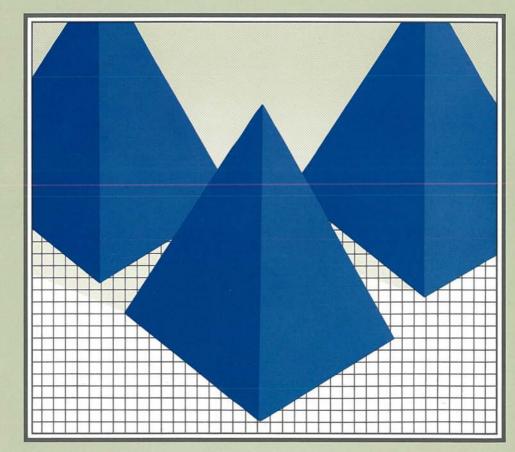
HEWLETT-PACKARD

HP 9000 Series 300 Computers

Using ARPA Services



Using ARPA Services

HP 9000 Series 300



Manual Part Number: 50952-90001

Printed in U.S.A., January 1989

Notice

Hewlett-Packard makes no warranty of any kind with regard to this material, including, but not limited to, the implied warranties of merchantability and fitness for a particular purpose. Hewlett-Packard shall not be liable for errors contained herein or for incidental or consequential damages in connection with the furnishing, performance, or use of this material.

Hewlett-Packard assumes no responsibility for the use or reliability of its software on equipment that is not furnished by Hewlett-Packard.

© Copyright 1987, 1988, 1989 Hewlett-Packard Company.

This document contains proprietary information, which is protected by copyright. All rights are reserved. No part of this document may be photocopied, reproduced or translated to another language without the prior written consent of Hewlett-Packard Company. The information contained in this document is subject to change without notice.

Restricted Rights Legend

Use, duplication or disclosure by the Government is subject to restrictions as set forth in paragraph (b)(3)(B) of the Rights in Technical Data and Software clause in DAR 7-104.9(a).

© Copyright 1980, 1984, AT&T, Inc.

© Copyright 1979, 1980, 1983, The Regents of the University of California.

© Copyright, 1979, 1987, Sun Microsystems, Inc.

This software and documentation is based in part on the Fourth Berkeley Software Distribution under license from the Regents of the University of California.

DEC and VAX are registered trademarks of Digital Equipment Corp.

MS-DOS[®] is a trademark of Microsoft Corp.

 $\text{UNIX}^{\ensuremath{\mathbb{R}}}$ is a U.S. registered trademark of AT&T in the U.S.A. and in other countries.

NFS is a trademark of Sun Microsystems, Inc.

Hewlett-Packard Co. 3404 E. Harmony Rd. Fort Collins, CO 80525 U.S.A.

Printing History

December 1986 ... Edition 1.

December 1987 ... Edition 2.

January 1989 . . . Edition 3.

Table of Contents

Chapter 1: Documentation Overview

Manual Overview
Who Should Read This Manual 1-1
What Is in This Manual 1-2
Conventions in This Manual 1-5
Reference Manual Guide 1-6
Chapter 2: Services Overview
Introduction
Getting Started
The ARPA/Berkeley Services
Services Listed by Function
Sending Mail to a Remote Host
Listing Information about a Remote Host
Logging into a Remote Host 2-5
Transferring Files to or from a Remote Host 2-6
Executing Commands on a Remote Host 2-6
Obtaining General Information 2-7
Interprocess Communication 2-8
Chapter 3: Sending Mail
Introduction
Using Sendmail 3-2
Executing Sendmail
Mailing to Files

Sendmail Operations Overview 3-3 Collecting Messages 3-3 Routing the Messages 3-3 More Mail System Information 3-6
Chapter 4: Listing Hosts with Ruptime
Using Ruptime
Displaying Ruptime Status Lines
Sorted by Host Name in Alphabetical Order
Sorted by Host Name in Reverse Alphabetical Order 4-5
Sorted by Decreasing Uptime
Sorted by Increasing Uptime
Sorted by Decreasing Number of Users 4-8
Sorted by Increasing Number of Users 4-9
Sorted by Decreasing Load Average 4-10
Sorted by Increasing Load Average 4-11
Chapter 5: Listing Users with Rwho
Using Rwho
Listing Active and Likely Active Users of Network Hosts 5-3
Listing All Users of Network Hosts
Chapter 6: Logging into a Host with Telnet
Using Telnet
1. Invoke Telnet
2. Change the Telnet Escape Character If Necessary 6-3
3. Connect to a Remote Host 6-7
4. Log into the Remote Host 6-8
Giving Telnet Commands When Telnet Is in Its Input State 6-9
Checking the Behavior of Carriage Returns from a Remote Host . 6-10
Changing the Carriage Return Mode Setting 6-10
Disconnecting from a Remote Host and/or Exiting Telnet 6-11
Disconnecting from a Remote Host and Remaining in Telnet 6-12
Exiting from Telnet When Telnet Is in Its Input State 6-13
Exiting from Telnet When Telnet Is in Its Command State 6-13
Obtaining Help 6-14

Listing the Telnet Commands	
Temporarily Returning to HP-UX on Your Local Host	
Executing a Single HP-UX Command on Your Local Host	
Working for an Extended Time on Your Local Host	6-18
Obtaining Telnet Status	6-20
Changing Where User Input Is Echoed	
Changing User Input Mode	
Connecting to a Remote Host When You Invoke Telnet	
Chapter 7: Logging into a Host with Rlogin	
Determining If You Need to Change the Rlogin Escape Character .	
Caution: What Not to Change the Rlogin Escape Character To	
Determining What Size Characters to Send with Rlogin	
When You Can Send Eight-Bit Characters	
When You Must Send Seven-Bit Characters	
Using Rlogin	
Automatic Login	
Manual Login	
If You Get Unexpected Results after Logging into a Remote Host .	
Logging Out of the Remote Host and Exiting Rlogin	
Temporarily Returning to HP-UX on Your Local Host	
Executing a Single HP-UX Command on Your Local Host	
Working for an Extended Time on Your Local Host	7-18
Passing the Rlogin Escape Character to a Remote Program	7-19
Logging into a Remote Host as Someone Else	
Giving Other Remote Users Rlogin Access to Your Local Account	7-24
Protecting Your .rhosts File	
Using Rlogin's "Shorthand" Syntax	7-26

Chapter 8: Transferring Files with Ftp

Using Ftp	8-1
1. Invoke Ftp	8-2
2. Choose Whether to Display Responses from a Remote Host	8-2
3. Connect to a Remote Host	8-4
4. Log into the Remote Host	8-5
Disconnecting from a Remote Host and Exiting Ftp	8-7
Exiting Ftp to Return to HP-UX on Your Local Host	8-7
Disconnecting from a Remote Host and Remaining in Ftp	8-8
Obtaining Help	8-9
Listing the Ftp Commands	8-9
Getting Information about a Specific Ftp Command	8-9
Temporarily Returning to HP-UX on Your Local Host	8-10
Executing a Single HP-UX Command on Your Local Host	8-10
Working for an Extended Time on Your Local Host	8-11
How Ftp Treats "Wild Card" Characters, or Metacharacters	8-12
Turning Globbing On or Off	8-12
Performing Directory Operations with Ftp	8-14
Changing the Local Working Directory	8-15
Changing the Remote Working Directory	8-16
Listing the Contents of the Remote Working Directory	8-17
Listing the Contents of a Remote Directory	8-20
Listing the Contents of Multiple Remote Directories	8-23
Displaying the Name of the Remote Working Directory	8-28
Creating a Remote Directory	8-29
Deleting a Remote Directory	8-30
Changing the Name of a Remote Directory	8-31
Transferring Files with Ftp	8-33
1. Set the Local and Remote Working Directories	8-33
2. Set the File Transfer Type	8-33
3. Choose Options to Monitor File Transfer Progress	8-36
4. Turn on Interactive Mode for Selective File Transfers	8-37
5. Perform One or More File Transfers	8-39
Performing Other File Operations with Ftp	8-62

Displaying the Contents of a Remote File	8-62
Creating a Remote File	
Appending Text to the End of a Remote File	8-65
Deleting a Remote File	8-68
Deleting Multiple Remote Files	8-69
Changing the Name of a Remote File	8-72
Obtaining Ftp Status	
Setting Up Automatic Remote Login for Ftp	8-75
Protecting Your .netrc File	
Logging into a Remote Host with a Login Not in Your .netrc File	
The Public Ftp Account	8-81
Logging into the Public (Anonymous) Ftp Account	
Specifying Ftp Settings and Connecting to a Remote Host When You	
InvokeFtp	8-83
Chapter 9: Transferring Files with Rcp	
Chapter 5. Transferring Files with http	
FILO O	0.1
File Copy Concepts	
Using Rcp	. 9-3
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host	. 9-3 . 9-3
Using Rcp	. 9-3 . 9-3
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host	. 9-3 . 9-3 . 9-6
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp	. 9-3 . 9-3 . 9-6 . 9-8
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer	. 9-3 . 9-3 . 9-6 . 9-8 9-12
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer .	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer From One or More Remote Producers to a Remote Consumer From Local and Remote Producers to a Local Consumer	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16 9-20
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer From One or More Remote Producers to a Remote Consumer From Local and Remote Producers to a Local Consumer From Local and Remote Producers to a Remote Consumer	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16 9-20 9-24
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer From One or More Remote Producers to a Remote Consumer From Local and Remote Producers to a Local Consumer From Local and Remote Producers to a Remote Consumer From Local and Remote Producers to a Remote Consumer Rep's Effect on File Attributes	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16 9-20 9-24 9-28
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer From One or More Remote Producers to a Remote Consumer From Local and Remote Producers to a Local Consumer From Local and Remote Producers to a Remote Consumer Rcp's Effect on File Attributes	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16 9-20 9-24 9-28
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer From One or More Remote Producers to a Remote Consumer From Local and Remote Producers to a Local Consumer From Local and Remote Producers to a Remote Consumer From Local and Remote Producers to a Remote Consumer Rcp's Effect on File Attributes Using "Wild Card" Characters, or Metacharacters with Rcp Copying Remote Files and Directories as Someone Else on the	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16 9-20 9-24 9-28 9-29
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer From One or More Remote Producers to a Remote Consumer From Local and Remote Producers to a Local Consumer From Local and Remote Producers to a Remote Consumer Rcp's Effect on File Attributes Using "Wild Card" Characters, or Metacharacters with Rcp Copying Remote Files and Directories as Someone Else on the RemoteHost	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16 9-20 9-24 9-28 9-28 9-29
Using Rcp Creating a \$HOME/.rhosts File on a Remote Host Performing Copy Operations with Rcp From a Local Producer to a Remote Consumer From One or More Remote Producers to a Local Consumer From One or More Remote Producers to a Remote Consumer From Local and Remote Producers to a Local Consumer From Local and Remote Producers to a Remote Consumer From Local and Remote Producers to a Remote Consumer Rcp's Effect on File Attributes Using "Wild Card" Characters, or Metacharacters with Rcp Copying Remote Files and Directories as Someone Else on the	. 9-3 . 9-3 . 9-6 . 9-8 9-12 9-16 9-20 9-24 9-28 9-29 9-29 9-30 9-31

ix

Chapter 10: Executing Commands with Remsh

Setting Up Permission to Use Remsh on a Remote Host	10-2
Creating a \$HOME/.rhosts File on a Remote Host	10-3
Executing Commands on a Remote Host as Yourself	10-5
Executing Commands on a Remote Host as Someone Else	10-7
Giving Other Remote Users Remsh Access to Your Local Account	10-8
Protecting Your .rhosts File	10-9
Executing More Than One Remote Command with Remsh	10-10
Using Shell Metacharacters with Remsh	10-12
Stdin, Stdout, and Stderr for Remsh	10-12
Using Remsh with Remote Commands That Do Not Take Input .	10-14
Using Remsh's "Shorthand" Syntax	10-17
Chapter 11: Interprocess Communication	
Overview of IPC	11-2
How You Can Use IPC	11-3
The Client-Server Model	11-4
IPC Library Routines	11-6
Key Terms and Concepts	11-7
IPC Using Internet Stream Sockets	11-10
Preparing Address Variables	11-11
Writing the Server Process	11-16
Writing the Client Process	11-22
Sending and Receiving Data	11-25
Closing a Socket	11-29
Example Using Stream Sockets	11-30
BSD IPC Using UNIX Domain Stream Sockets	11-44
Preparing Address Variables	11-46
Writing the Server Process	11-47
Writing the Client Process	11-54
Sending and Receiving Data	11-57
Closing a Socket	11-61
Examples Using UNIX Domain Stream Sockets	11-62
Advanced Topics for Stream Sockets	11-71

Socket Options	11-71
Synchronous I/O Multiplexing with Select	11-78
Sending and Receiving Data Asynchronously	11-80
Nonblocking I/O	11-83
Using Shutdown	11-84
Using Read and Write to Make Stream Sockets Transparent	11-86
Sending and Receiving Out of Band Data	11-86
IPC Using Internet Datagram Sockets	11-92
Preparing Address Variables	11-94
Writing the Server and Client Processes	11-99
Binding Socket Addresses to Datagram Sockets	11-100
Sending and Receiving Messages	11-102
Closing a Socket	11-107
Example Using Datagram Sockets	11-108
Advanced Topics for Internet Datagram Sockets	11-117
Specifying a Default Socket Address	11-117
Synchronous I/O Multiplexing with Select	11-119
Sending and Receiving Data Asynchronously	11-119
Nonblocking I/O	11-120
Using Broadcast Addresses	11-121
Programming Hints	11-122
Troubleshooting	11-122
Port Addresses	11-123
Using Diagnostic Utilities as Troubleshooting Tools	11-124
Adding a Server Process to the Internet Daemon	11-125
Summary Tables for System and Library Calls	11-130
Appendix A: Portability Issues	
Porting Issues for IPC Functions and Library Calls	. A-2
Porting Issues for Other Functions and Library Calls Typically Used	by
IPC	
Glossary	
In days	

Documentation Overview

Note

Before you read this manual, read the *Networking Overview: NS-ARPA and NFS Services* booklet for an introduction to important terms and concepts. The booklet positions HP 9000 Series 300 networking products relative to each other and lists and describes the components of each product. The booklet also contains specific network connectivity diagrams and a detailed documentation map.

Manual Overview

Who Should Read This Manual

This manual is written primarily for people who have:

- some experience with the HP-UX environment; and
- access to and familiarity with the HP-UX Reference manuals.

What Is in This Manual

The list below briefly describes the contents of each chapter in this manual.

Chapter 1: Introduction

The remainder of this chapter describes what you need to get started and provides a list of reference manuals.

Chapter 2: Services Overview

To aid you in finding the service that best suits your needs, this chapter lists and briefly describes the ARPA/Berkeley Services according to their function. The "Services Overview" chapter also briefly describes the Interprocess Communication package and how to obtain general information about your system once the NS-ARPA Services/300 product has been installed.

Chapter 3: Sending Mail

This chapter briefly describes the internetwork mail routing facility provided with the ARPA/Berkeley Services. Since this facility is automatically installed when the product is installed, but is **not** executable, your node manager may choose whether to make this facility executable on your local host. Plan to read this chapter only if your node manager has made this facility executable on your local host.

Chapter 4: Listing Hosts with Ruptime

This chapter explains how to list the names and condition of network hosts. The chapter also explains how to sort the list based on various items in the list.

Chapter 5: Listing Users with Rwho

This chapter explains how to list information about users logged into network hosts.

Chapter 6: Logging Into a Host with Telnet

This chapter explains how to use *telnet* to log into a remote host.

Chapter 7: Logging Into a Host with Rlogin

This chapter explains how to use *rlogin* to log into a remote host. It also describes how to give other network users *rlogin* access to your local account.

Chapter 8: Transferring Files with Ftp

This chapter explains how to use *ftp* to transfer files between your local host and a remote host and to perform remote file management operations.

Chapter 9: Transferring Files with Rcp

This chapter explains how to use *rcp* to transfer files and directories among network hosts. The chapter also describes how to give other network users *rcp* access to your local account.

Chapter 10: Executing Commands with Remsh

This chapter explains how to use *remsh* to execute a command on a remote host. The chapter also describes how to give other network users *remsh* access to your local account.

Chapter 11: Interprocess Communication

This chapter is for programmers who intend to use Interprocess Communication (IPC) based on 4.2 BSD programming development tools. The chapter describes how to use datagram and stream sockets. An overview of IPC and example programs are also included.

Appendix A: Portability Issues

This appendix explains portability issues and differences between HP's implementation of 4.2 BSD Interprocess Communication (sockets) and the Berkeley sockets.

Glossary

The glossary lists and defines terms used in this manual.

Index

The index provides a quick page-reference for subjects contained within this manual.

Conventions in This Manual

Notation	Description
boldface	Boldfacing is used for emphasis.
computer_font	Words in syntax statements that are not in italics must be entered exactly as shown. Punctuation characters other than brackets, braces and ellipses must also be entered exactly as shown.
italics	Words in syntax statements that are in italics denote a parameter that you must supply.
Return	This font is used to indicate a key on the computer's keyboard.
CTRL-D	This convention is used to indicate a combina- tion of keys to press simultaneously for a desired function.
	An ellipsis in a syntax statement indicates that a previous element may be repeated. In addition, vertical and horizontal ellipses are used in examples to indicate that portions of the example have been omitted.
-1	This is a command argument that appears fre- quently in the manual. We mention it here so you will not confuse the letter "1" with the number "1."

Reference Manual Guide

For more information on the following subjects, refer to the publications listed in the right column.

For information on:	Read:
HP-UX Operating System (HP 9000)	Introducing UNIX System V HP-UX Concepts and Tutorials Beginner's Guide series for HP-UX HP-UX Reference manuals
ARPA/Berkeley Manual Reference Pages	ARPA/Berkeley Services Reference Pages
C Programming Language	The C Programming Language, Brian W. Kernighan, Dennis M. Ritchie; Prentice-Hall, Inc.
	C Programming Guide, Jack Purdum, Que Corporation, Indianapolis
Networking	Computer Networks, Andrew S. Tanenbaum; Prentice-Hall, Inc.
	Networking Overview: NS-ARPA and NFS Services and X.25
NS Part of the NS- ARPA Services/300 Product	Using Network Services
NS-ARPA Services/300 Installation, Configura- tion, Maintenance, and Troubleshooting	Installing and Maintaining NS-ARPA Services

Services Overview

Introduction

The ARPA Services part of the NS-ARPA Services/300 product enables your HP 9000 Series 300 to transfer files, log into remote hosts, execute commands remotely, and send mail to and receive mail from remote hosts that are either on your network or accessible by your network.

The ARPA Services part of the NS-ARPA Services/300 product is a subset of networking services originally developed by the University of California at Berkeley (UCB) for the Advanced Research Projects Agency (ARPA) and UCB. The services originally developed for ARPA are called "ARPA Services." The services originally developed for UCB are called "Berkeley Services."

UCB developed the services based on the Berkeley Software Distribution of UNIX¹, version 4.2 (4.2 BSD).

4.2 BSD programming development tools for interprocess communication are also provided with the ARPA Services part of the NS-ARPA Services/300 product.

This chapter briefly explains the ARPA/Berkeley Services. To guide you to the service you need for a desired task, the services are listed by functionality. For tutorial information about individual services, read this manual. For specific details about individual services, read the reference page for the service.

⁽¹⁾ UNIX is a U.S. registered trademark of AT&T in the U.S.A. and other countries.

This chapter also lists and describes:

- the sources from which you can obtain additional information about your local host or network; and
- the 4.2 BSD-based ARPA Services/300 Interprocess Communication package.

Getting Started

Before you begin, make sure that:

- your node manager has installed the NS-ARPA Services/300 product on your local host and has brought up the network;
- you have asked your node manager for all the login names you may be associated with;
- you have asked your node manager what other hosts or nodes your Series 300 can communicate with.

Note

The computer you are working on is referred to as your **local host**, and all other computers (hosts) on the network are remote in relation to your local host.

The ARPA/Berkeley Services

HP's implementation of the ARPA/Berkeley Services is actually a combination of services originating from the Advanced Research Projects Agency (ARPA) and from the University of California at Berkeley (UCB).

The services originating from the ARPA environment are used to communicate in HP-UX, UNIX and non-UNIX environments.

Services originating from UCB are used for HP-UX or UNIX operations only.

The services from each environment are shown in the figure below.

ARPA Services	Berkeley Services
File Transfer Protocol (ftp)	Remote copy (rcp)
Telnet (telnet)	Remote login (rlogin)
Simple Mail Transfer Protocol (SMTP)	Remote execution (rexec)
	Remote shell (remsh)
	Remote uptime (ruptime)
	Remote who (rwho)
	Internetwork Mail Routing (sendmail)
	Interprocess Communication ("IPC" or "Berkeley sockets")

ARPA/Berkeley Services

Services Listed by Function

The number in parentheses next to the service name, e.g., *rlogin(1)*, corresponds to the section in the *ARPA/Berkeley Services Reference Pages* that documents the service.

Sending Mail to a Remote Host

sendmail(1M) originates from UCB and, when installed, works with your network's mailers to perform internetwork mail routing among HP-UX, UNIX and non-UNIX hosts on the network. When used in the command line, sendmail does not provide a friendly user interface. Sendmail supports mail aliasing and forwarding and uses ARPA's standard Simple Mail Transfer Protocol (SMTP).

Listing Information about a Remote Host

ruptime(1) is a Berkeley Service. It is used to list information about HP-UX or UNIX hosts on the network that are running the *rwho* daemon, *rwhod*. The information that *ruptime* displays includes host names, whether the hosts are up or down, the number of active users on the remote host and three numeric fields containing the 1-, 5- and 15-minute load averages for the number of processes in the remote host's run queue. rwho(1) is a Berkeley Service. It is used to list information about HP-UX or UNIX hosts on the network that are running the rwho daemon, rwhod. The information that rwho displays includes the user names of those who are actively logged into remote or local hosts on the network, the remote or local hosts' names, the users' terminal lines, the users' login times, and the amount of time each user has been idle.

Logging into a Remote Host

Rlogin and *telnet* allow you to log into a remote host on the network if you have an account on the remote host.

telnet(1)	originates from the ARPA environment and is used to log into a remote, HP-UX, UNIX or non-UNIX host. You must use <i>telnet</i> if the remote host is a non- UNIX host (i.e., VAX/VMS) and you must supply a login name and password to log into the remote host.
rlogin(1)	originates from UCB and is used to log in from a local HP-UX or UNIX host to a remote HP-UX or UNIX host without being prompted for a login name and password.

Transferring Files to or from a Remote Host

ftp(1)	originates from the ARPA environment and allows you to transfer files among HP-UX, UNIX and non-UNIX hosts on the network. For example, if you want to transfer a UNIX file to a remote host using an MS-DOS format, you would use <i>ftp. Ftp</i> is the file transfer program which uses the ARPA standard File Transfer Protocol (FTP).
rcp(1)	originates from UCB and allows you to transfer files between only HP-UX or UNIX hosts on the network.

Executing Commands on a Remote Host

remsh(1)	originates from UCB and allows you to execute commands on a remote HP-UX or UNIX host on the network. <i>Remsh</i> is the same command as <i>rsh</i> from the Berkeley Software Distribution of UNIX, version 4.2 (4.2 BSD). No login names or passwords are used although your account permis- sions are verified on the remote host before you can remotely execute a command.
rexec(3X)	originates from UCB and is a library routine used to execute commands on a remote HP-UX or UNIX host on the network. A customized program must be written to use <i>rexec</i> . The advantage of using <i>rexec</i> is that it can be used in programs and passwords can be specified.

Obtaining General Information

This section describes the sources from which you can obtain additional information.

hosts(4)	The <i>letc/hosts</i> file contains host names and internet addresses of remote hosts. You may not be limited to communicating only with those hosts listed in <i>letc/hosts</i> . Check with your node manager to verify which hosts are available for communication with your local host.
hosts.equiv(4)	The <i>/etc/hosts.equiv</i> file lists remote hosts on the net- work that are "equivalent" to your local host.
intro(3N)	The <i>intro</i> HP-UX reference page lists and briefly describes the Internet network library functions.
netrc(4)	The <i>netrc(4)</i> reference page describes the <i>.netrc</i> file used by <i>rexec</i> and <i>ftp</i> to determine login names and passwords to remote hosts.
networks(4)	The <i>letc/networks</i> file lists the official network name, number and aliases for networks that your local host recognizes.

\$HOME/.rhosts	The <i>\$HOME/.rhosts</i> file may be created by each user on the local host to specify remote user names that are equivalent to the local user. The local host then permits equivalent remote users to access the local user's account without requiring a password. <i>\$HOME/.rhosts</i> is described on the <i>hosts.equiv(4)</i> reference page.
protocols(4)	The <i>/etc/protocols</i> file contains the official protocol names, protocol numbers and protocol aliases recognized by your local host.
services(4)	The <i>letc/services</i> file contains the official service names, port numbers, protocol names used by the service and service name aliases that correspond to the service.

Interprocess Communication

The NS-ARPA Services/300 product provides programmers with an Interprocess Communication (IPC) package that allows processes to communicate with other local and remote processes through system calls. HP's IPC implementation is based on the IPC in 4.2 BSD.

Two transport protocols are available:

- Transmission Control Protocol (TCP), which is the transport protocol for stream sockets, and
- User Datagram Protocol (UDP), which is the transport protocol for datagram sockets.

For details on IPC, refer to the "Interprocess Communication" chapter in this manual.

Sending Mail

Introduction

Sendmail(1M), the internetwork mailing facility supplied with the ARPA/Berkeley Services, acts as a central post office that determines the internetwork routing needed for mail delivery to local or remote users. It routes messages to local users, files and programs. Sendmail also enables your local host to send mail to and receive mail from other hosts on a local area network or through a gateway. In addition, message aliasing and forwarding can also be specified.

Because *sendmail* is typically used in environments where internetwork communications are frequent or heavy, your node manager **may** have installed *sendmail* on your system. Ask your node manager if *sendmail* has been installed on your system before you continue reading this section.

Note

For details about *sendmail* and message aliasing and forwarding, refer to the *Installing and Maintaining NS-ARPA Services* manual.

Using Sendmail

Executing Sendmail

You can execute sendmail in two ways:

- Whenever a standard mailer is accessed, *sendmail* is automatically invoked. Standard HP-UX mailing programs are *mail* and *mailx*.
- You can use the *sendmail* command with arguments on the command line. Because this method does **not** provide a friendly user interface, it is typically used only in programs.

Mailing to Files

If you want to send a message to a local file, you must specify the filename as an absolute path (i.e., you must begin the filename with a slash "/").

If the file does not exist, you must own and have search (execute) permission in the directory in which the file is to be created.

If the file already exists, is not executable and is writable by all users, the message will be appended to the file.

Note

Sendmail does not write to executable files.

Sendmail Operations Overview

Sendmail performs its task in two phases: it collects messages and then routes them. While collecting and routing messages, message-address interpretation is controlled by a production system that manages both network-style addressing (e.g., user@host) and UUCP-style addressing (e.g., host!user). This production system is defined by the contents of the *sendmail* configuration file.

Collecting Messages

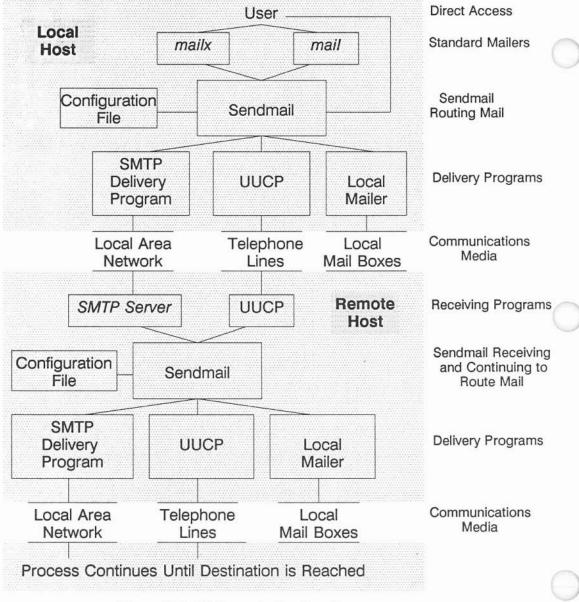
If *sendmail* is invoked via a standard mail program or by using the *sendmail* command on the command line, it collects the message from *stdin* and the argument list.

Routing the Messages

Once *sendmail* collects the message, it routes the information. To route the message, *sendmail*:

- rewrites the mail addresses of the recipients to conform to the standards of the target network;
- if necessary, adds lines to the message header so that the information is available for a recipient to use in a reply; and
- passes the mail to one of several specialized delivery agents for delivery.

Sendmail is also executed by the program that receives mail from the network. When incoming mail arrives, a receiving program passes the mail to sendmail for routing in the same way that a mailer invokes sendmail.



The figure below outlines the flow of messages through sendmail.

Flow of Mail Through Sendmail

Routing to Remote Hosts

If a recipient's host is on the LAN, *sendmail* uses the SMTP delivery module to send the message to the remote host on the network.

If the recipient has a UUCP address, *sendmail* calls the *uux* program to deliver the message on the remote system.

Routing to Local Destinations

If the recipient is a local user, *sendmail* calls the local mailer to deposit the message in the recipient's mailbox. The HP-UX local mailer is *rmail*.

If the recipient is a local file, *sendmail* writes the message to the file. This is the only case in which *sendmail* directly **delivers** a message to a destination.

Routing Error Messages

During mail transfer processes, *sendmail* creates a transcript of each mail transaction to send to the originator if the message is permanently undeliverable. This transcript contains any error messages that occur during the attempted mail delivery.

If an error status indicates that the delivery failed but might be successful if re-tried, *sendmail* stores the message on a queue for later delivery. *Sendmail* attempts to send the message again when it next processes the queue.

More Mail System Information

If *sendmail* is installed, the commands listed below or the corresponding HP-UX reference pages provide additional information about the *sendmail* program or your mail system.

mailq	is described on the <i>sendmail(1M)</i> reference page and prints a list of mail messages that are in the mail queue.
mailstats(1)	prints the mail traffic statistics.
praliases(1)	displays any aliases that your system recognizes.
uupath(1)	expands UUCP-style and network-style addresses.

Listing Hosts with Ruptime

Ruptime is a Berkeley Service that lists information about HP-UX or UNIX hosts on the network. The status information that *ruptime* displays includes:

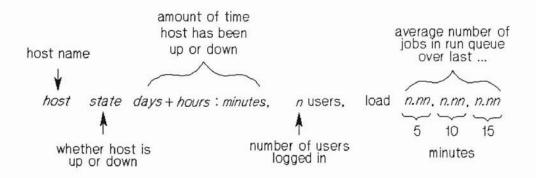
- network host names,
- whether each network host is up or down,
- the number of active users on each network host, and
- the average number of jobs in each network host's run queue over the last one, five, and fifteen minutes.

This information is useful in determining:

- · which network hosts you can perform work on,
- which network hosts are most heavily or least heavily loaded, and
- how responsive each network host is likely to be over the network.

Using Ruptime

For each network host, *ruptime* displays a status line with the following format:



Ruptime Status Line

Note

Ruptime does not count users who have not used the system for an hour or more. You can have ruptime count these idle users by invoking the command with the -a option as shown in the "Displaying Ruptime Status Lines" section.

Ruptime can display status lines (with or without idle users) sorted by:

- host name in alphabetical order,
- host name in reverse alphabetical order,
- decreasing uptime,
- increasing uptime,
- decreasing number of users,
- increasing number of users,
- · decreasing load average, and
- increasing load average.

The following sections tell how to display *ruptime* status lines sorted in the ways described above.

Displaying Ruptime Status Lines

Sorted by Host Name in Alphabetical Order

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime

Result:

Ruptime displays a list similar to the following:

hpabca	down 14+08:3	34				
hpabcb	down 1:1	13				
hpabcc	up 1+17:4	40, 4 u	users, loa	d 0.18,	0.13,	0.09
hpabcd	up 14+06:4	49, 3 1	users, loa	d 0.10,	0.38,	0.49

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a

Result:

Ruptime displays a list similar to the following:

hpabca	down 14+08	:34				
hpabcb	down 1	:13				
hpabcc	up 1+17	:40, 6 users,	load 0	.18,	0.13,	0.09
hpabcd	up 14+06	:49, 3 users,	load 0	.10,	0.38,	0.49

4-4 Displaying Ruptime Status Lines

Sorted by Host Name in Reverse Alphabetical Order

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime -r

Result:

Ruptime displays a list similar to the following:

hpabcd	up	14+06:49,	3	users,	load	0.10,	0.38,	0.49
hpabcc	up	1+17:40,	4	users,	load	0.18,	0.13,	0.09
hpabcb	down	1:13						
hpabca	down	14+08:34						

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a -r

Result:

Ruptime displays a list similar to the following:

hpabcd	up 14+06:49,	3 users,	load 0.10,	0.38, 0.49
hpabcc	up 1+17:40,	6 users,	load 0.18,	0.13, 0.09
hpabcb	down 1:13			
hpabca	down 14+08:34			

Sorted by Decreasing Uptime

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime -t

Result:

Ruptime displays a list similar to the following:

hpabcd	up 14+06:49	, 3 users,	load 0.10,	0.38, 0.49
hpabcc	up 1+17:40	, 4 users,	load 0.18,	0.13, 0.09
hpabcb	down 1:13			
hpabca	down 14+08:34			

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a -t

Result:

Ruptime displays a list similar to the following:

hpabcd	up 14+06:49,	3 users,	load 0.10,	0.38, 0.49
hpabcc	up 1+17:40,	6 users,	load 0.18,	0.13, 0.09
hpabcb	down 1:13			
hpabca	down 14+08:34			

4-6 Displaying Ruptime Status Lines

Sorted by Increasing Uptime

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime -r -t

Result:

Ruptime displays a list similar to the following:

hpabca	down	14+08:34				
hpabcb	down	1:13				
hpabcc	up	1+17:40,	4 users,	load 0.18,	0.13,	0.09
hpabcd	up	14+06:49,	3 users,	load 0.10,	0.38,	0.49

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a -r -t

Result:

Ruptime displays a list similar to the following:

hpabca	down 14+08:34				
hpabcb	down 1:13				
hpabcc	up 1+17:40,	6 users,	load 0.18,	0.13, 0.0	9
hpabcd	up 14+06:49,	3 users,	load 0.10,	0.38, 0.4	9

Sorted by Decreasing Number of Users

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime -u

Result:

Ruptime displays a list similar to the following:

hpabcc	up :	1+17:40,	4	users,	load	0.18,	0.13,	0.09
hpabcd	up 14	4+06:49,	3	users,	load	0.10,	0.38,	0.49
hpabcb	down	1:13						
hpabca	down 14	4+08:34						

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a -u

Result:

Ruptime displays a list similar to the following:

hpabcc	up	1+17:40,	6	users,	load	0.18,	0.13,	0.09
hpabcd	up	14+06:49,	3	users,	load	0.10,	0.38,	0.49
hpabcb	down	1:13						
hpabca	down	14+08:34						

4-8 Displaying Ruptime Status Lines

Sorted by Increasing Number of Users

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime -r -u

Result:

Ruptime displays a list similar to the following:

hpabca	down	14+08:34					
hpabcb	down	1:13					
hpabcd	up	14+06:49,	3 users,	load	0.10,	0.38,	0.49
hpabcc	up	1+17:40,	4 users,	load	0.18,	0.13,	0.09

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a -r -u

Result:

Ruptime displays a list similar to the following:

hpabca	down	14+08:34					
hpabcb	down	1:13					
hpabcd	up	14+06:49,	3 users,	load	0.10,	0.38,	0.49
hpabcc	up	1+17:40,	6 users,	load	0.18,	0.13,	0.09

Sorted by Decreasing Load Average

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime -1

Result:

Ruptime displays a list similar to the following:

hpabcc	up	1+17:40,	4	users,	load	0.18,	0.13,	0.09
hpabcd	up	14+06:49,	3	users,	load	0.10,	0.38,	0.49
hpabcb	down	1:13						
hpabca	down	14+08:34						

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a -1

Result:

Ruptime displays a list similar to the following:

hpabcc	up	1+17:40,	6	users,	load	0.18,	0.13,	0.09
hpabcd	up	14+06:49,	3	users,	load	0.10,	0.38,	0.49
hpabcb	down	1:13						
hpabca	down	14+08:34						

4-10 Displaying Ruptime Status Lines

Sorted by Increasing Load Average

Excluding Idle Users

At your HP-UX prompt, enter:

ruptime -r -1

Result:

Ruptime displays a list similar to the following:

hpabca	down	14+08:34					
hpabcb	down	1:13					
hpabcd	up	14+06:49,	3 users,	load	0.10,	0.38,	0.49
hpabcc	up	1+17:40,	4 users,	load	0.18,	0.13,	0.09

Including Idle Users

At your HP-UX prompt, enter:

ruptime -a -r -1

Result:

Ruptime displays a list similar to the following:

hpabca	down 14+08:34	
hpabcb	down 1:13	
hpabcd	up 14+06:49,	3 users, load 0.10, 0.38, 0.49
hpabcc	up 1+17:40,	6 users, load 0.18, 0.13, 0.09

4-12 Displaying Ruptime Status Lines

Listing Users with Rwho

Rwho is a Berkeley Service that lists information about HP-UX or UNIX hosts on the network. The information that *rwho* displays includes:

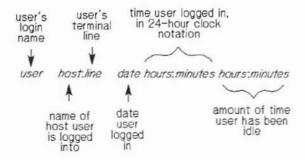
- the login name of each user who is logged into a host on the network,
- the name of the host each user is logged into,
- each user's terminal line,
- the date and time each user logged in, and
- the amount of time (if any) each user has been idle (has not used the system for one minute or more).

This information is useful in determining:

- who is logged into the hosts on the network and
- who is likely to be at their terminal or workstation.

Using Rwho

For each user logged into a network host, *rwho* displays an information line with the following format:



Rwho Status Line

Rwho reports information in the following way:

If	Then rwho		
a user has not used the system for one to fifty-nine minutes,	reports the amount of idle time for that user.		
a user has not used the system for one hour or more,	omits the user from its list, unless you invoke <i>rwho</i> with the -a option.		
a host has not broadcast an <i>rwho</i> status report to the network lately,	assumes that the host is down and does not list any users on that host.		

Note

Rwho's list of users can become excessively long when the number of users on network hosts becomes large.

With rwho, you can list either:

- users on network hosts who are active or who have been idle for less than one hour or
- all users logged into network hosts, regardless of the amount of time any of them have been idle.

Listing Active and Likely Active Users of Network Hosts

At your HP-UX prompt, enter:

rwho

Result:

Rwho displays a list similar to the following:

acb	hpabcd:ttyp3	Jun	2 08:32 :19
bjt	hpabcf:tty3p3	Jun	2 09:35 · Active User
chas	hpabcd:tty3p3	Jun	2 07:47 :27
cjc	hpabcd:tty1p2	Jun	2 07:55 - Active User
dae	hpabcf:ttyp2	Jun	2 08:28 :57

Listing All Users of Network Hosts

At your HP-UX prompt, enter:

rwho -a

Result:

Rwho displays a list similar to the following:

acb	hpabcd:ttyp3	Jun	2 08:32	:19	
bjt	hpabcf:tty3p3	Jun	2 09:35	•	Active User
chas	hpabcd:tty3p3	Jun	2 07:47	:27	
cjc	hpabcd:tty1p2	Jun	2 07:55	•	Active User
dae	hpabcf:ttyp2	Jun	2 08:28	:57	
gen	hpabcd:ttyp4	Jun	2 08:45	5:59	
kgj	hpabcd:ttyp0	Jun	2 08:09	1:02	
scb	hpabce:tty3p1	Jun	2 12:12	3:24	

Logging into a Host with Telnet

Telnet is an ARPA Service that you use to log into a remote HP-UX, UNIX, or non-UNIX host. *Telnet* is the virtual terminal program that uses the ARPA standard TELNET protocol.

Using Telnet

To use telnet you:

- invoke telnet,
- change the *telnet* escape character if necessary,
- connect to a remote host, and then
- log into that host.

After that, you can use *telnet* to do work on the remote host as if your terminal or workstation were physically connected to that host.

1. Invoke Telnet

At your HP-UX prompt, enter:

telnet

Result:

Telnet displays its prompt:

telnet>

Telnet's Two States

Telnet has two states: input state and command state. The telnet> prompt means that *telnet* is in its command state. *Telnet*'s command state allows you to execute individual *telnet* commands to get help, get status information, change characteristics of your *telnet* session or exit from *telnet*.

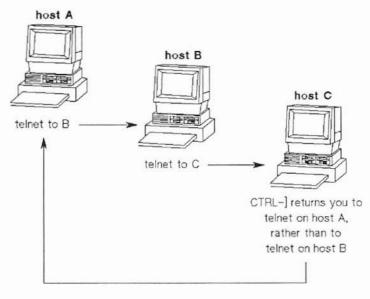
Connecting to and logging into a remote host puts *telnet* into its input state. Everything you type, with the exception of the *telnet* escape character, goes to the remote host. When *telnet* is in its input state, you can do work on the remote host as if it were your local host.

When *telnet* is in its input state, entering the *telnet* escape character puts *telnet* back into its command state. After you execute a *telnet* command, *telnet* returns to its input state.

2. Change the Telnet Escape Character If Necessary

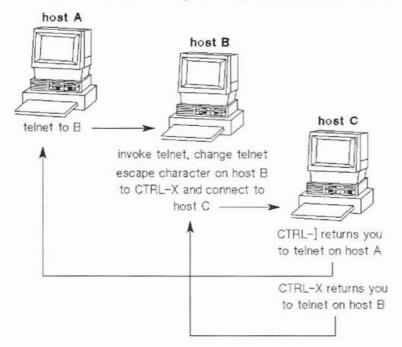
When you connect to a remote host, *telnet* presets its escape character to **CTRL-]** (sometimes shown as ^]). You need to set the *telnet* escape character to something else if:

- the character **CTRL-]** performs a particular function for a program within which you are running *telnet*,
- that character **CTRL-]** performs a particular function for a program you plan to run from within *telnet*, or
- you want to be able to return to an intermediate remote host when you nest a series of *telnet* commands over a chain of remote hosts. This last case is illustrated below:



Nested Telnet Commands

If you want to return to *telnet* on host B, you must change the escape character for *telnet* on host B when you invoke *telnet* to connect to host C:



Nested Telnet Commands with Different Escape Characters

You can change the *telnet* escape character with *telnet*'s escape command.

Caution: What Not to Change the Telnet Escape Character To

Do not change the *telnet* escape character to:

- a character with a particular function in a program you may run from within *telnet* or
- a character that may conflict with your terminal configuration. HP-UX associates certain characters with specific functions. These characters are listed in the following table.

Function	Predefined Character in HP-UX		
	ers, so they may be different in your l configuration:		
End of File	CTRL-D (EOT)		
Interrupt	CTRL-C (DEL or Rubout)		
Quit	CTRL-\ (FS)		
Erase	CTRL-H (#)		
Kill	CTRL-U (@)		
End of Line	CTRL-@ (NUL)		
Shell Layers Switch	CTRL-Z		
You can not change these char	acters in your terminal configuration:		
New Line/Line Feed	CTRL-J (LF)		
Stop or XOFF (Transmit Off)	CTRL-S (DC3)		
Start or XON (Transmit On)	CTRL-Q (DC1)		
Enquire	CTRL-E (ENQ)		
Acknowledge	CTRL-F (ACK)		

To find out what characters are used in your terminal configuration, at your HP-UX prompt, enter:

stty –a

Changing the Telnet Escape Character

At the telnet > prompt, enter:

escape new_telnet_escape_character

where *new_telnet_escape_character* is the character you want to change the *telnet* escape character to.

Note

If *new_telnet_escape_character* is a control character, you must enter it as character where represents **CTRL**. (For example, you would enter **CTRL–X** as X .)

Example Entry:

escape ^[

Result:

Telnet changes its escape character to the character you specify and displays:

Escape character is 'new_telnet_escape_character'.

If you are **not** connected to a remote host, *telnet* then redisplays its telnet> prompt.

If you are connected to a remote host, *telnet* returns you to the remote host.

Note

When connected to a remote host, press **Return** to redisplay the remote host's prompt.

3. Connect to a Remote Host

At the telnet > prompt, enter:

open remote_host

where remote host is the name or alias of a host listed in /etc/hosts.

Note

The file *letc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

To connect to a host **not** listed in */etc/hosts*, you can give the host's internet address where *remote_host* appears above. The internet address must be in dot notation (for example, 192.6.21.9).

Example Entry:

open hpabsa

Result:

Telnet connects you to the remote host and prompts for a remote login name, displaying:

Trying... Connected to remote_host. Escape character is '^]'. remote_host_identification_message remote host login prompt

4. Log into the Remote Host

Note

You must be able to supply *telnet* with a valid login name and password (if required) on the remote host.

Log into the remote host by supplying a valid remote login name and password, if required. The login name and password you supply may be yours or someone else's.

Result:

If the login name and password you supplied are valid on the remote host, the remote host logs you in and displays its login message and its prompt.

Giving Telnet Commands When Telnet Is in Its Input State

To give a command to *telnet* from its input state, you need to enter the *telnet* escape character. This character tells *telnet* that what follows is a *telnet* command, **not** information you are sending to the remote host.

The preset *telnet* escape character is **CTRL-]** (sometimes shown as ^]). (You may have changed this to something else, as explained earlier in this chapter.)

When you enter the *telnet* escape character, *telnet* responds with its telnet> prompt. The telnet> prompt lets you know that *telnet* is ready to accept a command. (*Telnet* is in its command state.)

You can execute only one *telnet* command at a time; *telnet* returns to its input state after each command completes.

The remaining sections of this chapter discuss what *telnet* commands are available in addition to *escape* and *open*.

Checking the Behavior of Carriage Returns from a Remote Host

When some remote hosts send a carriage return to your local host, your local host may need to change the carriage return into a carriage returnline feed combination. *Telnet's crmod* command allows you to enable or disable this behavior.

The following behavior indicates that *telnet*'s carriage return mode setting is wrong for the type of remote host you are communicating with:

- If pressing **Return** produces double-spaced lines (indicating an extra line feed), you need to **disable** carriage return mode.
- If pressing **Return** moves the cursor to the beginning of the same line so that the same line keeps getting overwritten (indicating no line feed), you need to **enable** carriage return mode.

Changing the Carriage Return Mode Setting

1. If you are not already at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

crmod

Result:

If carriage return mode was on, telnet turns it off and displays:

Wont map carriage return on output.

If carriage return mode was off, telnet turns it on and displays:

Will map carriage return on output.

If you are connected to a remote host, *telnet* returns you to the remote host.

Note

To redisplay the remote host's prompt, press Return.

If you are **not** connected to a remote host, *telnet* redisplays its telnet> prompt.

Disconnecting from a Remote Host and/or Exiting Telnet

You can disconnect from a remote host but remain in *telnet* if you connected to the remote host with *telnet*'s *open* command. This is useful if you want to connect to other remote hosts during the same *telnet* session.

If you want to exit from *telnet* and return to HP-UX on your local host, there are two ways to do so, depending on which state (input or command) *telnet* is in. Exiting from *telnet* disconnects from a remote host if a connection exists.

Disconnecting from a Remote Host and Remaining in Telnet

Note

This is possible only if you connected to the remote host with *telnet*'s *open* command.

1. If you are not at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

close

Result:

If you connected to the remote host with *telnet*'s *open* command, *telnet* disconnects from the remote host and returns the telnet> prompt, displaying:

Connection closed.

telnet>

If you connected to the remote host when you invoked *telnet* (as explained later in this chapter), *telnet* disconnects from the remote host and returns you to HP-UX on your local host. (You exit from *telnet*.) *Telnet* displays:

Connection closed.

6-12 Disconnecting from a Remote Host and/or Exiting Telnet

local_HP-UX_prompt

If no connection exists to a remote host, the *close* command has no effect. *Telnet* just redisplays its prompt:

telnet>

Exiting from Telnet When Telnet Is in Its Input State

Log out of the remote host as you normally would (such as with CTRL-D).

Result:

Telnet disconnects from the remote host and returns you to HP-UX on your local host, displaying:

Connection closed by foreign host.

local_HP-UX_prompt

Exiting from Telnet When Telnet Is in Its Command State

At the telnet > prompt, enter:

quit

Result:

Telnet disconnects from the remote host and returns you to HP-UX on your local host, displaying:

Connection closed.

local_HP-UX_prompt

Obtaining Help

You can obtain summary information about *telnet* commands with *telnet*'s ? command. You can either list the *telnet* commands or get information about a specific *telnet* command.

Listing the Telnet Commands

1. If you are not at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

?

Result:

Telnet lists its commands, displaying:

Commands may be abbreviated. Commands are: connect to a site open close close current connection exit telnet auit set escape character escape print status information status options toggle viewing of options processing toggle mapping of received carriage returns crmod debug toggle debugging shell escape ! ? print help information

telnet>

Note

If you were connected to a remote host and want to redisplay its prompt, press **Return** twice.

Getting Information about a Specific Telnet Command

1. If you are not already at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

? telnet_command

Example Entry:

? open

Result:

Telnet displays:

brief_description_of_command

telnet>

Note

If you were connected to a remote host and want to redisplay its prompt, press **Return** twice.

Temporarily Returning to HP-UX on Your Local Host

From within *telnet*, you can temporarily invoke a local HP-UX shell. This is a new shell, descended from the one started when you logged into your local host. This allows you to work on your local host and then return to *telnet*.

You can either:

- execute a single HP-UX command on your local host and automatically return to *telnet* or
- work on your local host for as long as you need to before you return to *telnet*.

Executing a Single HP-UX Command on Your Local Host

1. If you are not already at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

! HP-UX_command

where HP-UX_command is an HP-UX command line.

Example Entries:

! pwd

! hostname

Result:

A local HP-UX shell executes the command and returns you to the remote host, displaying:

[Returning to remote]

Note

To redisplay the remote host's prompt, press Return.

Working for an Extended Time on Your Local Host

1. If you are not already at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

!

Result:

Telnet gives you a local HP-UX shell to work in, displaying:

local_HP-UX_prompt

3. Enter HP-UX commands.

Result:

The local shell executes each command you enter and then redisplays the local HP-UX prompt.

4. Exit the local HP-UX shell as you normally would (such as with CTRL-D).

Result:

The local shell returns you to the remote host, displaying:

[Returning to remote]

Note

To redisplay the remote host's prompt, press Return.

Obtaining Telnet Status

You can display:

- whether or not a connection to a remote host exists and the name of the host to which a connection exists (if any),
- the current status of echo and mode (if a connection exists), and
- the current telnet escape character.
- 1. If you are not already at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

status

Result:

Telnet displays its status information. If you are connected to a remote host, *telnet* returns you to the remote host.

Note

To redisplay the remote host's prompt, press Return.

If you are not connected to a remote host, *telnet* redisplays its telnet> prompt.

Changing Where User Input Is Echoed

You can choose whether user input is echoed locally or remotely. In local echo, user input is echoed to the terminal by the local *telnet* before being transmitted to the remote host. In remote echo, the remote host echoes user input. By default, *telnet* starts a connection in local echo, and requests that the TELNET server do remote echo. If the server refuses the request, you will see an error message. You can check the status of *echo* with the *telnet status* command.

Local echo produces less network traffic than remote echo, because the server need not transmit user input back to the local system, and will transmit only the output of the remote application. When communication between the local and remote systems is slow, local echo will appear to provide better system response. However, note that remote applications that expect to handle echoing of user input themselves, such as csh(1), ksh(1), and vi(1), will not work correctly with local echo.

1. If you are not at the telnet > prompt, enter the *telnet* escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

echo local

or

echo remote

Result:

Telnet now echoes input locally or remotely, depending on which you specified.

Changing User Input Mode

You can set *telnet*'s user input mode to **character** or **line**. In **character** mode, *telnet* sends each character to the remote host as it is typed. In **line** mode, *telnet* gathers user input into lines and transmits each line to the remote host when the user types a carriage return, linefeed, or EOF. By default, *telnet* uses character mode. Note that setting line mode also sets local echo. You can check the status of *mode* with *telnet*'s *status* command.

In line mode, *telnet* transmits fewer packets over the network than it does in character mode, as it sends a packet only when the user terminates a line rather than sending each character in its own packet. This is particularly useful if you are connecting over some X.25 networks that charge users on a per-packet basis. However, note that remote applications that expect to interpret user input character by character, such as more(1), csh(1), ksh(1), and vi(1), will not work correctly in line mode.

1. If you are not already at the *telnet* > prompt, enter the telnet escape character.

Result:

Telnet displays its prompt:

telnet>

2. At the telnet > prompt, enter:

mode character

or

modeline

Result:

Telnet sets user input mode to character or line, depending on which you specified.

Connecting to a Remote Host When You Invoke Telnet

The *telnet* command you give from your local HP-UX prompt can take the following form:

telnet remote host

Specifying a remote host's name or alias (as listed in your local host's */etc/hosts* file) on the *telnet* command line causes *telnet* to connect to that remote host without your having to use *telnet*'s *open* command.

If you connect to a remote host in this way, *telnet*'s *close* command exits from *telnet*, instead of disconnecting from the remote host and remaining in *telnet*. Therefore, you can **not** connect to other remote hosts during the same *telnet* session. (You must reinvoke *telnet* to connect to another host.)

6-24 Connecting to a Remote Host When You Invoke Telnet

7

Logging into a Host with Rlogin

Rlogin is a Berkeley Service that you use to log into a remote HP-UX or UNIX host from your local host.

Before you use *rlogin*, you should:

- determine if you need to change the *rlogin* escape character and
- determine what size characters to send using rlogin.

You should determine these things ahead of time because:

- they can affect whether *rlogin* operates properly and communicates properly with a remote host, and
- you can change the settings associated with these only when you invoke *rlogin*.

Determining If You Need to Change the Riogin Escape Character

The *rlogin* escape character, when combined with other particular characters, allows you to exit *rlogin* and allows you to temporarily return to HP-UX on your local host.

Rlogin presets its escape character to a tilde (\sim). You need to set the *rlogin* escape character to something else if:

- that character performs a function for a program you are running now on your local host, and
- you plan to run *rlogin* from within that local program.

Otherwise when you enter the character, the program you are currently running will respond to it, instead of *rlogin*.

Caution: What Not to Change the Rlogin Escape Character To

If you must change the *rlogin* escape character, do **not** change it to a character that may conflict with your terminal configuration. HP-UX associates certain characters with specific functions. These characters are listed below:

Function	Predefined Character in HP-UX
	aracters, so they may be different in your minal configuration:
End of File	CTRL-D (EOT)
Interrupt	CTRL-C (DEL or Rubout)
Quit	CTRL-\ (FS)
Erase	CTRL-H (#)
Kill	CTRL-U (@)
End of Line	CTRL-@ (NUL)
Shell Layers Switch	CTRL-Z

Function

Predefined Character in HP-UX

You can not change these characters in your terminal configuration:

New Line/Line Feed	CTRL-J (LF)	
Stop or XOFF (Transmit Off)	CTRL-S (DC3)	
Start or XON (Transmit On)	CTRL-Q (DC1)	
Enquire	CTRL-E (ENQ)	
Acknowledge	CTRL-F (ACK)	

To find out what characters are used in your terminal configuration, at your HP-UX prompt, enter:

stty -a

Also, do **not** change the *rlogin* escape character to a period (.) or an exclamation mark (!), because you combine these characters with the *rlogin* escape character either to exit *rlogin* or to return temporarily to HP-UX on your local host.

Note that it is inconvenient to change the *rlogin* escape character to one that you may frequently enter at the beginning of lines for a program running on a remote host. (This is because you would need to enter the character twice to allow the program to respond to the character, instead of *rlogin*.)

You change the *rlogin* escape character by invoking *rlogin* with the -e option, as described after the next section of this chapter.

Determining What Size Characters to Send with Rlogin

Before you run *rlogin*, you must determine what size characters to send with *rlogin*. You must decide this ahead of time because:

- the character size can affect whether you can communicate properly with a remote host you login into with *rlogin*, and
- you can change the character size only when you invoke rlogin.

Rlogin sends eight-bit characters to a remote host unless you tell *rlogin* to send seven-bit characters instead.

In general, send seven-bit characters with *rlogin* if sending eight-bit characters with *rlogin* (the preset behavior) causes problems communicating with a remote host.

When You Can Send Eight-Bit Characters

For communication between your local host and a remote host to work properly with eight-bit characters, all of the following must be configured for eight-bit characters:

- the remote host's tty driver,
- the local host's tty driver, and
- your local terminal hardware.

The following instructions tell how to check if these are configured for eight-bit characters.

- 1. Check whether the remote host is configured for, and can support, eight-bit characters. If the remote host runs HP-UX, you can do this by performing the same steps on the remote host as shown below for checking this on your local host.
- **2.** Check whether your local host is configured for eight-bit characters as follows:

a. At your local HP-UX prompt, enter:

stty –a

Result:

The command displays all of your local host's tty driver settings.

b. Check the output for the setting:

cs8

This means that the character size is set to eight bits.

c. If the character size is not set to eight bits, at your local HP-UX prompt, enter:

stty cs8

Result:

This sets the character size for your local host's tty driver to eight bits.

3. Check whether your local terminal hardware is configured for eight-bit characters in its configuration menu or in its switch settings.

When You Must Send Seven-Bit Characters

You must tell rlogin to send seven-bit, instead of eight-bit, characters if:

- you can not configure your local terminal hardware to send eight-bit characters or seven-bit characters with high bit 0 (null parity),
- you can not configure a remote host to receive eight-bit characters (if it is not already configured to do so), or
- you might send eight-bit characters that a remote host interprets differently than your local host would. This can cause unpredictable results.

To send seven-bit characters with *rlogin*, you invoke *rlogin* with the -7 option, as explained in the next section of this chapter.

Using Rlogin

If you have an account on a remote host, then with rlogin, you can either:

- log into the remote host automatically (without supplying your remote login name and password) if the remote host is configured to allow this or
- log into the remote host manually by supplying your remote login name and password.

The following sections explain each of these options for logging into a remote host with *rlogin*.

Automatic Login

Rlogin allows you to log into a remote host without supplying your remote login name and password if the remote host is configured in either of two ways:

Either:

- you must have an account on the remote host with the same login name as your local login name, and
- the name of your local host must be in the remote host's /etc/hosts.equiv file,

or:

- you must have an account on the remote host, and
- the name of your local host and your local login name must be in a *.rhosts* file in your home directory on the remote host.

The next section explains how to create a remote *\$HOME/.rhosts* file for yourself, if you need to do so. Otherwise, skip the next section.

Creating a \$HOME/.rhosts File on a Remote Host

If you have an account on a remote host, you can set up the account so that you can log into the remote host without having to supply your remote login name and password. To do this, you create a file named *.rhosts* in your remote home directory. You can find out what your remote home directory is by entering:

echo \$HOME

on the remote host. You must place the name of your local host and your local login name in the *.rhosts* file you create.

Caution

A *\$HOME*/.*rhosts* file creates a significant security risk. Be sure to follow the directions below for "Protecting Your .rhosts File."

The entry you place in your .rhosts file must have the following format:

your_local_host's_name your_local_login_name

You can separate your_local_host's_name and your_local_login_name with any number of tabs or spaces. Put any comments after your_local_login_name.

Example \$HOME/.rhosts File Entry

If your local host's name were *hpabsa* and your local login name were *richard*, on the remote host you would create a *\$HOME/.rhosts* file with the following entry:

hpabsa richard

Protecting Your \$HOME/.rhosts File

It is important to protect your remote *.rhosts* file and home directory to prevent unauthorized users from gaining *rlogin* access to your remote account and host. Only you should be able to create a *.rhosts* file in your remote home directory and write entries to the file. To do so:

- 1. Insure that your remote .rhosts file is owned by you, the user.
- 2. Use the HP-UX or UNIX *chmod* command to protect your remote *.rhosts* file with 0400 (-r-----) permission.
- 3. Use the HP-UX or UNIX *chmod* command to protect your remote home directory so that no one else can read it or write to it. For example, you should protect your remote home directory with at least 0711 (-rwx--x--x) permission.

Logging into the Remote Host Automatically

At your HP-UX prompt, enter:

rlogin remote_host [-e character] [-7]

where:

• remote_host is the name or alias of a host listed in /etc/hosts.

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple aliases. The *letc/hosts* file may contain comments and other information as well.

- the brackets ([]) mean that the enclosed option is optional. Omitting the -e option sets the *rlogin* escape character to a tilde (~), and omitting the -7 option sets the character size to eight bits.
- *character* is the character you want to change the *rlogin* escape character to. If you want to enter a control character for the *rlogin* escape character, hold down **CTRL** while pressing another character key. (Control characters are not displayed.)
- -7 is an option that sets the character size to seven bits, with the eighth bit set to zero.

Example Entries:

rlogin hpabsb

rlogin hpabsb-7

rlogin hpabsb -e=

rlogin hpabsb -e= -7

Result:

The remote host logs you in, displaying:

remote_host's_login_message

remote_host_prompt

Note

Rlogin does not send **Break** to the remote host. Therefore, you cannot **Break** out of a program on the remote host when you are logged into the remote host with *rlogin*.

If you are now logged into the remote host, skip to the section called "If You Get Unexpected Results after Logging Into a Remote Host."

Manual Login

You must log into a remote host manually by supplying your remote login name and password if the remote host is not configured to allow automatic login.

Logging into the Remote Host Manually

1. At your HP-UX prompt, enter:

rlogin remote_host [-e character] [-7] -1 remote_login_name

where:

• remote_host is the name or alias of a host listed in /etc/hosts.

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple aliases. The *letc/hosts* file may contain comments and other information as well.

- the brackets ([]) mean that the enclosed option is optional. Omitting the -e option sets the *rlogin* escape character to a tilde (~), and omitting the -7 option sets the character size to eight bits.
- *character* is the character you want to change the *rlogin* escape character to. If you want to enter a control character for the *rlogin* escape character, hold down **CTRL** while pressing another character key. (Control characters are not displayed.)

- -7 is an option that sets the character size to seven bits, with the eighth bit set to zero.
- remote login name is your login name on the remote host.

Example Entries:

rlogin hpabsb –1 peter

rlogin hpabsb -7 -1 peter

rlogin hpabsb -e= -1 peter

rlogin hpabsb -e= -7 -l peter

Result:

The remote host prompts you for your remote password, displaying:

Password:

2. Enter your remote password.

Result:

The remote host logs you in, displaying:

remote_host's_login_message

remote_host_prompt

Note

Rlogin does not send **Break** to the remote host. Therefore, you cannot **Break** out of a program on the remote host when you are logged into the remote host with *rlogin*.

If You Get Unexpected Results after Logging into a Remote Host

The values set in the remote host's *letc/profile* file (for *sh* users) or *letc/csh.login* file (for *csh* users) may not match those you are accustomed to on your local host. For example, the terminal type or erase character on the remote host may be different from what you have set up on your local host. Therefore, you may get unexpected results while working on the remote host.

To set values on the remote host to match the values you are accustomed to on the local host, create or edit your own *\$HOME/.profile* file (if you use *sh*) or *\$HOME/.login* file (if you use *csh*) on the remote host. An easy way to do this is to copy your local *\$HOME/.profile* file or *\$HOME/.login* file to your home directory on the remote host.

The values in your own remote *.profile* file or *.login* file take precedence over the values in the remote host's default */etc/profile* or */etc/csh.login* file.

Logging Out of the Remote Host and Exiting Rlogin

There are two ways to log out of the remote host and exit rlogin:

• Log out of the remote host as you normally would (such as with CTRL-D).

Result:

Rlogin logs you out of the remote host, disconnects from the remote host and returns you to HP-UX on your local host, displaying:

Connection closed.

local_HP-UX_prompt

• At the beginning of a new line, enter:

rlogin_escape_character.

(That is, the *rlogin* escape character followed by a period.)

Note

Rlogin does not display its escape character until you enter the period, or second character. *Rlogin* recognizes its escape character only at the beginning of a new line.

Example Entry:

~ .

Logging into a Host with Rlogin 7-15

Result:

Rlogin logs you out of the remote host, disconnects from the remote host and returns you to HP-UX on your local host, displaying:

Closed connection.

local_HP-UX_prompt

Temporarily Returning to HP-UX on Your Local Host

From within *rlogin*, you can temporarily invoke a local HP-UX shell. This is a new shell, descended from the one started when you logged into your local host. This allows you to work on your local host and then return to *rlogin*.

You can either:

- execute a single HP-UX command on your local host and automatically return to *rlogin* or
- work on your local host for as long as you need to before you return to *rlogin.*

Executing a Single HP-UX Command on Your Local Host

At the beginning of a new line, enter:

rlogin_escape_character! HP-UX_command

Note

Rlogin does not display its escape character until you enter the exclamation mark, or second character. *Rlogin* recognizes its escape character only at the beginning of a new line.

Example Entries:

~! pwd

~! hostname

Result:

A local HP-UX shell executes the command and returns you to the remote host, displaying:

[Returning to remote]

Note

To redisplay the remote host's prompt, press Return.

Logging into a Host with Rlogin 7-17

Working for an Extended Time on Your Local Host

1. At the beginning of a new line, enter:

rlogin_escape_character!

Note

Rlogin does not display its escape character until you enter the exclamation mark, or second character. *Rlogin* recognizes its escape character only at the beginning of a new line.

Example Entry:

~!

Result:

Rlogin gives you a local HP-UX shell to work in, displaying:

local_HP-UX_prompt

2. Enter HP-UX commands.

Result:

The local shell executes each command you enter and then redisplays the local HP-UX prompt.

3. Exit the local HP-UX shell as you normally would (such as with CTRL-D).

Result:

The local shell returns you to the remote host, displaying:

[Returning to remote]

Note

To redisplay the remote host's prompt, press Return.

Passing the Rlogin Escape Character to a Remote Program

If you are running a program on the remote host, and you want to send the program the same character as the *rlogin* escape character, you can always do this as long as you do not need to enter the character at the beginning of a new line.

However, if you need to enter the character at the beginning of a new line, you must enter the character twice because *rlogin* reads the first one as its own escape character. If this is frequently necessary, you may want to change the *rlogin* escape character.

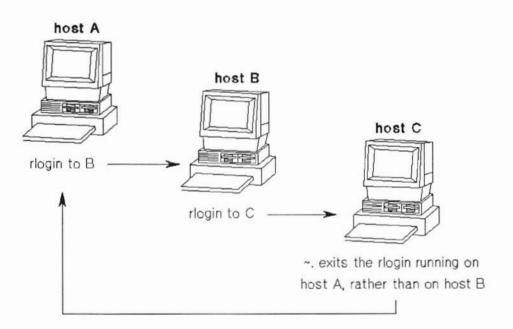
For example, suppose that you have logged into a remote host with *rlogin*, and you run the *vi* editor on the remote host from within *rlogin*. The current *rlogin* escape character would be a tilde (\sim) if you did not change it. Normally, if you wanted to capitalize a character at the beginning of a new line with *vi*, you would position the cursor at that character and enter:

However, when you are running vi from within *rlogin*, your local *rlogin* interprets the tilde (~) first, causing it not to have any effect in the remote vi. To accomplish the capitalization at the beginning of a new line, you must enter:

~~

instead, so that *rlogin* understands that you want to pass a tilde (\sim) to the remote host.

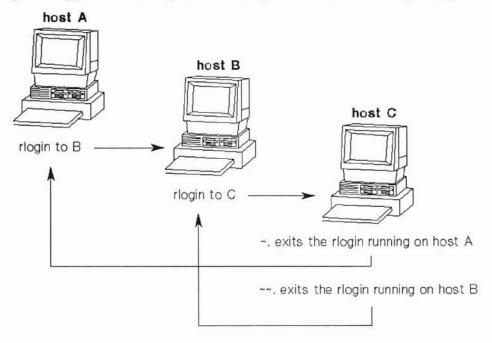
Another example of when you would need to enter the *rlogin* escape character more than once at the beginning of a new line is if you nested a series of *rlogin* commands over a chain of remote hosts. This case is illustrated below:



If you want to exit to host B, you must enter:

~~.

The first tilde (\sim) is interpreted by *rlogin* on host A, while the second one is passed to the program run from within *rlogin* on host B. That program happens to be *rlogin*, which responds to its exit escape sequence.



As a final example, if you wanted to invoke an interactive shell on host B instead of host A, you would enter:

~~!

(instead of \sim !) at the beginning of a new line.

Logging into a Remote Host as Someone Else

You can use *rlogin* to log into a remote host as someone else if you know that user's login name on the remote host, and either of the following two requirements are met:

Either that user must have your local host name and local login name in a *.rhosts* file in his or her home directory on the remote host,

or you must know that user's password on the remote host.

Note

If the remote user's account has no password, you can use *rlogin* to log into the remote host as that user without meeting either of the requirements above.

1. At your HP-UX prompt, enter:

rlogin remote_host [-e character] [-7] -1 remote_login_name

where:

- remote_host is the name or alias of a host listed in /etc/hosts,
- the brackets ([]) mean that the enclosed option is optional,
- *character* is the character you may want to change the *rlogin* escape character to,
- -7 is an option that sets the character size to seven bits, with the eighth bit set to zero, and
- *remote_login_name* is the login name of the remote user who you want to log into the remote host as.

Example Entries:

rlogin hpabsb -l alan

rlogin hpabsb -7 -1 alan

rlogin hpabsb -e= -l alan

rlogin hpabsb -e= -7 -1 alan

Result:

If the remote user has your local host name and local login name in his or her *\$HOME/.rhosts* file on the remote host, then the remote host logs you in and displays:

remote_host's_login_message

remote_host_prompt

Otherwise, the remote host prompts you for the remote user's password, displaying:

Password:

2. If the Password: prompt is displayed, enter the remote user's password.

Result:

The remote host logs you in and displays:

remote_host's_login_message

remote_host_prompt

Giving Other Remote Users Rlogin Access to Your Local Account

You can give remote users *rlogin* access to your local account by creating a *.rhosts* file. You place remote users' host names and login names in this file so that *rlogin* lets them log into your local host as you.

Caution

A *\$HOME*/.*rhosts* file creates a significant security risk. Be sure to follow the instructions below on "Protecting Your .rhosts File."

Your *.rhosts* file must be in your home directory on your local host. You can find out what your local home directory is by entering:

echo \$HOME

on your local host.

Each entry you place in your .rhosts file must have the following format:

remote_host name remote login name

Follow these rules when creating a .rhosts file:

- Each entry must contain a valid remote host name and remote login name.
- Separate the host name and login name with any number of tabs or blanks.
- Put any comments after the login name in any entry.

Example .rhosts File Entry

If you wanted to give user *cdm* on remote host *hpabsc rlogin* access to your local account, you would create a *\$HOME/.rhosts* file on your local host with the following entry:

hpabsc cdm

Protecting Your .rhosts File

It is important to protect your *.rhosts* file and your local home directory to prevent unauthorized users from gaining *rlogin* access to your local account. Only you should be able to create a *.rhosts* file in your home directory and write entries to it. To do this:

- 1. Insure that your .rhosts file is owned by you, the user.
- 2. Use the HP-UX *chmod* command to protect your *.rhosts* file with 0400 (-r-----) permission.
- **3.** Use the HP-UX *chmod* command to protect your local home directory so that no one else can read it or write to it. For example, you should protect your local home directory with at least 0711 (-rwx--x--x) permission.
- 4. Insure that your account has a password. Otherwise, remote users can log into your local host (with your login name) as you.

Using Rlogin's "Shorthand" Syntax

If your local host is configured properly, you can enter an *rlogin* command line without the *rlogin* command. That is, an *rlogin* command line can start with the name of a remote host, omitting the *rlogin* command.

You can use this shorthand syntax only if two conditions hold true. You can omit the command *rlogin* from the *rlogin* command line if:

- you add the path */usr/hosts* to your command search path in your *.login*, *.cshrc*, or *.profile* file. Which file contains your *\$PATH* variable depends on which shell you use.
- the super-user or node manager has linked */usr/bin/remsh* to */usr/hosts/host*, where *host* is the name or alias of a remote host (listed in */etc/hosts*) on which you want to execute a command.

Note

Remsh knows whether you mean to invoke *rlogin* or *remsh* by parsing what you give on the command line.

To find out which hosts you can use *rlogin*'s shorthand syntax for, list the contents of the directory */usr/hosts*.

Transferring Files with Ftp

Ftp is an ARPA Service that allows you to transfer files among HP-UX, UNIX, and non-UNIX network hosts that support ARPA Services. *Ftp* is the file transfer program that uses the ARPA standard File Transfer Protocol (FTP).

Ftp has a one-line command syntax. This service not only allows you to perform file transfers, but also file management operations such as changing, listing, creating, and deleting remote directories.

Using Ftp

To use ftp, you:

- invoke ftp,
- choose whether or not to display responses from any remote host you connect to,
- connect to a remote host, and then
- log into that host.

After that, you can use *ftp* to perform file management operations and file transfers.

1. Invoke Ftp

At your HP-UX prompt, enter:

ftp

Result:

Ftp displays its prompt:

ftp>

2. Choose Whether to Display Responses from a Remote Host

Ftp can display all responses from any remote host you connect to. These responses tell you whether *ftp* commands completed successfully. This feature is called **verbose mode**.

You can also choose **not** to have *ftp* display all responses from the remote host so that in most cases, on completing a command, *ftp* just returns its ftp> prompt. Even if verbose mode is off, when you change one of *ftp*'s settings, *ftp* displays the resulting state of the setting.

If *ftp*'s input comes from your keyboard (HP-UX terminal), *ftp* initially has verbose mode on. (This is the usual setting.) Otherwise, *ftp* has verbose mode off (for example, if *ftp*'s input is coming from a file).

Note

This chapter shows both verbose and non-verbose *ftp* responses.

Turning Verbose Mode On or Off

At the ftp > prompt, enter:

verbose

Result:

If verbose mode was on, *ftp* turns it off and displays:

Verbose mode off.

ftp>

If verbose mode was off, *ftp* turns it on and displays:

Verbose mode on.

ftp>

3. Connect to a Remote Host

At the ftp > prompt, enter:

open remote_host

where remote host is the name or alias of a host listed in /etc/hosts.

Note

The file *letc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias...

The ellipsis (...) means that a host may have multiple *aliases*. The *letc/hosts* file may contain comments and other information as well.

To connect to a host **not** listed in */etc/hosts*, you can give the host's internet address where *remote_host* appears above. The internet address must be in dot notation (for example, 192.6.21.9).

Example Entry:

open hpabsa

Result:

Ftp connects you to the remote host and prompts for a remote login name.

• If verbose mode is off, *ftp* displays:

Name (remote_host:remote_login_name):

• If verbose mode is on, *ftp* displays:

Connected to remote_host.

remote_host FTP server...ready.

Name (remote_host:remote_login_name):

4. Log into the Remote Host

Note

For security reasons, you can only log into accounts that have passwords associated with them. You must be able to supply *ftp* with a valid login name and password on the remote host.

1. To log in with the same remote login name as your local login name, press Return at the Name (...): prompt.

To log in with a different remote login name, enter the remote login name at the Name (...): prompt .

Result:

Ftp prompts for the remote password associated with the login name you gave, displaying:

Password (remote_host:remote_login_name):

2. Enter the password associated with the remote login name you gave.

Result:

Ftp logs you into the remote host if the password is valid.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

Password required for remote_login_name.

User remote_login_name logged in.

ftp>

Problems? *Ftp* may also prompt you for an account name, if the remote host you are logging into requires one.

The remote host may be configured to refuse *ftp* connections to specific users for security reasons.

Disconnecting from a Remote Host and Exiting Ftp

You have two options for disconnecting from a remote host:

- You can exit *ftp*. This disconnects from a remote host if a connection exists and returns you to HP-UX.
- You can disconnect from a remote host and remain in *ftp*. This is useful if you want to connect to other remote hosts during the same *ftp* session.

Exiting Ftp to Return to HP-UX on Your Local Host

At the ftp > prompt, enter:

quit

or

bye

Result:

If a connection exists to a remote host, *ftp* disconnects from the remote host and returns you to HP-UX on your local host.

- If verbose mode is off, your local host redisplays its HP-UX prompt.
- If verbose mode is on, ftp displays:

Goodbye.

and your local host redisplays its HP-UX prompt.

If a connection does not exist to a remote host, *ftp* returns you to HP-UX on your local host, and your local host redisplays its HP-UX prompt.

Disconnecting from a Remote Host and Remaining in Ftp

At the ftp > prompt, enter:

close

Result:

Ftp disconnects from the remote host.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

Goodbye.

ftp>

Obtaining Help

You can obtain summary information about *ftp* commands with *ftp*'s *help* command. You can either list the *ftp* commands or get information about a specific *ftp* command.

Listing the Ftp Commands

At the ftp> prompt, enter: help or ? Result: Ftp lists its commands, displaying: Commands may be abbreviated. Commands are: command_list ftp>

Getting Information about a Specific Ftp Command

At the ftp > prompt, enter:

help ftp_command

or

? ftp_command

Result:

Ftp displays:

command

brief_description_of_command

ftp>

Temporarily Returning to HP-UX on Your Local Host

From within *ftp*, you can temporarily invoke a local HP-UX shell. This is a new shell, descended from the one started when you logged into the local host. This allows you to work on the local host and then return to *ftp*.

You can either:

- execute a single HP-UX command on your local host and automatically return to *ftp* or
- work on your local host for as long as you need to before you return to *ftp*.

Executing a Single HP-UX Command on Your Local Host

At the ftp > prompt, enter:

! HP-UX_command

where command is an HP-UX command line.

Example Entries:

- ! pwd
- ! hostname

8-10 Temporarily Returning to HP-UX on Your Local Host

Result:

A local HP-UX shell executes the command and returns you to *ftp* and the ftp> prompt.

Working for an Extended Time on Your Local Host

1. At the ftp > prompt, enter:

!

Result:

Ftp gives you a local HP-UX shell to work in, displaying:

local_HP-UX_prompt

2. Enter HP-UX commands.

Result:

The local shell executes each command you enter and then redisplays the local HP-UX prompt.

3. Exit the local HP-UX shell as you normally would (such as with CTRL-D).

Result:

The local shell returns you to ftp and the ftp> prompt.

How Ftp Treats "Wild Card" Characters, or Metacharacters

You can use csh(1) metacharacters in the directory and file names you specify in *ftp* commands. These metacharacters, or wild card characters, represent a set of characters or character strings and are a "shorthand" way of specifying a set of directory or file names. The following table is a quick reference to the meaning of the csh(1) metacharacters supported by *ftp*:

Character	Matches
*	any string, including a null string
?	any single character
[]	any of the enclosed characters
~	your home directory
~ login_name	login_name's home directory
{,}	any of the enclosed character strings separated by commas

Csh Metacharacters Supported by Ftp

The expansion of metacharacters into the directory and file names they match is called **globbing**. Globbing is on when you first invoke *ftp*. You must turn it off if you need to specify a directory or file name containing one of the characters listed above. That way, *ftp* interprets the character literally, instead of trying to match it to a set of characters. You can turn globbing on or off while in *ftp* at the ftp> prompt.

Turning Globbing On or Off

At the ftp > prompt, enter:

glob

If globbing was on, ftp turns it off and displays:

Globbing off.

ftp>

If globbing was off, *ftp* turns it on and displays:

Globbing on.

ftp>

For some of its commands *ftp* always expands metacharacters, even if globbing is off. If you want a local or remote host to interpret a metacharacter literally when given one of these commands, precede the character with a backslash (\). For example, with some *ftp* commands, you would need to enter a directory named *my*?s as my\?s.

The following table summarizes globbing behavior for applicable *ftp* commands. These commands are discussed in more detail later in this chapter.

Ftp Command	Metacharacters Are
dir	always expanded
ls	always expanded
mdelete	expanded if globbing is on
mdir	expanded if globbing is on
mget	expanded if globbing is on
mls	expanded if globbing is on
mput	expanded if globbing is on

Globbing Behavior for Ftp Commands

In single-file *ftp* commands:

- *ftp* expands metacharacters in a remote file specification only if it begins with the tilde (~) metacharacter. *Ftp* then performs the operation on only the first file that matches the expanded specification.
- *ftp* always expands metacharacters in a local file specification. *Ftp* then performs the operation on only the first file that matches the expanded specification.

Performing Directory Operations with Ftp

From within ftp, you can:

- change local and remote working directories,
- list the contents of remote directories,
- display the name of the remote working directory,
- create a remote directory,
- · delete a remote directory, and
- change the name of a remote directory.

The following sections tell how to perform these directory operations.

Changing the Local Working Directory

To Your Local Home Directory

At the ftp > prompt, enter:

lcd

Result:

Ftp changes the local working directory to your local home directory and displays:

Local directory now your_local_home_directory_path

ftp>

To Another Local Directory

At the ftp > prompt, enter:

lcd local_directory_path

where *local_directory_path* is the full or relative path to a local directory.

Example Entries:

lcd /users/richard/projects

lcd projects

Ftp changes the local working directory to the path you specify and displays:

Local directory now full_local_working_directory_path

ftp>

Changing the Remote Working Directory

At the ftp > prompt, enter:

cd remote_directory_path

where *remote_directory_path* is the full or relative path to a remote directory.

Example Entries:

```
cd /users/lab/richard/xfers
```

cd xfers

Result:

Ftp changes the remote working directory to the path you specify.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, ftp displays:

CWD command okay.

ftp>

8-16 Performing Directory Operations with Ftp

Listing the Contents of the Remote Working Directory

With *ftp*, you can get an extended or abbreviated listing of the contents of the remote working directory. You can send the listing to the display (actually, *stdout*) or to a local file.

To Stdout (Usually the Display)

Note

Metacharacters are always expanded for *ftp*'s *dir* and *ls* commands.

For an extended listing, at the ftp > prompt, enter:

dir

For an abbreviated listing, at the ftp > prompt, enter:

1s

Result:

Ftp lists the contents of the remote working directory to the display (if *stdout* is not redirected).

• If verbose mode is off, *ftp* displays:

```
directory_contents_listing
ftp>
```

• If verbose mode is on, ftp displays:

PORT command okay.

Opening data connection for /bin/ls -1...

directory_contents_listing

Transfer complete.

number bytes received in number seconds...

ftp>

To a Local File

Note

Metacharacters are always expanded for *ftp*'s *dir* and *ls* commands.

For an extended listing, at the ftp > prompt, enter:

dir . local_file_path

For an abbreviated listing, at the ftp > prompt, enter:

1s . local_file_path

where the period (.) represents the remote working directory, and *local_file_path* is the full or relative path to a local file.

8-18 Performing Directory Operations with Ftp

Example Entries:

dir . /users/richard/projects/dirconts

1s . dirconts

Result:

Ftp lists the contents of the remote working directory to the local file you specify.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for /bin/ls -l...

Transfer complete.

number bytes received in number seconds...

Listing the Contents of a Remote Directory

With *ftp*, you can get an extended or abbreviated listing of the contents of a remote working directory. You can send the listing to the display (actually, *stdout*) or to a local file.

To Stdout (Usually the Display)

Note

Metacharacters are always expanded for *ftp*'s *dir* and *ls* commands.

For an extended listing, at the ftp > prompt, enter:

dir remote_directory_path

For an abbreviated listing, at the ftp > prompt, enter:

1s remote_directory_path

where *remote_directory_path* is the full or relative path to a remote directory.

Example Entries:

dir /users/lab/richard/doc

ls doc

Result:

Ftp lists the contents of the remote directory you specify to the display (actually, *stdout*).

8-20 Performing Directory Operations with Ftp

• If verbose mode is off, *ftp* displays:

directory_contents_listing

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for /bin/ls -l...

directory_contents_listing

Transfer complete.

number bytes received in number seconds...

ftp>

To a Local File

Note

Metacharacters are always expanded for *ftp*'s *dir* and *ls* commands.

For an extended listing, at the ftp > prompt, enter:

dir remote_directory_path local_file_path

For an abbreviated listing, at the ftp > prompt, enter:

1s remote_directory_path local_file_path

where *remote_directory_path* is the full or relative path to a remote directory, and *local_file_path* is the full or relative path to a local file.

Example Entries:

dir /users/lab/richard/doc /users/richard/projects/dirconts

ls doc dirconts

Result:

Ftp lists the contents of the remote directory you specify to the local file you specify.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for /bin/ls -1...

Transfer complete.

number bytes received in number seconds...

Listing the Contents of Multiple Remote Directories

With *ftp*, you can get an extended or abbreviated listing of the contents of multiple remote directories. You can send the listing to the display (actually, *stdout*) or to a local file.

To Stdout (Usually the Display)

Note

Metacharacters are expanded for *ftp*'s *mdir* and *mls* commands if globbing is on.

For extended listings, at the ftp > prompt, enter:

mdir remote_directory_path ... -

For abbreviated listings, at the ftp > prompt, enter:

mls remote_directory_path ... -

where *remote_directory_path* is the full or relative path to a remote directory, and – represents *stdout*. The ellipsis (...) means that you can specify multiple *remote_directory_paths*.

Example Entries:

mdir /users/lab/richard /users/lab/richard/doc -

mls doc code code/source -

If interactive mode is on, *ftp* prompts you to verify that *stdout* (the display) is really where you want the contents of the directories listed. *Ftp* displays:

local-file-?

• If this is not what you want to do, enter:

Ν

and *ftp* cancels the directory listings and redisplays its ftp> prompt.

• If this is what you want to do, enter:

Y

and continue with the description below. Interactive mode is explained in detail in a later section of this chapter.

Ftp displays (lists to *stdout*) the contents of the remote directories you specify.

• If verbose mode is off, *ftp* displays: directory_contents_listing directory contents listing . (repeated for each directory listed) ftp> • If verbose mode is on, *ftp* displays: PORT command okay. Opening data connection for /bin/ls... directory contents listing Transfer complete. number bytes received in number seconds... PORT command okay. Opening data connection for /bin/ls... directory contents listing Transfer complete. number bytes received in number seconds... . (repeated for each directory listed) ftp>

To a Local File

With *ftp*, you can send an extended or abbreviated listing of the contents of multiple remote directories to a local file.

Note

Metacharacters are expanded for *ftp*'s *mdir* and *mls* commands if globbing is on.

For extended listings, at the ftp > prompt, enter:

mdir remote_directory_path. local_file_path

For abbreviated listings, at the ftp > prompt, enter:

mls remote_directory_path ... local_file_path

where *remote_directory_path* is the full or relative path to a remote directory, and *local_file_path* is the full or relative path to a local file. The ellipsis (...) means that you can specify multiple *remote_directory_paths*.

Example Entries:

mdir /users/lab/richard /users/lab/richard/doc dirlists

mls doc code code/source dirlists

Result:

If interactive mode is on, *ftp* prompts you to verify that the local file you specified is really where you want the contents of the directories listed. *Ftp* displays:

```
local-file local_file_path ?
```

8-26 Performing Directory Operations with Ftp

• If local_file_path is not correct, enter:

Ν

and *ftp* cancels the directory listings and redisplays its ftp> prompt.

• If local_file_path is correct, enter:

Y

and continue with the description below. Interactive mode is explained in detail in a later section of this chapter.

Ftp lists the contents of the remote directories you specify to the local file you specify.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for /bin/ls...

Transfer complete.

number bytes received in number seconds...

PORT command okay.

Opening data connection for /bin/ls...

Transfer complete.

number bytes received in number seconds...

. (repeated for each directory listed)

ftp>

•

Displaying the Name of the Remote Working Directory

At the ftp > prompt, enter:

pwd

Ftp displays:

"full_remote_working_directory_path" is the current working directory.

ftp>

Creating a Remote Directory

At the ftp > prompt, enter:

mkdir remote_directory_path

where *remote_directory_path* is a full or relative path to a remote directory.

Example Entries:

mkdir /users/lab/richard/temp

mkdir temp

Result:

Ftp creates the remote directory you specify.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

MKDIR command okay.

Deleting a Remote Directory

Note

To delete a remote directory, the directory must be empty.

At the ftp > prompt, enter:

rmdir remote_directory_path

or

delete remote_directory_path

where *remote_directory_path* is a full or relative path to a remote directory.

Example Entries:

rmdir /users/lab/richard/temp

delete temp

Result:

Ftp deletes the remote directory you specify.

8-30 Performing Directory Operations with Ftp

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

RMDIR command okay.

ftp>

or

DELE command okay.

ftp>

Changing the Name of a Remote Directory

Note

You can only rename a remote directory. You can not change the path to (move) a remote directory.

At the ftp > prompt, enter:

rename old_remote_directory_path new_remote_directory_path

where *old_remote_directory_path* is a full or relative path to an existing remote directory, and *new_remote_directory_path* is the same as *old_remote_directory_path*, except for the directory name.

Note

If you do not specify *new_remote_directory_path*, *ftp* prompts you for it by displaying:

(to-name)

Example Entries:

rename /users/lab/rich/doc /users/lab/rich/comment

rename doc comment

Result:

Ftp changes the name of the existing directory to the new name you specify.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

File exists, ready for destination name.

RNTO command okay.

Transferring Files with Ftp

Ftp lets you set up the file transfer environment before you actually transfer any files. This includes:

- setting the local and remote working directories,
- setting the file transfer type,
- turning on (or off) options to monitor file transfer progress, and
- turning on (or off) prompting to confirm each file transfer.

Once you have set up the file transfer environment, you can transfer files between your local host and a remote host.

1. Set the Local and Remote Working Directories

Use *ftp*'s *lcd* and *cd* commands to set the local and remote working directories before you begin the file transfer. These commands are discussed earlier in this chapter.

2. Set the File Transfer Type

The file transfer type you choose depends on the other operating system involved in the file transfer, **not** on the type of file you are transferring.

When you transfer files to or from a host with the HP-UX or UNIX operating system, you can use either the *ascii* or *binary* file transfer type. However, the file transfer is faster if you set the type to *binary*.

When you transfer files to or from a host with a non-HP-UX or non-UNIX operating system, you must use the *ascii* file transfer type. This insures that the files are transferred in a format appropriate for the other (different) operating system. The only time you may want to use the faster *binary* file transfer type for transferring files to or from a non-HP-UX or non-UNIX host is to archive them (for storage only, not access). The files would not be in a meaningful format if accessed on the destination host after the transfer. You would need to transfer the files back to a host with the same type of operating system from which they originated to access them in a meaningful form.

When you first invoke ftp, the file transfer type is preset to ascii.

Displaying the Current File Transfer Type

At the ftp > prompt, enter:

type

Result:

If the current file transfer type is ascii, *ftp* displays:

Using ascii mode to transfer files.

ftp>

If the current file transfer type is binary, *ftp* displays:

Using binary mode to transfer files.

ftp>

Changing the File Transfer Type to Binary

At the ftp > prompt, enter:

binary or

type binary

8-34 Transferring Files with Ftp

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

```
Type set to I.
```

```
ftp>
```

Changing the File Transfer Type to Ascii

At the ftp > prompt, enter:

ascii

or

type ascii

Result:

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

Type set to A.

3. Choose Options to Monitor File Transfer Progress

Within *ftp*, there are two ways to monitor the progress of a file transfer.

- You can have ftp display a hash (#) sign for each 1024 bytes transferred.
- You can have *ftp* sound a bell after each file transfer completes.

Both of these options are off when you first invoke *ftp*.

Turning Hash Sign Displaying On or Off

At the ftp > prompt, enter:

hash

Result:

If hash sign displaying was off, *ftp* turns it on and displays:

Hash mark printing on (1024 bytes/hash mark).

ftp>

If hash sign displaying was on, *ftp* turns it off and displays:

Hash mark printing off.

ftp>

Turning Bell Sounding On or Off

At the ftp > prompt, enter:

bell

If bell sounding was off, *ftp* turns it on and displays:

Bell mode on.

ftp>

If bell sounding was on, *ftp* turns it off and displays:

Bell mode off.

ftp>

4. Turn on Interactive Mode for Selective File Transfers

When you perform an operation involving multiple files, *ftp* can list the files one by one, asking for each whether you want to perform the operation or not. This allows you to perform the operation on only the files you select. *Ftp* calls this feature **interactive mode**.

Use interactive mode when you want to select from among multiple files to transfer during a single file transfer operation.

You can also use interactive mode when deleting multiple remote files and listing the contents of multiple remote directories to a local destination. However, in the latter case, *ftp* prompts you to confirm the local destination, not each remote directory listing.

Interactive mode is on when you first invoke *ftp*.

Turning Interactive Mode On or Off

At the ftp > prompt, enter:

prompt

If interactive mode was on, *ftp* turns it off and displays:

Interactive mode off.

ftp>

If interactive mode was off, *ftp* turns it on and displays:

Interactive mode on.

5. Perform One or More File Transfers

Note

When you transfer files with *ftp*, you are really copying those files from one place to another. Transfers do not move or delete the original files.

With Ftp, You Can Transfer	То
a remote file	a local file with the same directory path and name
	a local file with a different direc- tory path and/or name
multiple remote files	multiple local files with the same directory paths and names
a local file	a remote file with the same direc- tory path and name
	the end of a remote file with the same directory path and name
	a remote file with a different direc- tory path and/or name
	the end of a remote file with a dif- ferent directory path and/or name
multiple local files	multiple remote files with the same directory paths and names

Ftp File Transfer Options

The following sections explain all of these *ftp* options for transferring files.

Caution

In an *ftp* file transfer, insure that the source file and destination file are **not** the **same file** on the same host. Otherwise, *ftp* destroys the contents of the file.

From a Remote File To a Local File with the Same Directory Path and Name

This special case of an *ftp* file transfer requires a special syntax. The syntax you use causes *ftp* to behave in the following way:

- *Ftp* copies a file from the remote working directory to a file with the **same name** in the local working directory. The remote and local working directory paths may be different.
- *Ftp* copies a file from any other remote directory to a local file with the same directory path and name you identify the remote file with.

At the ftp > prompt, enter:

get remote_file_path

or

recv remote_file_path

where remote file path is either:

- the name of the source file if the file is in the remote working directory or
- the full or relative path to the source file if the file is in another remote directory.

Note

Any remote **directory path** you specify as part of a *remote_file_path* **must also exist on the local host**. Otherwise, *ftp* will not transfer the file.

Example Entries:

get program1

recv /usr/bin/game

Result:

If the file is in the remote working directory, *ftp* copies the file to a file with the same name in the local working directory.

Otherwise, *ftp* copies the remote file to a local file with the same directory path and name you identified the remote file with.

If the destination file already exists, *ftp* replaces its contents with the source file's contents.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file_name...

Transfer complete.

number bytes received in number seconds...

From a Remote File To a Local File with a Different Directory Path and/or Name

At the ftp > prompt, enter:

get remote_file_path local_file_path

or

recv remote_file_path local_file_path

where *remote_file_path* is a full or relative path to the source file in a remote directory and *local_file_path* is a full or relative path to the destination file in a local directory.

Note

Local_file_path must include the destination file's name.

Example Entries:

get program1 firstprog

recv program1 code/source/program1

get /users/lab/richard/graphics/designl designl

recv graphics/design1 graphics/work/firstdsgn

8-42 Transferring Files with Ftp

Ftp copies the file from the remote directory to the file in the local directory. If the destination file already exists, *ftp* replaces its contents with the source file's contents.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, ftp displays:

PORT command okay.

Opening data connection for remote_file_name...

Transfer complete.

number bytes received in number seconds...

From Multiple Remote Files To Multiple Local Files with the Same Directory Paths and Names

Ftp behaves in the following way when you transfer multiple files from a remote host to a local host:

- *Ftp* copies files from the remote working directory to files with the same **names** in the local working directory. The remote and local working directories may be different.
- *Ftp* copies files from any other remote directory to local files with the **same directory paths** and **names** you identify the remote files with.

Note

Metacharacters are expanded for *ftp*'s *mget* command if globbing is on.

If Interactive Mode Is On

1. At the ftp > prompt, enter:

mget remote_file_path ...

where *remote_file_path* is either:

- the name of a file if the file is in the remote working directory or
- the full or relative path to a file if the file is in another remote directory.

The ellipsis (...) means that you can specify multiple remote file paths.

Note

Any remote directory path you specify as part of a *remote_file_path* must also exist on the local host. Otherwise, *ftp* will not transfer the file(s).

Example Entries:

```
mget /bin/c*
```

mget program1 program2 program3

Result:

Ftp asks if you want to transfer the first remote file that matches what you specified. This gives you the option of **not** transferring the remote file.

Ftp displays:

mget first_remote_file_path?

2. To not transfer the remote file, enter:

Ν

To transfer the remote file, enter:

Y

Result:

If you enter N, *ftp* does **not** transfer the remote file and asks if you want to transfer the next remote file that matches what you specified, displaying:

mget next_remote_file_path?

If you enter Y, and:

- the file is in the remote working directory, *ftp* copies the file to a file with the same name in the local working directory.
- the file is in another remote directory, *ftp* copies the remote file to a local file with the same directory path and name.

Then *ftp* asks if you want to transfer the next remote file that matches what you specified.

• If verbose mode is off, *ftp* displays:

mget next_remote_file_path?

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file_path...

Transfer complete.

number bytes received in number seconds...

mget next_remote_file_path?

3. Repeat the previous step until ftp redisplays its ftp > prompt instead of the mget ...? prompt.

This means that there are no more files that match what you originally entered.

If Interactive Mode Is Off

At the ftp > prompt, enter:

mget remote_file_path ...

where remote_file_path is either:

- the name of a source file if the file is in the remote working directory or
- the full or relative path to a source file if the file is in another remote directory.

The ellipsis (...) means that you can specify multiple remote_file_paths.

Note

Any remote **directory path** you specify as part of a *remote_file_path* **must also exist on the local host**. Otherwise, *ftp* will not transfer the file(s).

Example Entries:

mget /bin/c*

mget program1 program2 program3

Result:

Ftp copies any files in the remote working directory to files with the same names in the local working directory.

Ftp copies any other remote files to local files with the same directory paths and names you identified the local files with.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file_path...

Transfer complete.

number bytes received in number seconds...

PORT command okay.

Opening data connection for remote_file_path...

Transfer complete.

number bytes received in number seconds...

. (repeated for each file transferred)

From a Local File To a Remote File with the Same Directory Path and Name

This special case of an *ftp* file transfer requires a special syntax. The syntax you use causes *ftp* to behave in the following way:

- *Ftp* copies a file from the local working directory to a file with the same **name** in the remote working directory. The local and remote working directory paths may be different.
- *Ftp* copies a file from any other local directory to a remote file with the **same directory path** and **name** you identify the local file with.

At the ftp > prompt, enter:

put local_file_path

or

send local_file_path

where *local_file_path* is either:

- the name of the source file if the file is in the local working directory or
- the full or relative path to the source file if the file is in another local directory.

Note

Any local directory path you specify as part of a *local_file_path* must also exist on the remote host. Otherwise, *ftp* will not transfer the file.

Example Entries:

put program2

send /tmp/printout

Result:

If the file is in the local working directory, *ftp* copies the file to a file with the same name in the remote working directory.

Otherwise, *ftp* copies the local file to a remote file with the same directory path and name you identified the local file with.

If the destination file already exists, *ftp* replaces its contents with the source file's contents.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file_name...

Transfer complete.

number bytes received in number seconds...

From a Local File To the End of a Remote File with the Same Directory Path and Name

This special case of an *ftp* file transfer requires a special syntax. The syntax you use causes *ftp* to behave in the following way:

- *Ftp* copies a file from the local working directory to the end of a file with the same name in the remote working directory. The local and remote working directory paths may be different.
- *Ftp* copies a file from any other local directory to the end of a remote file with the same name and directory path you identify the local file with.

At the ftp > prompt, enter:

append local_file_path

where *local_file_path* is either:

- the name of the source file if the file is in the local working directory or
- the full or relative path to the source file if the file is in another local directory.

Note

Any local **directory path** you specify as part of a *local_file_path* **must also exist on the remote host.** Otherwise, *ftp* will not transfer the file.

Example Entries:

append form

append /tmp/testdata

Result:

If the file is in the local working directory, *ftp* copies the file to the end of a file with the same name in the remote working directory.

Otherwise, *ftp* copies the local file to the end of a remote file with the same directory path and name you identified the local file with.

If the destination file does not exist, ftp creates it.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file_name...

Transfer complete.

number bytes received in number seconds...

From a Local File To a Remote File with a Different Directory Path and/or Name

At the ftp > prompt, enter:

put local_file_path remote_file_path

or

send local_file_path remote_file_path

where *local_file_path* is a full or relative path to the source file in a local directory and *remote_file_path* is a full or relative path to the destination file in a remote directory.

Note

Remote_file_path **must** include the destination file's name.

Example Entries:

put report results

send doc/internal/issues issues

put status /users/lab/richard/mail/update

send email/urgent/schedule project/schedule

Result:

Ftp copies the file from the local directory to the file in the remote directory. If the destination file already exists, *ftp* replaces its contents with the source file's contents.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file_name...

Transfer complete.

number bytes received in number seconds...

From a Local File To the End of a Remote File with a Different Directory Path and/or Name

At the ftp > prompt, enter:

append local_file_path remote_file_path

where *local_file_path* is a full or relative path to the source file in a local directory and *remote_file_path* is a full or relative path to the destination file in a remote directory.

Note

Remote_file_path **must** include the destination file's name.

Example Entries:

append /users/richard/doc/section2 comment/chapter1

append aliases .login

append email/bugreport /users/lab/richard/project/defects

Result:

Ftp appends the file from the local directory to the end of the file in the remote directory. If the destination file does not exist, *ftp* creates it.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file_name...

Transfer complete.

number bytes received in number seconds...

From Multiple Local Files To Multiple Remote Files with the Same Directory Paths and Names

Ftp behaves in the following way when you transfer multiple files from a local host to a remote host:

- *Ftp* copies files from the local working directory to files with the same **names** in the remote working directory. The local and remote working directory paths may be different.
- *Ftp* copies files from any other local directory to remote files with the same directory paths and names you identify the local files with.

Note

Metacharacters are expanded for *ftp*'s *mput* command if globbing is on.

If Interactive Mode Is On

1. At the ftp > prompt, enter:

mput local_file_path ...

where *local_file_path* is either:

- the name of a file if the file is in the local working directory or
- the full or relative path to a file if the file is in another local directory.

The ellipsis (...) means that you can specify multiple local_file_paths.

Note

Any local **directory path** you specify as part of a *local_file_path* **must also exist on the remote host**. Otherwise, *ftp* will not transfer the file(s).

Example Entries:

mput /etc/*
mput memo1 memo2 memo3

Result:

Ftp asks if you want to transfer the first local file that matches what you specified. This gives you the option of **not** transferring the local file.

Ftp displays:

mput first_local_file_path?

2. To not transfer the local file, enter:

Ν

To transfer the local file, enter:

Y

Result:

If you enter N, *ftp* does **not** transfer the local file and asks if you want to transfer the next local file that matches what you specified, displaying:

mput next_local_file_path?

If you enter Y, and:

- the file is in the local working directory, *ftp* copies the file to a file with the same name in the remote working directory.
- the file is in another local directory, *ftp* copies the local file to a remote file with the same directory path and name.

Then *ftp* asks if you want to transfer the next local file that matches what you specified.

• If verbose mode is off, *ftp* displays:

mput next_local_file_path?

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for local_file_path...

Transfer complete.

number bytes received in number seconds...

mput next_local_file_path?

3. Repeat the previous step until ftp redisplays its ftp > prompt instead of the mput ...? prompt.

This means that there are no more files that match what you originally entered.

If Interactive Mode Is Off

At the ftp > prompt, enter:

mput local_file_path ...

where *local_file_path* is either:

- the name of a source file if the file is in the local working directory or
- the full or relative path to a source file if the file is in another local directory.

The ellipsis (...) means that you can specify multiple local_file_paths.

Note

Any local **directory path** you specify as part of a *local_file_path* **must also exist on the remote host**. Otherwise, *ftp* will not transfer the file(s).

Example Entries:

mput /etc/*

mput memol memo2 memo3

Result:

Ftp copies any files in the local working directory to files with the same names in the remote working directory.

Ftp copies any other local files to remote files with the same directory paths and names you identified the local files with.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, ftp displays:

PORT command okay.

Opening data connection for remote_file_path...

Transfer complete.

number bytes received in number seconds...

PORT command okay.

Opening data connection for remote_file_path...

Transfer complete.

number bytes received in number seconds...

. (repeated for each file transferred)

Performing Other File Operations with Ftp

From within ftp, you can:

- display the contents of a remote file,
- create a remote file,
- append text to the end of a remote file,
- delete one or more remote files, and
- change the name of a remote file.

The following sections tell how to perform these operations.

Displaying the Contents of a Remote File

To display the contents of a remote file, you transfer the remote file to *stdout* (usually your display, or HP-UX terminal).

At the ftp > prompt, enter:

get remote_file_path -

or

recv remote_file_path -

where *remote_file_path* is a full or relative path to a remote file, and – represents *stdout* (usually the display).

Example Entries:

get /users/lab/richard/comment/readme -

recv comment/readme -

8-62 Performing Other File Operations with Ftp

Result:

Ftp sends the contents of the remote file you specify to the display (actually, *stdout*).

• If verbose mode is off, *ftp* displays:

```
contents_of_file
```

ftp>

• If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file...

contents_of_file

Transfer complete.

number bytes received in number seconds...

Creating a Remote File

To create a remote file, you transfer from *stdin* (usually your keyboard input, or HP-UX terminal) to the remote file.

1. At the ftp > prompt, enter:

put - remote_file_path

or

send - remote_file_path

where – represents *stdin* (usually keyboard input), and *remote_file_path* is a full or relative path to a remote file.

Example Entries:

put - /users/lab/richard/comment/note

send - comment/note

Result:

Ftp creates the file you specify and waits for you to enter what you want to put into the file.

- If verbose mode is off, *ftp* displays nothing.
- If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file...

8-64 Performing Other File Operations with Ftp

2. At the keyboard, enter what you want to put into the file.

Example Entry:

This is a test.

I am entering words into the file I am creating.

This is the last line.

Result:

Ftp displays what you enter as you enter it.

3. When you finish entering the contents of the file, press Return.

Result:

The cursor moves to a new line.

4. Press CTRL-D.

Result:

This signals the end of the file, and *ftp* adds what you entered to the file.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

Transfer complete.

number bytes received in number seconds...

ftp>

Appending Text to the End of a Remote File

To append text to the end of a remote file, you append from *stdin* (usually your keyboard input, or HP-UX terminal) to the end of the remote file.

Transferring Files with Ftp 8-65

1. At the ftp> prompt, enter:

append - remote_file_path

where – represents *stdin* (usually keyboard input) and *remote_file_path* is a full or relative path to a remote file.

Example Entries:

```
append - /users/lab/richard/comment/note
append - comment/note
```

Result:

Ftp waits for you to enter what you want to append to the end of the file.

- If verbose mode is off, *ftp* displays nothing.
- If verbose mode is on, *ftp* displays:

PORT command okay.

Opening data connection for remote_file...

2. At the keyboard, enter what you want to append to the end of the file.

Example Entry:

These are words that I want appended to the end of a file.

This is the last line.

Result:

Ftp displays what you enter as you enter it.

3. When you finish entering what you want to append to the file, press Return.

Result:

The cursor moves to a new line.

8-66 Performing Other File Operations with Ftp

4. Press CTRL-D.

Result:

This signals the end of the file, and *ftp* appends what you entered to the end of the file.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

Transfer complete.

number bytes received in number seconds...

Deleting a Remote File

At the ftp > prompt, enter:

delete remote_file_path

where remote_file_path is a full or relative path to a remote file.

Result:

Ftp deletes the remote file you specify.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

DELE command okay.

Deleting Multiple Remote Files

Note

Metacharacters are expanded for *ftp*'s *mdelete* command if globbing is on.

If Interactive Mode Is On

1. At the ftp > prompt, enter:

mdelete remote_file_path ...

where *remote_file_path* is a full or relative path to a remote file. The ellipsis (...) means that you can specify multiple *remote_file_paths*.

Example Entries:

mdelete /users/lab/richard/doc/spec?

mdelete doc/spec1 doc/spec2 doc/spec3

Result:

Ftp asks if you want to delete the first remote file that matches what you specified. This gives you the option of keeping the remote file.

Ftp displays:

mdelete first remote file path?

2. To keep the remote file, enter:

Ν

To delete the remote file, enter:

Y

Result:

If you enter N, *ftp* keeps the remote file and asks if you want to delete the next remote file that matches what you specified, displaying:

mdelete next_remote_file_path?

If you enter Y, *ftp* deletes the remote file and asks if you want to delete the next remote file that matches what you specified.

• If verbose mode is off, *ftp* displays:

mdelete next_remote_file_path?

• If verbose mode is on, *ftp* displays:

DELE command okay.

mdelete next_remote_file_path?

3. Repeat the previous step until ftp redisplays its ftp > prompt instead of the mdelete ...? prompt.

This means that there are no more files that match what you originally entered.

If Interactive Mode Is Off

At the ftp > prompt, enter:

mdelete remote_file_path ...

where *remote_file_path* is a full or relative path to a remote file. The ellipsis (...) means that you can specify multiple *remote_file_paths*.

Example Entries:

mdelete /users/lab/richard/doc/spec?

mdelete doc/spec1 doc/spec2 doc/spec3

Result:

Ftp deletes all of the files that match what you specified.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

DELE command okay.

DELE command okay.

. (repeated for each file deleted)

Changing the Name of a Remote File

Note

You can use *ftp*'s *rename* command to change the path to (move) a remote file. You can **not** use *ftp*'s *rename* command to change the path to (move) a remote directory.

At the ftp > prompt, enter:

rename old_remote_file_path new_remote_file_path

where *old_remote_file_path* is a full or relative path to an existing remote file, and *new_remote_file_path* is a full or relative path to a new file.

Note

If you do not specify *new_remote_file_path*, *ftp* prompts you for it by displaying:

(to-name)

Example Entries:

rename /users/lab/richard/doc/note /users/lab/richard/memo

rename memo memo.tmp

Result:

Ftp changes the name and/or the path to the remote file.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

File exists, ready for destination name.

RNTO command okay.

Obtaining Ftp Status

You can display the status of all of *ftp*'s feature settings. These include:

- whether or not a connection to a remote host exists and the name of the host to which a connection exists (if any),
- the file transfer type,
- whether or not *ftp* is set to sound a bell after each file transfer completes,
- whether or not interactive mode is on for multiple-file operations,
- whether or not globbing is on,
- whether or not *ftp* is set to display a hash mark for each 1024 bytes transferred during a file transfer, and
- other status information for settings that you cannot change or that need not be used in HP's implementation of *ftp*.

At the ftp > prompt, enter:

status

Result:

Ftp displays the status of all of its feature settings.

Setting Up Automatic Remote Login for Ftp

You can create a file named *.netrc* that lets *ftp* log you into a remote host automatically. You place remote login names and passwords in this file so that *ftp* need not prompt you for these. This feature can be useful for programs that need to perform *ftp* operations unattended.

Caution

Having a *.netrc* file is a serious security risk. Passwords in this file are unencrypted. Be sure to follow the directions below for "Protecting Your .netrc File."

Your *.netrc* file must be in your home directory on your local host. You can find out what your home directory is by entering:

echo \$HOME

Each entry you place in your .netrc file must have the following format:

machine remote_host_name login remote_login_name password remote_password

Follow these rules when creating a .netrc file:

- Each entry must contain a valid remote host name, remote login name and remote password. Valid remote host names are listed in the */etc/hosts* file on your local host.
- Separate each word in an entry with tabs, commas or blanks.
- Follow each remote host name with only one remote login name and one remote password.

Example \$HOME/.netrc File Entry:

If you wanted to set up automatic login to the remote host *hpabsa*, and your remote login name and password on that host were *carolyn* and *driveway*, respectively, you would create a *\$HOME/.netrc* file with the following entry:

machine hpabsa login carolyn password driveway

Protecting Your .netrc File

It is important to protect your *.netrc* file and your home directory to prevent unauthorized users from gaining access to the remote hosts and accounts in your *.netrc* file. To do so:

- 1. Insure that your .netrc file is owned by you, the user.
- 2. Use the HP-UX *chmod* command to protect your *.netrc* file with 0400 (-r-----) permission.
- **3.** Use the HP-UX *chmod* command to protect your home directory so that no one else can read it or write to it. For example, you should protect your home directory with at least 0711 (-rwx--x--x) permission.
- 4. Insure that your local account has a password.

To automatically log into a remote host once your *.netrc* file is set up, just invoke *ftp* and open a connection to the remote host. *Ftp* then automatically logs you into that host.

Logging into a Remote Host with a Login Not in Your .netrc File

If you need to log into a remote host as someone else, you can override the automatic login (*.netrc* file) you set up for *ftp*. To do so, you:

- invoke *ftp* with the -n option to disable automatic login,
- choose whether you want *ftp* to display responses from the remote host (choose whether you want verbose mode on or off),
- connect to the remote host, and then
- log into the remote host with *ftp*'s user command.

1. Invoke Ftp with the -n Option

At your HP-UX prompt, enter:

ftp -n

Result:

Ftp displays its prompt:

ftp>

Note

Automatic login remains disabled for your entire *ftp* session.

2. Choose Whether You Want Verbose Mode On or Off

If you want *ftp* to display responses from the remote host, insure that verbose mode is on. Otherwise turn it off.

Use *ftp*'s *status* command to check the verbose mode setting, and use *ftp*'s *verbose* command to change the setting if you want to. These two *ftp* commands are explained in earlier sections of this chapter.

3. Connect to the Remote Host

Connect to the remote host as you normally would with *ftp*'s *open* command. This command is explained earlier in this chapter.

4. Log into the Remote Host with Ftp's User Command

There are two ways to log into a remote host with *ftp*'s *user* command. The faster way allows you to enter all of the login information on one line, but displays the remote password as you enter it. The other way causes *ftp* to prompt you for the remote password and does not display the remote password as you enter it.

Logging into a Remote Host with a Single Command Line

Caution

This method displays the remote password as you enter it.

At the ftp > prompt, enter:

user remote_login_name remote_password [account]

where *remote_login_name* and *remote_password* must be valid on the remote host. Some remote hosts require you to enter a valid account name.

Example Entry:

user richard soccer

Result:

Ftp checks the remote login name, password, and account (if applicable) for validity and logs you into the remote host if they are valid.

• If verbose mode is off, ftp displays:

ftp>

• If verbose mode is on, *ftp* displays:

Password required for remote_login_name.

User remote_login_name logged in.

ftp>

Logging into a Remote Host without Displaying the Remote Password

1. At the ftp > prompt, enter:

user remote_login_name

where remote login name must be valid on the remote host.

Example Entry:

user richard

Result:

Ftp prompts you for the remote password associated with the remote login name you entered.

• If verbose mode is off, *ftp* displays:

Password:

• If verbose mode is on, ftp displays:

Password required for remote_login_name.

Password:

2. Enter the remote password associated with the remote login name you gave. (The password is not displayed as you enter it.)

Result:

Ftp checks the remote password for validity and logs you into the remote host if your password is valid. Some remote hosts may require you to enter a valid account name before you are logged in.

• If verbose mode is off, *ftp* displays:

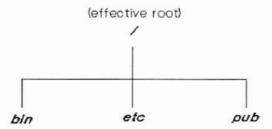
ftp>

• If verbose mode is on, *ftp* displays:

User remote_login_name logged in.

The Public Ftp Account

Some remote hosts may have a public (guest or anonymous) *ftp* account. When you log into this account, it becomes your **effective** root directory, and you cannot access anything above it. A public *ftp* account has the following directory structure:



User's View of Public Ftp Directory Structure

Directory	Description
bin	This directory contains copies of the $ls(1)$, $sh(1)$ and $csh(1)$ programs to support <i>ftp</i> 's <i>dir</i> , <i>ls</i> , and <i>pwd</i> commands.
etc	This directory contains copies of the files $passwd(1)$ and $group(1)$. These files must be present for ftp 's dir , ls , and pwd commands to work properly.
pub	Users can place files in this directory for public access. For example, this directory might contain announcements, re- quests for comment, or host tables.

Logging into the Public (Anonymous) Ftp Account

- **1.** Invoke *ftp*, insure that verbose mode is set to what you want, and connect to a remote host as you normally would.
- 2. When *ftp* prompts for a remote login name with the Name (...): prompt, enter:

ftp

Result:

Ftp prompts for the remote password associated with the *ftp* remote login name, displaying:

Password (remote_host:ftp):

3. Enter the name of your local host.

Example Entry:

hpabsa

Result:

If a public (guest or anonymous) *ftp* account exists on the remote host, *ftp* logs you into that account and makes the *ftp* directory your working and effective root directory.

• If verbose mode is off, *ftp* displays:

ftp>

• If verbose mode is on, *ftp* displays:

Guest login ok, send ident as password.

Guest login ok, access restrictions apply.

ftp>

8-82 The Public Ftp Account

Specifying Ftp Settings and Connecting to a Remote Host When You Invoke Ftp

You can change some of *ftp*'s settings when you first invoke *ftp*, if you know ahead of time how you want *ftp* set up. You can also connect to a remote host when you invoke *ftp*.

The *ftp* command you give from your local HP-UX prompt can take the following form:

ftp [-g] [-i] [-n] [-v] [remote_host]

Anything in brackets is optional. The following table explains the effect of each option above.

Option	Effect
-g	This option turns off globbing (metacharacter expansion).
-i	This option turns off interactive mode for multiple-file operations. This is useful if, for example, you know ahead of time that you want to use <i>ftp</i> to perform bulk, instead of selective, file deletions or file transfers.
-n	This option disables automatic login (as set up by a \$HOME/. <i>netrc</i> file). Use of this option is explained in an earlier section of this chapter.
-v	This option turns on verbose output (the display of responses from any remote host you connect to). If you invoke <i>ftp</i> from your keyboard (HP-UX terminal), <i>ftp</i> turns on verbose output without this option. This option is only useful if <i>ftp</i> is invoked indirectly. For example, if a program invokes <i>ftp</i> with this option, and <i>ftp</i> 's output is going to a file, the output file will contain a "log" of <i>ftp</i> 's results.

Specifying a remote host's name or alias (as listed in your local host's *letc/hosts* file) on the *ftp* command line causes *ftp* to connect to that remote host without your having to use *ftp*'s *open* command.

8-84 Specifying Ftp Settings and Connecting to a Remote Host When You Invoke Ftp

Transferring Files with Rcp

Rcp is a Berkeley Service that allows you to copy files between only HP-UX or UNIX hosts on the network.

Rcp can copy the contents of an entire directory. This includes all files and the contents of all subdirectories within that directory. From your local host, you can also copy files between two remote hosts.

Rcp allows file transfers to and from other hosts only if the configuration files that this service uses are set up properly. These files are mentioned later in this chapter.

File Copy Concepts

With rcp you can copy files and directories

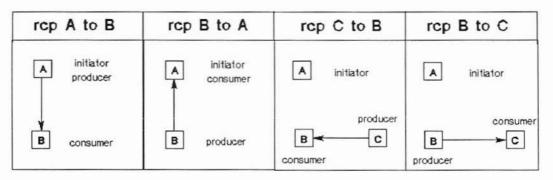
From	То
your local host	a remote host
a remote host	your local host
a remote host	another remote host

These scenarios can be represented by the initiator/producer/consumer model.

In any *rcp* operation, there may be up to three hosts involved: an initiator, a producer and a consumer.

- The **initiator** is your local host, the host on which you make the *rcp* request. It receives requests from you and starts file copies.
- The **producer** is the host that the source file or directory is on. It accesses the source file or directory and produces the data to be copied.
- The **consumer** is the host that the destination file or directory resides on or will reside on. It receives the incoming data and writes it to the destination file or directory.

In the diagram below, the arrows represent data being copied from a source file to a destination file. If host A is the initiator in each case, the location of the producer and consumer depends on the location of the source and destination files.



Initiator/Producer/Consumer Model

Using Rcp

Rcp allows you to copy files to or from a remote host if the remote host is configured in either of two ways.

Either:

- you must have an account on the remote host with the same login name as your local login name, and
- the name of your local host must be in the remote host's */etc/hosts.equiv* file,

or:

- you must have an account on the remote host, and
- the name of your local host and your local login name must be in a *.rhosts* file in your home directory on the remote host.

The next section explains how to create a remote *\$HOME/.rhosts* file for yourself, if you need to do so. Otherwise, skip the next section.

Creating a \$HOME/.rhosts File on a Remote Host

If you have an account on a remote host, you can give yourself *rcp* access to your remote account by creating a file named *.rhosts* in your remote home directory. You can find out what your remote home directory is by entering:

echo \$HOME

on the remote host. You must place the name of your local host and your local login name in the *.rhosts* file you create.

Note

A *\$HOME/.rhosts* file creates a significant security risk. Be sure to follow the directions below for "Protecting Your .rhosts File."

The entry you place in the remote *.rhosts* file must have the following format:

your_local_host's_name your_local_login_name

You can separate your_local_host's_name and your_local_login_name with any number of tabs or spaces. Put any comments after your_local_login_name.

Example \$HOME/.rhosts File Entry

If your local host's name were *hpabsa* and your local login name were *richard*, on the remote host you would create a *\$HOME/.rhosts* file with the following entry:

hpabsa richard

Protecting Your \$HOME/.rhosts File

It is important to protect your remote *.rhosts* file and home directory to prevent unauthorized users from gaining *rcp* access to your remote account and host. Only you should be able to create a *.rhosts* file in your remote home directory and write entries to the file. To do so:

1. Insure that your remote .rhosts file is owned by you, the user.

- 2. Use the HP-UX or UNIX *chmod* command to protect your remote *.rhosts* file with 0400 (-r-----) permission.
- 3. Use the HP-UX or UNIX *chmod* command to protect your remote home directory so that no one else can read it or write to it. For example, you should protect your remote home directory with at least 0711 (-rwx--x--x) permission.

Performing Copy Operations with Rcp

Note

When you copy remote files and directories, the working directory for *rcp* on the remote host is your remote *\$HOME* directory.

With *rcp*, you can copy from:

- a single local or remote file,
- multiple local and/or remote files,
- a single local or remote directory,
- multiple local and/or remote directories,
- any combination of local and/or remote files and directories.

What *rcp* can copy to depends on what *rcp* is copying from.

Note

Rcp can copy **from** only ordinary files and directories, **not** special files and directories (such as */dev* files). However, *rcp* can copy **to** special files (such as */dev* files).

Note

In *rcp* file transfers you **must** explicitly specify the destination file or directory.

Any output generated by commands in a *.login*, *.profile*, or *.cshrc* file on the remote host can cause *rcp* errors.

Caution

Do not specify the same source and destination files. This can corrupt the file's contents.

When *rcp* completes a copy operation, your local host redisplays its prompt.

The following sections explain all of *rcp*'s copy options.

Source	Allowed Destinations	Result of Copy
A Local File	A New Remote File	<i>Rcp</i> creates the destination file and copies the source file's contents into the destination file.
	An Existing Remote File	<i>Rcp</i> overwrites the destination file's contents with the source file's contents.
	An Existing Remote Directory	<i>Rcp</i> copies the source file into the des- tination directory.
Local Files	An Existing Remote Directory	<i>Rcp</i> copies the source files into the des- tination directory.

From a Local Producer to a Remote Consumer

If the destination is a link to a file, the file to which the destination is linked is overwritten, and all links to the file remain the same.

At your HP-UX prompt, enter:

rcp local_path ... remote_host:remote_path

where:

- *local_path* is the path relative to your local working directory or the full path from the local root directory,
- the ellipsis (...) means that you can specify multiple local_paths,

• remote host is the name or alias of a host listed in /etc/hosts, and

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

• *remote_path* is the path relative to your remote home directory or the full path from the remote root directory.

Example Entry:

rcp /users/alan/program mail/defects hpabsb:project

Source	Allowed Destinations	Result of Copy
A Local Directory	A New Remote Direc- tory	<i>Rcp</i> creates the destination directory and copies the contents of the source directory into the destination directory.
	An Existing Remote Directory	<i>Rcp</i> copies the source directory itself along with its contents into the destina- tion directory. <i>Rcp</i> overwrites any exist- ing files.
Local Directories	An Existing Remote Directory	<i>Rcp</i> copies the source directories themselves along with their contents into the destination directory.
Any Combination of Local Files and Directories	An Existing Remote Directory	<i>Rcp</i> copies the source files and the source directories themselves along with their contents into the destination directory.

At your HP-UX prompt, enter:

rcp -r local_path ... remote_host:remote_path

where:

- -r (recursive option) causes *rcp* to copy the contents of any source directories,
- *local_path* is the path relative to your local working directory or the full path from the local root directory,
- the ellipsis (...) means that you can specify multiple local_paths,

• remote host is the name or alias of a host listed in /etc/hosts, and

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

• *remote_path* is the path relative to your remote home directory or the full path from the remote root directory.

Example Entry:

rcp -r mail /users/alan/memos hpabsb:correspondence

Source	Allowed Destinations	Result of Copy
A Remote File	A New Local File	<i>Rcp</i> creates the destination file and copies the source file's contents into the destination file.
	An Existing Local File	<i>Rcp</i> overwrites the destination file's contents with the source file's contents.
	An Existing Local Directory	<i>Rcp</i> copies the source file into the des- tination directory.
Remote Files	An Existing Local Directory	<i>Rcp</i> copies the source files into the des- tination directory.

From One or More Remote Producers to a Local Consumer

If the destination is a link to a file, the file to which the destination is linked is overwritten, and all links to the file remain the same.

At your HP-UX prompt, enter:

rcp remote_host:remote_path ... local_path

where:

• remote_host is the name or alias of a host listed in /etc/hosts,

Note

The file *letc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

- *remote_path* is the path relative to your remote home directory or the full path from the remote root directory,
- the ellipsis (...) means that you can specify multiple remote paths, and
- *local_path* is the path relative to your local working directory or the full path from the local root directory.

Example Entry:

rcp hpabsb:/users/alan/graphics/logo hpabsb:form templates

Source	Allowed Destinations	Result of Copy
A Remote Directory	A New Local Direc- tory	<i>Rcp</i> creates the destination directory and copies the contents of the source directory into the destination directory.
	An Existing Local Directory	<i>Rcp</i> copies the source directory itself along with its contents into the destina- tion directory. <i>Rcp</i> overwrites any exist- ing files.
Remote Directories	An Existing Local Directory	<i>Rcp</i> copies the source directories themselves along with their contents into the destination directory.
Any Combination of Remote Files and Directories	An Existing Local Directory	<i>Rcp</i> copies the source files and the source directories themselves along with their contents into the destination directory.

At your HP-UX prompt, enter:

rcp -r remote_host:remote_path ... local_path

where:

-r (recursive option) causes *rcp* to copy the contents of any source directories,

• remote_host is the name or alias of a host listed in /etc/hosts,

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

- *remote_path* is the path relative to your remote home directory or the full path from the remote root directory,
- the ellipsis (...) means that you can specify multiple local_paths, and
- *local_path* is the path relative to your local working directory or the full path from the local root directory.

Example Entry:

rcp -r hpabsb:/users/alan/document hpabsb:paper textfiles

From One or More Remote Producers to a Remote Consumer

Rcp allows you to copy files between two remote hosts if the remote consumer host is configured in either of two ways:

Either:

- you must have an account on the remote consumer host with the same login name you have on the remote producer host, and
- the name of the remote producer host must be in the remote consumer host's /etc/hosts.equiv file,

or:

- you must have an account on the remote consumer host, and
- the name of the remote producer host and your login name on the remote producer host must be in a *.rhosts* file in your home directory on the remote consumer host.

Source	Allowed Destinations	Result of Copy
A Remote File	A New Remote File	<i>Rcp</i> creates the destination file and copies the source file's contents into the destination file.
	An Existing Remote File	<i>Rcp</i> overwrites the destination file's contents with the source file's contents.
	An Existing Remote Directory	<i>Rcp</i> copies the source file into the des- tination directory.
Remote Files	An Existing Remote Directory	<i>Rcp</i> copies the source files into the des- tination directory.

At your HP-UX prompt, enter:

rcp remote_host:remote_path ... remote_host:remote_path

where:

• remote_host is the name or alias of a host listed in /etc/hosts,

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

- *remote_path* is the path relative to your remote home directory or the full path from the remote root directory, and
- the ellipsis (...) means that you can specify multiple remote paths.

Example Entry:

rcp hpabsb:graphics/logo hpabsb:form hpabsc:templates

Source	Allowed Destinations	Result of Copy
A Remote Directory	A New Remote Directory	<i>Rcp</i> creates the destination directory and copies the contents of the source directory into the destination directory.
	An Existing Remote Directory	<i>Rcp</i> copies the source directory itself along with its contents into the destina- tion directory. <i>Rcp</i> overwrites any exist- ing files.
Remote Directories	An Existing Remote Directory	<i>Rcp</i> copies the source directories themselves along with their contents into the destination directory.
Any Combination of Remote Files and Directories	An Existing Remote Directory	<i>Rcp</i> copies the source files and the source directories themselves along with their contents into the destination directory.

At your HP-UX prompt, enter:

rcp -r remote_host:remote_path ... remote_host:remote_path

where:

• -r (recursive option) causes *rcp* to copy the contents of any source directories,

• remote_host is the name or alias of a host listed in /etc/hosts,

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The *letc/hosts* file may contain comments and other information as well.

- *remote_path* is the path relative to your remote home directory or the full path from the remote root directory, and
- the ellipsis (...) means that you can specify multiple remote _paths.

Example Entry:

rcp -r hpabsb:document hpabsb:paper hpabsc:textfiles

Source	Allowed Destinations	Result of Copy	
Local and Remote	An Existing Local	Rcp copies the source files into the	
Files	Directory	destination directory.	

From Local and Remote Producers to a Local Consumer

At your HP-UX prompt, enter:

rcp local_path ... remote_host:remote_path ... local_path

where:

- *local_path* is the path relative to your local working directory or the full path from the local root directory,
- the ellipses (...) mean that you can specify multiple *local_paths* or *remote_paths*,
- remote_host is the name or alias of a host listed in /etc/hosts, and

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

• *remote_path* is the path relative to your remote home directory or the full path from the remote root directory.

Example Entry:

rcp module1 hpabsb:module2 /code/integration

Source	Allowed Destinations	Result of Copy
Local and Remote Directories	An Existing Local Directory	<i>Rcp</i> copies the source directories themselves along with their contents into the destination directory.
Any Combination of Local and Remote Files and Directories	An Existing Local Directory	<i>Rcp</i> copies the source files and the source directories themselves along with their contents into the destination directory.

At your HP-UX prompt, enter:

rcp -r local_path... remote_host:remote_path... local_path

where:

- -r (recursive option) causes *rcp* to copy the contents of any source directories,
- *local_path* is the path relative to your local working directory or the full path from the local root directory,
- the ellipsis (...) means that you can specify multiple local_paths,
- remote_host is the name or alias of a host listed in /etc/hosts, and

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

• *remote_path* is the path relative to your remote home directory or the full path from the remote root directory.

Example Entry:

rcp -r /users/alan/drawings hpabsc:charts /lib/graphics

From Local and Remote Producers to a Remote Consumer

Source	Allowed Destinations	Result of Copy	
Local and Remote	An Existing Remote	<i>Rcp</i> copies the source files into the des-	C
Files	Directory	tination directory.	

At your HP-UX prompt, enter:

rcp local_path... remote_host:remote_path... remote_host:remote_path

where:

- *local_path* is the path relative to your local working directory or the full path from the local root directory,
- the ellipses (...) mean that you can specify multiple *local_paths* or *remote_paths*,
- remote_host is the name or alias of a host listed in /etc/hosts, and

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The *letc/hosts* file may contain comments and other information as well.

• *remote_path* is the path relative to your remote home directory or the full path from the remote root directory.

Example Entry:

rcp logfile1 hpabsb:logfile2 hpabsc:/tests/results

Source	Allowed Destinations	Result of Copy	
Local and Remote Directories	An Existing Remote Directory	<i>Rcp</i> copies the source directories themselves along with their contents into the destination directory.	
Any Combination of Local and Remote Files and Directories	An Existing Remote Directory	<i>Rcp</i> copies the source files and the source directories themselves along with their contents into the destination directory.	

At your HP-UX prompt, enter:

rcp -r local_path... remote_host:remote_path... remote_host:remote_path

where:

- -r (recursive option) causes *rcp* to copy the contents of any source directories,
- *local_path* is the path relative to your local working directory or the full path from the local root directory,
- the ellipses (...) mean that you can specify multiple *local_paths*,
- remote_host is the name or alias of a host listed in /etc/hosts, and

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet_address official_name alias ...

The ellipsis (...) means that a host may have multiple *aliases*. The */etc/hosts* file may contain comments and other information as well.

• *remote_path* is the path relative to your remote home directory or the full path from the remote root directory.

Example Entry:

rcp -r /users/alan/reports hpabsb:article hpabsc:newsitems

Rcp's Effect on File Attributes

Rcp sets the last access time of any source files and/or directories to the time that the copy occurs.

If the destination file or directory	exists,	does not exist,
then the mode is	unchanged	same as source file
owner is	unchanged	same as user's
group is	unchanged	same as user's
last access time is	unchanged	set to the time when copy occurred
last modification time is	set to the time when copy occurred	set to the time when copy occurred

Using "Wild Card" Characters, or Metacharacters with Rcp

In an *rcp* command, a local file or directory path can contain any metacharacters that are allowed by the shell you are using. Metacharacters, or wild card characters, stand for a set of characters or character strings and are a "shorthand" way of specifying a set of directory or file names. Your local shell expands the metacharacters into the directory and file names they match before *rcp* performs the copy operation.

In an *rcp* command, for any metacharacters in a remote file or directory path to be expanded on the remote host, not on the local host, you must enclose each remote source path in single (") or double (" ") quotes. You can also escape individual metacharacters by preceding them with a back-slash (\) so that the remote host expands them.

Example Entries:

rcp -r hpabsb:"*.c" /users/alan/cprograms

rcp -r hpabsb:*.c /users/alan/cprograms

Remember that if the source specification includes any directories (not files only), you must use the rcp -r syntax.

Copying Remote Files and Directories as Someone Else on the Remote Host

With rcp you can assume the identity of another user on a remote host if:

- you know that user's login name on the remote host and
- that user has your local host name and local login name in a *.rhosts* file in his or her home directory on the remote host.

When you copy remote files and directories under these conditions, the working directory for *rcp* on the remote host is the user's remote home directory.

To assume the identity of another user on a remote host, you use the following syntax for a remote file or directory:

remote_host.remote_login_name:remote_path

where:

- remote_host is the name or alias of a host listed in /etc/hosts,
- remote login name is the login name of the remote user, and
- *remote_path* is the path relative to the remote user's home directory or the full path from the remote root directory.

Example File Syntax:

hpabsb.alan:cprograms/modulel.c

Giving Other Remote Users Rcp Access to Your Local Account

You can give remote users *rcp* access to your local account by creating a *.rhosts* file. You place remote users' host names and login names in this file so that *rcp* lets them assume your identity when copying files to or from your local host.

Caution

A *\$HOME/.rhosts* file creates a significant security risk. Be sure to follow the instructions below for "Protecting Your .rhosts File."

Your .rhosts file must be in your home directory on your local host.

Each entry you place in your local *.rhosts* file must have the following format:

remote_host_name remote_login_name

Follow these rules when creating a .rhosts file:

- Each entry must contain a valid remote host name and remote login name.
- Separate the host name and login name with any number of tabs or blanks.
- Put any comments after the login name in any entry.

Example \$HOME/.rhosts File Entry

If you wanted to give user *cdm* on remote host *hpabsb rcp* access to your local account, you would create a local *\$HOME/.rhosts* file with the following entry:

hpabsb cdm

Protecting Your .rhosts File

It is important to protect your local *.rhosts* file and your local home directory to prevent unauthorized users from gaining *rcp* access to your account and local host. Only you should be able to create a *.rhosts* file in your home directory and write entries to it. To do this:

- 1. Insure that your .rhosts file is owned by you, the user.
- 2. Use the HP-UX *chmod* command to protect your local *.rhosts* file with 0400 (-r-----) permission.
- **3.** Use the HP-UX *chmod* command to protect your local home directory so that no one else can read it or write to it. For example, you should protect your local home directory with at least 0711 (-rwx--x--x) permission.

Executing Commands with Remsh

Remsh is a Berkeley Service that allows you to execute commands on a remote HP-UX or UNIX host on the network. *Remsh* is the same command as *rsh* in 4.2 BSD and later versions.

This chapter will cover:

- Setting Up Permission to Use Remsh on a Remote Host
- · Executing Commands on a Remote Host as Yourself
- Executing Commands on a Remote Host as Someone Else
- Giving Other Remote Users Remsh Access to Your Local Account
- Executing More Than One Remote Command with Remsh
- Using Shell Metacharacters with Remsh
- Using Remsh with Remote Commands That Do Not Take Input
- Using Remsh's "Shorthand" Syntax

Setting Up Permission to Use Remsh on a Remote Host

Caution

Do **not** use *remsh* to run an interactive command, such as *vi* or *more*. With some interactive commands, *remsh* hangs. To run interactive commands, log into the remote host with *rlogin*.

Remsh allows you to execute a command on a remote host if the remote host is configured in either of two ways.

Either:

- you must have an account on the remote host with the same login name as your local login name, and
- the name of your local host must be in the remote host's *letc/hosts.equiv* file,

or:

- you must have an account on the remote host, and
- the name of your local host and your local login name must be in a *.rhosts* file in your home directory on the remote host.

The next section explains how to create a remote *\$HOME/.rhosts* file for yourself, if you need to do so. Otherwise, skip the next section.

For more information about remote hosts, see the *hosts.equiv* (4) reference pages.

Creating a \$HOME/.rhosts File on a Remote Host

If you have an account on a remote host, you can give yourself *remsh* access to your remote account by creating a file named *.rhosts* in your remote home directory. You can find out what your remote home directory is by entering:

echo \$HOME

on the remote host. You must place the name of your local host and your local login name in the *.rhosts* file you create.

Caution

A *\$HOME*/*.rhosts* file creates a significant security risk. Be sure to follow the directions below for "Protecting Your .rhosts File."

The entry you place in your .rhosts file must have the following format:

your_local_host's_name your_local_login_name

You can separate your_local_host's_name and your_local_login_name with any number of tabs or spaces. Put any comments after your_local_login_name.

Example \$HOME/.rhosts File Entry

If your local host's name were *hpabsa* and your local login name were *richard*, on the remote host you would create a *\$HOME/.rhosts* file with the following entry:

hpabsa richard

Protecting Your \$HOME/.rhosts File

It is important to protect your remote *.rhosts* file and home directory to prevent unauthorized users from gaining *remsh* access to your remote account and host. Only you should be able to create a *.rhosts* file in your remote home directory and write entries to the file. To do so:

1. Insure that your remote .rhosts file is owned by you, the user.

- 2. Use the HP-UX or UNIX *chmod* command to protect your remote *.rhosts* file with 0400 (-r-----) permission.
- 3. Use the HP-UX or UNIX *chmod* command to protect your remote home directory so that no one else can read it or write to it. For example, you should protect your remote home directory with at least 0711 (-rwx--x--x) permission.

Executing Commands on a Remote Host as Yourself

Note

When you execute a command on a remote host, the working directory for *remsh* on the remote host is your remote *\$HOME* directory.

Remsh passes interrupt, terminate, quit and hangup signals to the remote command you execute.

At your HP-UX prompt, enter:

remsh remote_host command

where *remote_host* is the name or alias of a host listed in */etc/hosts* and *command* is a non-interactive HP-UX or UNIX command to execute on the remote host.

Note

The file */etc/hosts* contains entries for hosts with which you can communicate using ARPA/Berkeley Services. For each host, the file has a line containing the host's:

internet address official name alias ...

The ellipsis (...) means that a host may have multiple aliases. The */etc/hosts* file may contain comments and other information as well.

Example Entry:

remsh hpabsb cp form form.bkp

Result:

Remsh searches for the command you specify in the following remote directories in the order shown:

1. /bin

2. /usr/bin

3. /usr/contrib/bin

4. /usr/local/bin

On finding the command, *remsh* executes the command on the remote host and then your local host redisplays its prompt.

Note that if you do not give any command on the *remsh* command line, *remsh* interprets any options in the command line as *rlogin* options and runs *rlogin*.

Executing Commands on a Remote Host as Someone Else

With *remsh* you can execute a command as another user on a remote host if that user has your local host name and local login name in a *.rhosts* file in his or her home directory on the remote host.

Note

When you execute a command under these conditions, the working directory for *remsh* on the remote host is the remote user's home directory.

If the remote user's account has no password, you can use *remsh* to execute remote commands as that user without having your local host name and local login name in the user's *\$HOME/.rhosts* file.

At your HP-UX prompt, enter:

remsh remote host -1 remote login name command

where:

- remote_host is the name or alias of a host listed in /etc/hosts,
- *remote_login_name* is the login name of the remote user who you want to execute the command as, and
- command is a command to execute on the remote host.

Giving Other Remote Users Remsh Access to Your Local Account

You can give remote users *remsh* access to your local account by creating a *.rhosts* file in your local home directory. You place remote users' host names and login names in this file so that *remsh* lets them execute commands as you on your local host. (For more information, see the *hosts.equiv* (4) reference pages.)

Caution

A *\$HOME/.rhosts* file creates a significant security risk. Be sure to follow the instructions below on "Protecting Your .rhosts File."

Your .rhosts file must be in your home directory on your local host.

Each entry you place in your *.rhosts* file must have the following format:

remote_host_name remote_login_name

Follow these rules when creating a .rhosts file:

- Each entry must contain a valid remote host name and remote login name.
- Separate the host name and login name with any number of tabs or blanks.
- Put any comments after the login name in any entry.

Example \$HOME/.rhosts File Entry

If you wanted to give user *cdm* on remote host *hpabsc remsh* access to your local account, you would create a *\$HOME/.rhosts* file on your local host with the following entry:

hpabsc cdm

Protecting Your .rhosts File

It is important to protect your *.rhosts* file and your local home directory to prevent unauthorized users from gaining *remsh* access to your local account. Only you should be able to create a *.rhosts* file in your home directory and write entries to it. To do this:

- 1. Insure that your .rhosts file is owned by you, the user.
- 2. Use the HP-UX *chmod* command to protect your *.rhosts* file with 0400 (-r-----) permission.
- **3.** Use the HP-UX *chmod* command to protect your local home directory so that no one else can read it or write to it. For example, you should protect your local home directory with at least 0711 (-rwx--x--x) permission.
- 4. Insure that your account has a password. Otherwise, anyone can execute commands as you (with your login name) on your local host.

Executing More Than One Remote Command with Remsh

When you use *remsh* to execute more than one remote command, be aware of the following: a new remote shell executes the command(s) on each *remsh* command line. Therefore, when a remote command terminates, its process attributes (such as its environment and current working directory) disappear along with the shell that executed the command. For example, a remote *cd* command executed with *remsh* isolates the change of working directory to that instance of the command. A subsequent remote *pwd* command executed with *remsh* does not reflect the previous change in the working directory.

To execute more than one remote command without losing process attributes from one command to the next, you must:

- put the commands on a single remsh command line,
- separate the commands with semicolons (;), and
- enclose the string of commands in quotes (" ") so that the remote host executes every command. Otherwise, your local host executes any command(s) after the first one on the command line.

You can also have the remote host execute the string of commands by preceding each semicolon separator with a back slash (\backslash).

At your HP-UX prompt, enter:

remsh remote_host "command; ... command"

or

remsh remote_host command \;...command

where *remote_host* is the name or alias of a host listed in */etc/hosts*, and *command* is a non-interactive HP-UX or UNIX command to execute on the remote host. The ellipsis (...) means that you can specify more than one *command*; or *command* \;.

Example Entries:

remsh hpabsb "pwd; cd reports; pwd"
remsh hpabsb pwd \; cd reports \; pwd

Result:

Remsh executes the commands in sequence on the remote host, and your local host redisplays its prompt. Each remote command inherits the preceding one's process attributes. For example, the last *pwd* command in the example entries above would show *reports* as the remote working directory.

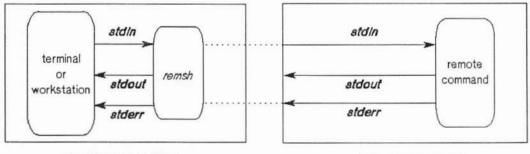
Using Shell Metacharacters with Remsh

Commands can contain metacharacters to be interpreted on either the local host or the remote host. The metacharacters you may use are those allowed by the shell you are using.

To Have Metacharacters Interpreted	Do This
On Your Local Host	Specify them as you normally would for local commands.
On a Remote Host	Enclose them in double quotes (" ").

Stdin, Stdout, and Stderr for Remsh

Remsh's stdin becomes the remote command's *stdin*, and the remote command's *stdout* and *stderr* become *remsh's stdout* and *stderr*, as illustrated below.



On Your Local Host

On a Remote Host

This means that you can use metacharacters to:

- redirect a remote command's input (stdin), output (stdout) and diagnostic output (stderr),
- pipe the output of a remote command into another remote or local command, and
- pipe the output of another remote or local command into a remote command.

For example, the command

remsh hpabsb cat remotefile > localfile

appends the remote file *remotefile* to the local file *localfile*. In contrast, the command

remsh hpabsb cat remotefile1 ">" remotefile2

appends the remote file remotefile1 to the other remote file remotefile2.

The command

```
remsh hpabsb cat /tmp/broadcastmsg | wall
```

displays the message in the remote file *broadcastmsg* on all of the local host's terminals. In contrast, the command

```
remsh hpabsb cat /tmp/broadcastmsg "|" wall
```

displays the message in the remote file *broadcastmsg* on all of the remote host's terminals.

Enclosing the metacharacters in double quotes causes the remote host to interpret them, instead of the local host.

Using Remsh with Remote Commands That Do Not Take Input

Remsh cannot determine if a remote command requires input. Therefore, *remsh* operates on the assumption that all remote commands require input. This behavior can cause problems if you use *remsh* to execute a remote command that does not require input. The *remsh* command attempts to read input (*stdin*) on the local host, even though the remote command requires none. The following examples illustrate this behavior.

Example 1:

Suppose you enter a local command while *remsh* is running a remote command that requires no input. Normally, the command would go into your type-ahead buffer and would be executed as soon as the *remsh* command finished. Instead, *remsh* reads the local command as input and the local command never executes (your local shell never gets the command).

Example 2:

This example involves shell scripts. Suppose you had a file named *text* with the following lines in it:

firstline secondline thirdline

and suppose that you wrote the following shell script named test:

```
#!/bin/sh
remsh hpabsb sleep 3
grep "second"
.
.
.
```

If you executed:

test < text

you would expect the shell script to find and display the line

second line

but instead the script displays nothing. This is because any command in the shell script (including *remsh*) inherits *stdin*, which is the input file *text*. Therefore *remsh* reads the file *text* as input and the following *grep* command never sees the file.

Example 3 (for ksh and csh only):

Suppose you put the following command in the background:

remsh hpabsb echo hello &

Instead of seeing

hello

you see the following message after you enter the next carriage return:

[1] + stopped (tty input) remsh hpabsb echo hello

The *remsh* in the background tries to read its *stdin* (your terminal input). Since the shell does not allow background processes to read your terminal, the shell stops the background process, and notifies you.

In all 3 examples, to prevent such mishaps, *remsh* provides an option, -n, that redirects *remsh*'s input from */dev/null*. Whenever you use *remsh* to run a remote command that requires no input, it is good practice to invoke *remsh* with the -n option.

At your HP-UX prompt, enter:

remsh remote_host -n command

where *remote_host* is the name or alias of a host listed in */etc/hosts* and *command* is a command requiring no input to execute on the remote host.

Example Entries:

remsh hpabsb -n who remsh hpabsb -n sleep 3 remsh hpabsb -n echo hello &

Result:

Remsh executes the command on the remote host, taking the command's input from /*dev/null*.

Using Remsh's "Shorthand" Syntax

Your local host can be configured so that you can enter a *remsh* command line without the *remsh* command. That is, a *remsh* command line can start with the name of a remote host, omitting the *remsh* command.

In order to do this:

- you must add the path */usr/hosts* to your command search path in your *.login, .cshrc*, or *.profile* file. Which file contains your *\$PATH* variable depends on which shell you use.
- the super-user or node manager must link */usr/bin/remsh* to */usr/hosts/host*, where *host* is the name or alias of a remote host (listed in */etc/hosts*) on which you want to execute a command.

To find out which hosts you can use *remsh*'s shorthand syntax for, list the contents of the directory */usr/hosts*.

10-18 Using Remsh's "Shorthand" Syntax

11

Interprocess Communication

This chapter describes HP's implementation of the 4.2 BSD Interprocess Communication (IPC) facilities. The chapter includes the following sections:

- an IPC overview using the Client-Server model;
- a description of important terms and concepts;
- the details of IPC using stream sockets;
- advanced IPC concepts for stream sockets;
- the details of IPC using datagram sockets;
- advanced IPC concepts for datagram sockets;
- a list of programming hints;
- how to add a server process to inetd; and
- tables of the available system and library calls.

Note

IPC is a program development tool. Before you attempt to use IPC, you may need to familiarize yourself with the C programming language and the HP-UX operating system. You could implement an IPC application using Fortran or Pascal, but all library calls and include files are implemented in C.

Overview of IPC

The IPC facility allows you to create distributed applications that pass data between programs (on the same computer or on separate computers on the network) without requiring an understanding of the many layers of networking protocols. This is accomplished by using a set of system calls. These system calls, when used in the correct sequence, allow you to create communication endpoints called **sockets** and transfer data between them.

This chapter describes the steps involved in establishing and using IPC connections. It also describes the protocols you must use and how the IPC system calls interact. The details of each system call are described in the section 2 entries of the *ARPA/Berkeley Services Reference Pages*.

To understand the general model for IPC, you need to understand what is meant by a **socket**, a **socket descriptor**, and **binding**. Read the following definitions before you read about the Client-Server model.

socket	Sockets are communication endpoints. A pair of con- nected sockets provides an interface similar to that of HP-UX pipes. A socket is identified by a socket descriptor.
socket descriptor	A socket descriptor is an HP-UX file descriptor that references a socket instead of an ordinary file. There- fore, it can be used for reading, writing, or most standard file system calls after an IPC connection is established. All IPC functions use socket descriptors as arguments.
binding	Before a socket can be accessed across the network, it must be bound to an address. Binding makes the socket accessible to other sockets on the network by

How You Can Use IPC

The best example of how IPC can be used is the ARPA/Berkeley Services themselves. The services use IPC to communicate between remote hosts. Using the IPC facility, you can write your own distributed application programs to do a variety of tasks.

detail throughout this chapter.

establishing its address. Binding is explained in more

For example, you can write distributed application programs to:

- access a remote database;
- · access multiple computers at one time; or
- spread subtasks across several hosts.

The Client-Server Model

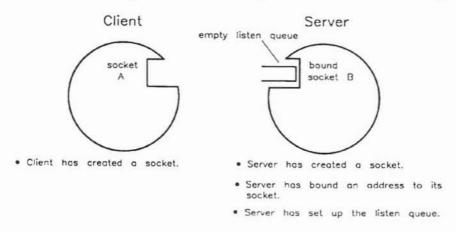
Typical IPC applications consist of two separate application level processes; one process (the **client**) requests a connection and the other process (the **server**) accepts it.

The server process creates a socket, binds an address to it, and sets up a mechanism (called a listen queue) for receiving connection requests. The client process creates a socket and requests a connection to the server process. Once the server process accepts a client process's request and a connection is established, full-duplex (two-way) communication can occur between the two sockets.

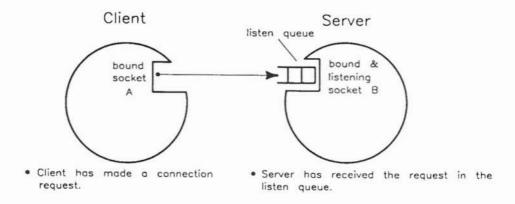
This set of conventions must be implemented by both processes. Depending on the needs of your application, your implementation of the model can be symmetric or asymmetric. In a symmetrical application of the model, either process can be a server or a client. In an asymmetrical application of the model, there is a clearly defined server process and client process. An example of an asymmetrical application is the *ftp* service.

Creating a Connection: the Client-Server Model

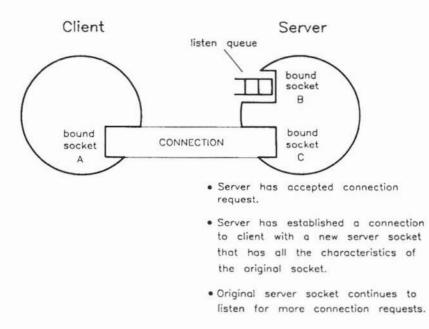
The following figures illustrate conceptual views of the client-server model at three different stages of establishing a connection. The steps that have been accomplished at each stage are listed below each figure.



Client-Server in a Pre-Connection State



Client-Server at Time of Connection Request



Client-Server When Connection Is Established

A detailed description of the Client-Server model is discussed in the "IPC Using Stream Sockets" section of this chapter.

IPC Library Routines

The library routines and system calls that you need to implement an IPC application are described throughout this chapter. In addition, a complete list of all these routines and system calls is provided in the "Summary Tables for Library and System Calls" section of this chapter.

The library routines are in the common "c" library named *libc.a.* Therefore, there is no need to specify any library name on the *cc* command line to use these library calls - *libc.a* is used automatically.

Key Terms and Concepts

The following list is meant to give you a basic understanding of the terms used to describe IPC. Many of the terms have more detailed explanations within this chapter in the places where the terms are used.

Communication Terms

packet	A message or data unit that is transmitted between communicating processes.	
message	The data sent in one UDP packet.	
channel	Communication path created by establishing a con- nection between sockets.	
peer	The remote process with which a process communicates.	

Addressing Terms

addressing	A means of labeling a socket so that it is distinguish- able from other sockets on a host.	
communication domain	A set of properties that describes the characteristics of processes communicating through sockets. The AF_INET (internet address family) domain is sup- ported. The AF_UNIX (UNIX address family) domain is also supported, for local communication only.	
address family	The address format used to interpret addresses specified in socket operations. The internet address family (AF_INET) and Berkeley UNIX address fami- ly (AF_UNIX) are supported.	
internet address	A four-byte address that identifies a node on the net- work.	

port	An address within a host that is used to differentiate between multiple sockets with the same internet ad- dress. You can use port address values 1024 through 65535. (Port addresses 1 through 1023 are reserved for the super-user.)
socket address	For AF_INET, the socket address consists of the in- ternet address, port address and address family of a socket. The internet and port address combination al- lows the network to locate a socket. For AF_UNIX, the socket address is the directory path name of the vnode bound to the socket.
binding	Associating a socket address with a socket. Once a socket address is bound, other sockets can connect to the socket and send data to or receive data from it.
association	An IPC connection is defined by an association. An AF_INET association contains the (protocol, local address, local port, remote address, remote port)-tuple. An AF_UNIX association contains the (protocol, local address, peer address)-tuple. Associa- tions must be unique; duplicate associations on the same host cannot exist. The tuple is created when the local and remote socket addresses are bound and connected. This means the association is created in two steps and there is a chance that two associations could be alike. The host prevents this by checking for uniqueness of the tuple at connection time and reporting an error if the tuple is not unique.

Protocols

There are two Internet transport layer protocols that can be used with IPC. They are TCP, which implements stream sockets, and UDP, which implements datagram sockets.

TCP Provides the underlying communication support for stream sockets. TCP is used to implement reliable, sequenced, flow-controlled two-way communication based on byte streams similar to pipes. Refer to the *TCP(7P)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on TCP.

UDP Provides the underlying communication support for datagram sockets. UDP is an unreliable protocol. A process receiving messages on a datagram socket could find messages are duplicated, out-of-sequence, or missing. Messages retain their record boundaries and are sent as individually addressed packets. There is no concept of a connection between the communicating sockets. Refer to the UDP(7P) entry in the ARPA/Berkeley Services Reference Pages for more information on UDP.

In addition, the UNIX domain protocol may be used with AF_UNIX sockets for interprocess communication on the same node. Refer to the unix(7p) entry in the LAN reference pages for more information on the UNIX domain protocol.

Using Socket Descriptors as File Descriptors

A socket descriptor is a special kind of HP-UX file descriptor; it can be used as though it were a file descriptor, but it references a socket instead of a file. System calls that use file descriptors (e.g. *read*, *write*, *select*) can be used with socket descriptors.

IPC Using Internet Stream Sockets

This section describes the steps involved in creating an Internet stream socket IPC connection using the AF_INET address family. If you want to use datagram sockets, skip to the section called "IPC Using Datagram Sockets."

As discussed in the "Protocols" section, Internet TCP stream sockets provide bidirectional, reliable, sequenced and unduplicated flow of data without record boundaries.

The following table lists the steps involved in creating and terminating an IPC connection using stream sockets. Each step is described in more detail in the sections that follow the table.

Client Process Activity	System call used	Server Process Activity	System call used
Activity	System can used	Activity	System can used
create a socket	socket()	create a socket	socket()
bind a socket address (optional)	bind()	bind a socket address	bind()
		listen for incoming connection requests	
request a connection	connect()		
		accept connection	accept()
send data	write() or send()		
		receive data	read() or recv()
		send data	write() or send()
receive data	read() or recv()		
disconnect socket (optional)	<pre>shutdown() or close()</pre>	disconnect socket (optional)	shutdown() or close()

Building an IPC Connection Using Stream Sockets

11-10 IPC Using Internet Stream Sockets

The following sections explain each of the activities mentioned in the previous table. The description of each activity specifies a system call and includes:

- what happens when the system call is used;
- when to make the call;
- what the parameters do;
- how the call interacts with other IPC system calls; and
- where to find details on the system call.

The stream socket program examples are at the end of these descriptive sections. You can refer to the example code as you work through the descriptions.

Preparing Address Variables

Before you begin to create a connection, establish the correct variables and collect the information that you need to request a connection.

Your server process needs to:

- declare socket address variables;
- assign a wildcard address; and
- get the port address of the service that you want to provide.

Your client process needs to:

- declare socket address variables;
- get the remote host's internet address; and
- get the port address for the service that you want to use.

These activities are described next. Refer to the program example at the end of the "IPC Using Stream Sockets" section to see how these activities work together.

Declaring Socket Address Variables

You need to declare a variable of type struct *sockaddr_in* to use for socket addresses.

For example, the following declarations are used in the example client program:

```
struct sockaddr_in myaddr; /* for local socket address */
struct sockaddr_in peeraddr; /* for peer socket address */
```

Sockaddr_in is a special case of sockaddr and is used with the AF_INET addressing domain. Both types are shown in this chapter, but sockaddr_in makes it easier to manipulate the internet and port addresses. Some of the IPC system calls are declared using a pointer to sockaddr, but it can also be a pointer to sockaddr_in.

The sockaddr_in address structure consists of the following fields:

short sin_family	Specifies the address family and should al- ways be set to AF_INET.
u_short <i>sin_port</i>	Specifies the port address. Assign this field when you bind the port address for the socket or when you get a port address for a specific service.
struct inaddr sin_addr	Specifies the internet address. Assign this field when you get the internet address for the remote host.

The server process only needs an address for its own socket. Your client process may not need an address for its local socket.

Refer to the *inet(7F)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *sockaddr_in*.

Getting the Remote Host's Internet Address

Gethostbyname obtains the internet address of the host and the length of that address (as the size of struct *in_addr*) from */etc/hosts*.

Gethostbyname and its parameters are described in the following table.

INCLUDE FILES: #include <netdb.h SYSTEM CALL: struct hostent *g char *name;</netdb.h 		>		
		<pre>struct hostent *gethostbyname(name) char *name;</pre>		
Parameter	Descrip	tion of Contents	INPUT Value	
name	pointer to a valid host name (null-terminated string)		host name	
FUNCTION RESULT:		address	r (0) if failure occurs	

EXAMPLE SYSTEM CALL:

```
#include <netdb.h>
struct hostent *hp; /* pointer to host info for remote host */
...
peeraddr.sin_family = AF_INET;
hp = gethostbyname (argv[1]);
peeraddr_in.sin_addr.s_addr = ((struct in_addr *)(hp->h_addr))->s_addr;
```

The *argv[1]* parameter is the host name specified in the client program command line.

Refer to the gethostent(3N) entry in the ARPA/Berkeley Services Reference Pages for more information on gethostbyname.

Getting the Port Address for the Desired Service

When a server process is preparing to offer a service, it must get the port address for the service from */etc/services* so it can bind that address to its "listen" socket. If the service is not already in */etc/services*, you must add it.

When a client process needs to use a service that is offered by some server process, it must request a connection to that server process's "listening" socket. The client process must know the port address for that socket.

Getservbyname obtains the port address of the specified service from *letc/services*.

Getservbyname and its parameters are described in the following table.

INCLUDE FI	LES:	#include <netdb.h></netdb.h>	
SYSTEM CA	LL:	struct servent *ge char *name, *proto	etservbyname(name, proto) o;
Parameter	Descrip	tion of Contents	INPUT Value
name	pointer to a valid service name		service name
proto	pointer to the protocol to be used		"tcp" or 0 if TCP is the only protocol for the service
FUNCTION	RESULT:	pointer to strue dress	ct servent containing port ad-

EXAMPLE SYSTEM #include <netdb.h> CALL: #include <netdb.h> struct servent *sp; /* pointer to service info */ ... sp = getservbyname ("example", "tcp"); peeraddr.sin_port = sp->s_port;

NULL pointer (0) if failure occurs

When to Get Server's Socket Address

Which Processes	When
server process	before binding the listen socket
client process	before client executes a connection request

Refer to the getservent(3N) entry in the ARPA/Berkeley Services Reference Pages for more information on getservbyname.

Using a Wildcard Local Address

Wildcard addressing simplifies local address binding. When an address is assigned the value of INADDR_ANY, the host interprets the address as any valid address. This is useful for your server process when you are setting up the listen socket. It means that the server process does not have to look up its own internet address.

For example, to bind a specific port address to a socket, but leave the local internet address unspecified, the following source code could be used:

```
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
...
struct sockaddr_in sin;
...
s = socket(AF_INET, SOCK_STREAM, 0);
sin.sin_family = AF_INET;
sin.sin_addr.s_addr = INADDR_ANY;
sin.sin_port = MYPORT;
bind (s, &sin, sizeof(sin));
```

Writing the Server Process

This section discusses the calls your server process must make to connect with and serve a client process.

Creating a Socket

The server process must call socket to create a communication endpoint.

Socket and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>s = socket(af, type, protocol)</pre>
	int af, type, protocol;

Parameter	Description of Contents		INPUT Value	
af	address family		AF_INET	
type	socket type		SOCK_STREAM	
protocol	underlying protocol to be used		0 (default) or value returned by <i>getprotobyname</i>	
FUNCTION RESULT:		socket number –1 if failure oc	(HP-UX file descriptor) curs	
EXAMPLE SYSTEM s = socket CALL:		s = socket (AF_INE	ET, SOCK_STREAM, 0);	

The socket number returned is the socket descriptor for the newly created socket. This number is an HP-UX file descriptor and can be used for reading, writing or any standard file system calls after an IPC connection is established. A socket descriptor is treated like a file descriptor for an open file.

When to Create Sockets

Which Processes	When
-----------------	------

server process

before any other IPC system calls

Refer to the *socket(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *socket*.

Binding a Socket Address to the Server Process's Socket

After your server process has created a socket, it must call *bind* to bind a socket address. Until an address is bound to the server socket, other processes have no way to reference it.

The server process must bind a specific port address to this socket, which is used for listening. Otherwise, a client process would not know what port to connect to for the desired service.

Set up the address structure with a local address (as described in the "Preparing Address Variables" section) before you make a *bind* call. Use a wildcard address so your server process does not have to look up its own internet address.

Bind and its parameters are described in the following table.

int addrlen:

INCLUDE FILES: #include <sys/types.h> #include <netinet/in.h> #include <sys/socket.h> SYSTEM CALL: bind (s, addr, addrlen) int s; struct sockaddr *addr;

INPUT Value Parameter **Description of Contents** S socket descriptor of local socket descriptor of socket to be bound socket addr socket address pointer to address to be bound to s addrlen length of socket address size of struct sockaddr in FUNCTION RESULT: 0 if bind is successful -1 if failure occurs EXAMPLE SYSTEM struct sockaddr in myaddr; CALL: bind (ls, myaddr, sizeof(struct sockaddr_in));

When to Bind Socket Addresses

Which Processes	When	
server process	after socket is created and before any other IPC system calls	

Refer to the *bind(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *bind*.

Setting the Server Up to Wait for Connection Requests

Once your server process has an address bound to it, it must call *listen* to set up a queue that accepts incoming connection requests. The server process then monitors the queue for requests (using *select(2)* or *accept*, which is described in "Accepting a Connection"). The server process cannot respond to a connection request until it has executed *listen*.

Listen and its parameters are described in the following table.

INCLUDE FILES: none SYSTEM CALL: listen(s, backlog) int s, backlog; Parameter Description of Contents **INPUT Value** socket descriptor of local server socket's descriptor S socket backlog maximum number of connecsize of queue (between 1 and 20) tion requests in the queue at any time FUNCTION RESULT: 0 if listen is successful -1 if failure occurs EXAMPLE SYSTEM listen (1s, 5); CALL:

Backlog is the number of unaccepted incoming connections allowed at a given time. Further incoming connection requests are rejected.

When to Set Server Up to Listen

Which Processes	When
server process	after socket is created and bound and before the server can respond
	to connection requests

Refer to the *listen(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *listen*.

Accepting a Connection

The server process can accept any connection requests that enter its queue after it executes *listen*. Accept creates a new socket for the connection and returns the socket descriptor for the new socket. The new socket:

- is created with the same properties as the old socket;
- has the same bound port address as the old socket; and
- is connected to the client process' socket.

Accept blocks until there is a connection request from a client process in the queue.

Accept and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>	
	<pre>#include <netinet in.h=""></netinet></pre>	
	<pre>#include <sys socket.h=""></sys></pre>	
SYSTEM CALL:	<pre>s = accept(s,addr,addrlen)</pre>	
	int s;	
	struct sockaddr *addr;	
	<pre>int *addrlen;</pre>	

Parameter	Description of Co	ntents	INPUT Value	OUTPUT Value
S	socket descriptor of local socket	of	socket descriptor of server socket	unchanged
addr	socket address		pointer to address structure where address will be put	pointer to socket address of client socket that server's new socket is connected to
addrlen	length of address		pointer to the size of struct sockaddr_in	pointer to the actual length of address returned in addr
succe		succe	et descriptor of new socket if accept is essful failure occurs	
EXAMPL CALL:	E SYSTEM	 addrler	<pre>sockaddr_in peeraddr; n = sizeof(sockaddr_in); cept (ls, peeraddr, &addr</pre>	·len);

There is no way for the server process to indicate which requests it can accept. It must accept all requests or none. Your server process can keep track of which process a connection request is from by examining the address returned by *accept*. Once you have this address, you can use *gethostbyaddr* to get the host name. You can close down the connection if you do not want the server process to communicate with that particular client host or port.

When to Accept a Connection

Writing the Client Process

This section discusses the calls your client process must make to connect with and be served by a server process.

Creating a Socket

The client process must call socket to create a communication endpoint.

Socket and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
OVOTEN (CALL	

SYSTEM CALL:

s = socket(af, type, protocol)

int af, type, protocol;

Parameter	Description of Contents		INPUT Value
af	address fan	nily	AF_INET
type	socket type		SOCK_STREAM
protocol	underlying used	protocol to be	0 (default) or value returned by <i>getprotobyname</i>
FUNCTION	RESULT:	socket numbe -1 if failure o	er (HP-UX file descriptor)
EXAMPLE S	YSTEM	s = socket (AF_1	INET, SOCK_STREAM, 0);

The socket number returned is the socket descriptor for the newly created socket. This number is an HP-UX file descriptor and can be used for reading, writing or any standard file system calls after an IPC connection is established. A socket descriptor is treated like a file descriptor for an open file.

When to Create Sockets

Which Processes	When

client process

before requesting a connection

Refer to the *socket(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *socket*.

Requesting a Connection

Once the server process is listening for connection requests, the client process can request a connection with the *connect* call.

Connect and its parameters are described in the following table.

Parameter	Description of Contents	INPUT Value
	int addrlen;	
	struct sockaddr *addr;	
	int s;	
SYSTEM CALL:	<pre>connect(s, addr, addrlen)</pre>	
	<pre>#include <sys socket.h=""></sys></pre>	
	<pre>#include <netinet in.h=""></netinet></pre>	
INCLUDE FILES	<pre>#include <sys types.h=""></sys></pre>	

Parameter	Description of Contents	INPUT Value	
S	socket descriptor of local socket	socket descriptor of socket request- ing a connection	
addr	pointer to the socket address	pointer to the socket address of the socket to which client wants to con- nect	
addrlen	length of addr	size of address structure pointed to by addr	
FUNCTION	RESULT: 0 if connect is	successful	

CTION RESULT: 0 if connect is successful -1 if failure occurs

EXAMPLE SYSTEM	<pre>struct sockaddr_in peeraddr;</pre>
CALL:	
	connect (s. peeraddr. sizeof(struct sockaddr in)):

Connect initiates a connection and blocks if the connection is not ready, unless you are using nonblocking I/O. (For information on nonblocking I/O, see the "Advanced Topics for Stream Sockets: Nonblocking I/O" section of this chapter.) When the connection is ready, the client process completes its *connect* call and the server process can complete its *accept* call.

Note

The client process does not get feedback that the server process has completed the *accept* call. As soon as the *connect* call returns, the client process can send data.

Note

Local internet and port addresses are bound when *connect* is executed if you have not already bound them yourself. These address values are chosen by the local host.

When to Request a Connection

Which Processes	When	
client process	after socket is created and after server socket has a listening socket	

Refer to the connect(2) entry in the ARPA/Berkeley Services Reference Pages for more information on connect.

Sending and Receiving Data

After the *connect* and *accept* calls are successfully executed, the connection is established and data can be sent and received between the two socket endpoints. Because the stream socket descriptors correspond to HP-UX file descriptors, you can use the *read* and *write* calls (in addition to *recv* and *send*) to pass data through a socket-terminated channel.

If you are considering the use of the *read* and *write* system calls instead of the *send* and *recv* calls described below, you should consider the following:

Advantage:	If you use <i>read</i> and <i>write</i> instead of <i>send</i> and <i>recv</i> , you can use a socket for <i>stdin</i> or <i>stdout</i> .
Disadvantage:	If you use <i>read</i> and <i>write</i> instead of <i>send</i> and <i>recv</i> , you cannot use the options specified with the <i>send</i> or <i>recv flags</i> parameter.

See the table called "Other System Calls," listed at the end of the chapter for more information on which of these system calls are best for your application.

Sending Data

Send and its parameters are described in the following table.

INCLUDE FILES	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>count = send(s,msg,len,flags)</pre>
	int s;
	char *msg;
	int len, flags;
Parameter	Description of Contents

Parameter	Descript	ion of Contents	INPUT Value
s	socket descr socket	iptor of local	socket descriptor of socket sending data
msg	pointer to d	ata buffer	pointer to data to be sent
len	size of data	buffer	size of msg
flags	settings for	optional flags	0 or MSG_OOB
FUNCTION	RESULT:	number of by -1 if failure of	ytes actually sent
EXAMPLE S	YSTEM	count = send (s	, buf, 10, 0);

TAIDLIT X7 1

Send blocks until the specified number of bytes have been queued to be sent, unless you are using nonblocking I/O. (For information on nonblocking I/O, see the "Advanced Topics for Stream Sockets: Nonblocking I/O" section of this chapter.)

When to Send Data

Which Processes	When	
server or client	after connection is established	
process		

Refer to the send(2) entry in the ARPA/Berkeley Services Reference Pages for more information on send.

Receiving Data

Recv and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>count = recv(s,buf,len,flags)</pre>
	int s;
	char *buf;
	int len, flags;

Parameter	Description of Contents socket descriptor of local socket		socket descriptor of socket receiv- ing data	
s				
buf	pointer to d	ata buffer	pointer to buffer that is to receive data	
len	maximum n that should	umber of bytes be received	size of data buffer	
flags	settings for	optional flags	0, MSG_OOB or MSG_PEEK	
FUNCTION RESULT:		number of bytes actually received –1 if failure occurs		
EXAMPLE SYSTEM CALL:		count = recv(s,	buf, 10, 0);	

Recv blocks until there is at least 1 byte of data to be received, unless you are using nonblocking I/O. (For information on nonblocking I/O, see the "Advanced Topics for Stream Sockets: Nonblocking I/O" section of this chapter.) The host does not wait for *len* bytes to be available; if less than *len* bytes are available, that number of bytes are received.

No more than *len* bytes of data are received. If there are more than *len* bytes of data on the socket, the remaining bytes are received on the next *recv*.

Flag Options

The flags options are:

- 0 for no options;
- MSG_OOB for out of band data; or
- MSG_PEEK for a nondestructive read.

Use the MSG_OOB option if you want to receive out of band data. Refer to the "Advanced Topics for Stream Sockets, Sending and Receiving Out of Band Data" section of this chapter for more information.

Use the MSG_PEEK option to preview incoming data. If this option is set on a *recv*, any data returned remains in the socket buffer as though it had not been read yet. The next *recv* returns the **same data**.

When to Receive Data

When	
after connection is established	

Refer to the recv(2) entry in the ARPA/Berkeley Services Reference Pages for more information on recv.

Closing a Socket

In most applications, you do not have to worry about cleaning up your sockets. When you exit your program and your process terminates, the sockets are closed for you.

If you need to close a socket while your program is still running, use the *close* system call. For example, you may have a daemon process that uses *fork* to create the server process. The daemon process creates the IPC connection and then passes the socket descriptor to the server. You then have more than one process with the same socket descriptor. The daemon process should do a *close* of the socket descriptor to avoid keeping the socket open once the server is through with it. Because the server performs the work, the daemon does not use the socket after the *fork*.

Close decrements the file descriptor count and the calling process can no longer use that file descriptor.

When the last *close* is executed on a socket descriptor, any unsent data are sent before the socket is closed. Any unreceived data are lost. This delay in closing the socket can be controlled by the socket option SO_LINGER. See the "Socket Options" section for information on the SO_LINGER and SO_DONTLINGER options.

For syntax and details on *close*, refer to the *close(2)* entry in the *HP-UX Reference* manual.

Additional options for closing sockets are discussed in the "Advanced Topics for Stream Sockets: Using Shutdown" section of this chapter.

Example Using Stream Sockets

These program examples demonstrate how to set up and use stream sockets. The client program is intended to run in conjunction with the server program. The client program requests a service called *example* from the server program.

The server process receives requests from the remote client process, handles the request and returns the results to the client process. Note that the server:

- uses the wildcard address for the listen socket;
- uses the *ntohs* address conversion call to show how to port to a host that requires it; and
- uses the SO_LINGER option for a graceful disconnect. The SO_LINGER options is discussed in the "Socket Options" section, which follows the example.

The client process creates a connection, sends requests to the server process and receives the results from the server process. Note that the client:

- uses *shutdown*, which is discussed in the "Advanced Topics for Stream Sockets" section of this chapter, to indicate that it is done sending requests;
- uses *getsockname* to see what socket address was assigned to the local socket by the host; and
- uses the *ntohs* address conversion call to show how to port to a host that requires it.

Before you run the example programs:

• make the following entry in the two host's /etc/services files:

example 22375/tcp

• compile the programs with the -lbsdipc option.

The source code for these two programs follows. It is also located in the directory */usr/netdemo/socket*.

```
/*
                        SERV.TCP
 *
 *
        This is an example program that demonstrates the use of
 *
        stream sockets as an IPC mechanism. This contains the server,
 *
        and is intended to operate in conjunction with the client
 *
        program found in client.tcp. Together, these two programs
 *
        demonstrate many of the features of sockets, as well as good
 *
        conventions for using these features.
 *
 *
        This program provides a service called "example". In order for
 *
        it to function, an entry for it needs to exist in the
 *
        /etc/services file. The port address for this service can be
 *
        any port number that is likely to be unused, such as 22375.
 *
        The host on which the client will be running
 *
        must also have the same entry (same port number) in its
 *
        /etc/services file.
 *
*/
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
#include <signal.h>
#include <stdio.h>
#include <netdb.h>
                                /* connected socket descriptor */
int s:
                               /* listen socket descriptor */
int ls:
struct hostent *hp;
                                /* pointer to host info for remote host */
                               /* pointer to service information */
struct servent *sp;
                               /* contains time returned by time() */
long timevar;
char *ctime();
                               /* declare time formatting routine */
long linger = 1:
                               /* allow a lingering, graceful close */
                                /* used when setting SO LINGER */
struct sockaddr_in myaddr_in; /* for local socket address */
struct sockaddr in peeraddr in; /* for peer socket address */
```

```
/*
                                        MAIN
        This routine starts the server. It forks, leaving the child
        to do all the work, so it does not have to be run in the
 *
        background. It sets up the listen socket, and for each incoming
        connection, it forks a child process to process the data. It
 *
        will loop forever, until killed by a signal.
 */
main(argc, argv)
int argc;
char *argv[];
        int addrlen;
                /* clear out address structures */
        memset ((char *)&myaddr_in, 0, sizeof(struct sockaddr_in));
        memset ((char *)&peeraddr in, 0, sizeof(struct sockaddr in));
                /* Set up address structure for the listen socket. */
        myaddr in.sin_family = AF_INET;
                /* The server should listen on the wildcard address,
                 * rather than its own internet address. This is
                 * generally good practice for servers, because on
                 * systems which are connected to more than one
                 * network at once will be able to have one server
                 * listening on all networks at once. Even when the
                 * host is connected to only one network, this is good
                 * practice, because it makes the server program more
                 * portable.
                 */
        myaddr in.sin addr.s addr = INADDR ANY;
                /* Find the information for the "example" server
                 * in order to get the needed port number.
                 */
        sp = getservbyname ("example", "tcp");
        if (sp == NULL) {
                fprintf(stderr, "%s: example not found in /etc/services\n",
                                argv[0]);
                exit(1);
        myaddr in.sin port = sp->s port;
                /* Create the listen socket. */
        1s = socket (AF INET, SOCK STREAM, 0);
        if (ls == -1) {
                perror(argv[0]);
                fprintf(stderr, "%s: unable to create socket\n", argv[0]);
                exit(1):
        }
```

```
/* Bind the listen address to the socket. */
if (bind(ls. &myaddr in, sizeof(struct sockaddr in)) == -1) {
        perror(argv[0]):
        fprintf(stderr, "%s: unable to bind address\n", argv[0]);
        exit(1);
}
        /* Initiate the listen on the socket so remote users
         * can connect. The listen backlog is set to 5. 20
         * is the currently supported maximum.
         */
if (listen(ls, 5) == -1) {
        perror(argv[0]);
        fprintf(stderr, "%s: unable to listen on socket\n", argv[0]);
        exit(1);
}
        /* Now, all the initialization of the server is
         * complete, and any user errors will have already
         * been detected. Now we can fork the daemon and
         * return to the user. We need to do a setpgrp
         * so that the daemon will no longer be associated
         * with the user's control terminal. This is done
         * before the fork, so that the child will not be
         * a process group leader. Otherwise, if the child
         * were to open a terminal, it would become associated
         * with that terminal as its control terminal. It is
         * always best for the parent to do the setpgrp.
         */
setpgrp();
switch (fork()) {
case -1:
                        /* Unable to fork. for some reason. */
        perror(argv[0]);
        fprintf(stderr, "%s: unable to fork daemon\n", argv[0]);
        exit(1);
case 0:
                        /* The child process (daemon) comes here. */
                /* Close stdin and stderr so that they will not
                 * be kept open. Stdout is assumed to have been
                 * redirected to some logging file, or /dev/null.
                 * From now on, the daemon will not report any
                 * error messages. This daemon will loop forever,
                 * waiting for connections and forking a child
                 * server to handle each one.
                 */
        fclose(stdin):
        fclose(stderr);
```

```
/* Set SIGCLD to SIG IGN, in order to prevent
                 * the accumulation of zombies as each child
                 * terminates. This means the daemon does not
                 * have to make wait calls to clean them up.
                 */
        signal(SIGCLD, SIG IGN);
        for(::) {
                        /* Note that addrlen is passed as a pointer
                         * so that the accept call can return the
                         * size of the returned address.
                         */
                addrlen = sizeof(struct sockaddr in);
                        /* This call will block until a new
                         * connection arrives. Then, it will
                         * return the address of the connecting
                         * peer, and a new socket descriptor, s,
                         * for that connection.
                         */
                s = accept(ls, &peeraddr_in, &addrlen);
                if ( s == -1) exit(1);
                switch (fork()) {
                                /* Can't fork, just continue. */
                case -1:
                        exit(1);
                                /* Child process comes here. */
                case 0:
                        server();
                        exit(0);
                                /* Daemon process comes here. */
                default:
                                /* The daemon needs to remember
                                 * to close the new accept socket
                                 * after forking the child. This
                                 * prevents the daemon from running
                                 * out of file descriptors. It
                                 * also means that when the server
                                 * closes the socket, that it will
                                 * allow the socket to be destroyed
                                 * since it will be the last close.
                                 */
                        close(s);
                }
        }
default:
                        /* Parent process comes here. */
       exit(0);
}
```

}

```
1*
 *
                                SERVER
 * *
        This is the actual server routine that the daemon forks to
*
        handle each individual connection. Its purpose is to receive
 *
        the request packets from the remote client, process them,
 *
        and return the results to the client. It will also write some
 *
        logging information to stdout.
 *
*/
server()
{
                                /* keeps count of number of requests */
        int reacht = 0;
                                /* This example uses 10 byte messages. */
        char buf[10];
        char *inet ntoa();
        char *hostname:
                                /* points to the remote host's name string */
        int len, len1;
                /* Close the listen socket inherited from the daemon. */
        close(ls):
                /* Look up the host information for the remote host
                 * that we have connected with. Its internet address
                 * was returned by the accept call, in the main
                 * daemon loop above.
                 */
        hp = gethostbyaddr ((char *) &peeraddr in.sin addr,
                                sizeof (struct in addr),
                                peeraddr in.sin family);
        if (hp == NULL) {
                        /* The information is unavailable for the remote
                         * host. Just format its internet address to be
                         * printed out in the logging information. The
                         * address will be shown in "internet dot format".
                         */
                hostname = inet ntoa(peeraddr in.sin addr);
        } else {
                hostname = hp->h name; /* point to host's name */
        }
                /* Log a startup message. */
        time (&timevar);
```

11-36 IPC Using Internet Stream Sockets

/* The port number must be converted first to host byte * order before printing. On most hosts, this is not * necessary, but the ntohs() call is included here so * that this program could easily be ported to a host * that does require it. */ printf("Startup from %s port %u at %s". hostname, ntohs(peeraddr in.sin port), ctime(&timevar)); /* Set the socket for a lingering, graceful close. * Since linger was set to 1 above, this will cause * a final close of this socket to wait until all of the * data sent on it has been received by the remote host. */ if (setsockopt(s, SOL SOCKET, SO LINGER, (char *)&linger. sizeof(long)) == -1) { printf("Connection with %s aborted on error\n", hostname): errout: exit(1): } /* Go into a loop, receiving requests from the remote * client. After the client has sent the last request. * it will do a shutdown for sending, which will cause * an end-of-file condition to appear on this end of the * connection. After all of the client's requests have * been received, the next recy call will return zero * bytes, signalling an end-of-file condition. This is * how the server will know that no more requests will * follow, and the loop will be exited. */ while (len = recv(s. buf, 10, 0)) { if (len == -1) goto errout: /* error from recv */ /* The reason this while loop exists is that there * is a remote possibility of the above recv returning * less than 10 bytes. This is because a recv returns * as soon as there is some data, and will not wait for * all of the requested data to arrive. Since 10 bytes * is relatively small compared to the allowed TCP * packet sizes, a partial receive is unlikely. If * this example had used 2048 bytes requests instead. * a partial receive would be far more likely. * This loop will keep receiving until all 10 bytes * have been received, thus guaranteeing that the * next recv at the top of the loop will start at * the beginning of the next request. */

```
while (len < 10) {
                len1 = recv(s, &buf[len], 10-len, 0);
                if (len1 == -1) goto errout;
                len += len1:
        }
                /* Increment the request count. */
        regcnt++;
                /* This sleep simulates the processing of the
                 * request that a real server might do.
                 */
        sleep(1);
                /* Send a response back to the client. */
        if (send(s, buf, 10, 0) != 10) goto errout;
}
        /* The loop has terminated, because there are no
         * more requests to be serviced. As mentioned above,
        * this close will block until all of the sent replies
        * have been received by the remote host. The reason
         * for lingering on the close is so that the server will
         * have a better idea of when the remote has picked up
         * all of the data. This will allow the start and finish
         * times printed in the log file to reflect more accurately
         * the length of time this connection was used.
        */
close(s):
        /* Log a finishing message. */
time (&timevar);
        /* The port number must be converted first to host byte
         * order before printing. On most hosts, this is not
         * necessary, but the ntohs() call is included here so
         * that this program could easily be ported to a host
         * that does require it.
         */
printf("Completed %s port %u, %d requests, at %s\n",
        hostname, ntohs(peeraddr in.sin port), reqcnt, ctime(&timevar));
```

}

```
CLIENT. TCP
```

/*

```
This is an example program that demonstrates the use of stream
        sockets as an IPC mechanism. This contains the client, and is
        intended to operate in conjunction with the server program found
 *
        in servitcp. Together, these two programs demonstrate many of the
 *
        features of sockets, as well as good conventions for using these
 *
        features.
 *
        This program requests a service called "example". In order for it
 *
        to function, an entry for it needs to exist in the /etc/services
 *
        file. The port address for this service can be any port number
 *
        that is likely to be unused, such as 22375. The host
 *
        on which the server will be running must also have the same entry
 *
        (same port number) in its /etc/services file.
 *
*/
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
#include <stdio.h>
#include <netdb.h>
int s:
                                /* connected socket descriptor */
struct hostent *hp:
                                /* pointer to host info for remote host */
struct servent *sp;
                                /* pointer to service information */
long timevar:
                                /* contains time returned by time() */
char *ctime():
                               /* declare time formatting routine */
struct sockaddr in myaddr in; /* for local socket address */
struct sockaddr in peeraddr in; /* for peer socket address */
1*
 *
                        MAIN
 *
 *
        This routine is the client which requests service from the remote
 *
        "example server". It creates a connection, sends a number of
 *
        requests, shuts down the connection in one direction to signal the
 *
        server about the end of data, and then receives all of the responses.
 *
        Status will be written to stdout.
 *
 *
        The name of the system to which the requests will be sent is given
        as a parameter to the command.
 */
```

```
main(argc, argv)
int argc:
char *argv[];
        int addrlen, i, i:
                /* This example uses 10 byte messages. */
        char buf[10]:
        if (argc != 2) {
                fprintf(stderr. "Usage: %s <remote host>\n". argv[0]);
                exit(1):
        }
                /* clear out address structures */
        memset ((char *)&myaddr in. 0. sizeof(struct sockaddr in));
        memset ((char *)&peeraddr in, 0, sizeof(struct sockaddr in));
                /* Set up the peer address to which we will connect. */
        peeraddr_in.sin_family = AF INET;
                /^* Get the host information for the hostname that the
                 * user passed in.
                 *1
        hp = gethostbyname (argv[1]);
                /* argv[1] is the host name. */
        if (hp == NULL) {
                fprintf(stderr, "%s: %s not found in /etc/hosts\n",
                                argv[0], argv[1]);
                exit(1);
        }
        peeraddr in.sin addr.s addr = ((struct in addr *)(hp->h addr))->s addr;
                /* Find the information for the "example" server
                 * in order to get the needed port number.
                 */
        sp = getservbyname ("example", "tcp");
        if (sp == NULL) {
                fprintf(stderr, "%s: example not found in /etc/services\n",
                                argv[0]):
                exit(1);
        3
        peeraddr in.sin port = sp->s port;
                /* Create the socket. */
        s = socket (AF INET, SOCK STREAM, 0);
        if (s == -1) {
                perror(argv[0]);
                fprintf(stderr, "%s: unable to create socket\n", argv[0]);
                exit(1):
        }
```

```
/* Try to connect to the remote server at the address
         * which was just built into peeraddr.
         */
if (connect(s, &peeraddr_in, sizeof(struct sockaddr in)) == -1) {
        perror(argv[0]);
        fprintf(stderr, "%s: unable to connect to remote\n", argv[0]);
        exit(1);
}
        /* Since the connect call assigns a random address
         * to the local end of this connection, let's use
         * getsockname to see what it assigned. Note that
         * addrlen needs to be passed in as a pointer,
         * because getsockname returns the actual length
         * of the address.
         */
addrlen = sizeof(struct sockaddr in);
if (getsockname(s, &myaddr_in, &addrlen) == -1) {
        perror(argv[0]);
        fprintf(stderr, "%s: unable to read socket address\n", argv[0]);
        exit(1);
}
        /* Print out a startup message for the user. */
time(&timevar);
        /* The port number must be converted first to host byte
         * order before printing. On most hosts, this is not
         * necessary, but the ntohs() call is included here so
         * that this program could easily be ported to a host
         * that does require it.
         */
printf("Connected to %s on port %u at %s",
                argv[1], ntohs(myaddr in.sin port), ctime(&timevar));
        /* This sleep simulates any preliminary processing
         * that a real client might do here.
         */
```

```
sleep(5);
```

```
/* Send out all the requests to the remote server.
                * In this case, five are sent, but any random number
                * could be used. Note that the first four bytes of
                * buf are set up to contain the request number. This
                * number will be returned in the reply from the server.
                */
       for (i=1; i<=5; i++) {
                *(int *)buf = i;
                if (send(s, buf, 10, 0) != 10) {
                       fprintf(stderr, "%s: Connection aborted on error ",
                                        argv[0]);
                        fprintf(stderr, "on send number %d\n", i);
                        exit(1):
                }
       }
                /* Now. shutdown the connection for further sends.
                * This will cause the server to receive an end-of-file
                * condition after it has received all the requests that
                * have just been sent, indicating that we will not be
                 * sending any further requests.
                */
        if (shutdown(s, 1) == -1) {
                perror(argv[0]):
                fprintf(stderr, "%s: unable to shutdown socket\n", argv[0]);
                exit(1);
       }
                /* Now, start receiving all of the replys from the server.
                 * This loop will terminate when the recv returns zero,
                 * which is an end-of-file condition. This will happen
                 * after the server has sent all of its replies, and closed
                 * its end of the connection.
                 *1
       while (i = recv(s, buf, 10, 0)) {
                if (i = -1) {
errout:
                        perror(argv[0]);
                        fprintf(stderr, "%s: error reading result\n", argv[0]);
                        exit(1);
                }
```

```
/* The reason this while loop exists is that there
                         * is a remote possibility of the above recv returning
                         * less than 10 bytes. This is because a recv returns
                         * as soon as there is some data, and will not wait for
                         * all of the requested data to arrive. Since 10 bytes
                         * is relatively small compared to the allowed TCP
                         * packet sizes, a partial receive is unlikely. If
                         * this example had used 2048 bytes requests instead,
                         * a partial receive would be far more likely.
                         * This loop will keep receiving until all 10 bytes
                         * have been received, thus guaranteeing that the
                         * next recv at the top of the loop will start at
                         * the beginning of the next reply.
                         */
                while (i < 10) {
                        j = recv(s, &buf[i], 10-i, 0);
                        if (j == -1) goto errout:
                        i += j:
                }
                        /* Print out message indicating the identity of
                         * this reply.
                         */
                printf("Received result number %d\n", *(int *)buf);
        }
                /* Print message indicating completion of task. */
        time(&timevar);
        printf("All done at %s", ctime(&timevar));
}
```

BSD IPC Using UNIX Domain Stream Sockets

This section describes the steps involved in creating a UNIX Domain stream socket BSD IPC connection between two processes executing on the same node. Datagram sockets are not currently supported for UNIX Domain.

UNIX Domain (AF_UNIX) stream sockets provide bidirectional, reliable, unduplicated flow of data without record boundaries. They offer significant performance increases when compared with the use of local Internet (AF_INET) sockets, due primarly to lower code execution overhead.

The following table lists the steps involved in creating and terminating a UNIX Domain BSD IPC connection using stream sockets. Each step is described in more detail in the sections that follow the table.

Client Process Activity	System call used	Server Process Activity	System call used
create a socket	socket()	create a socket	socket()
		bind a socket address	bind()
		listen for incoming connection requests	
request a connection	connect()		
		accept connection	accept()
send data	write() or send()		
		receive data	read() or recv()
		send data	write() or send()
receive data	read() or recv()		
disconnect socket (optional)	shutdown() or close()	disconnect socket (optional)	shutdown() or close()

Building a UNIX Domain BSD IPC Connection Using Stream Sockets

The following sections explain each of the activities mentioned in the previous table. The description of each activity specifies a system call and includes:

- what happens when the system call is used;
- when to make the call;
- what the parameters do;
- · how the call interacts with other BSD IPC system calls; and
- where to find details on the system call.

The UNIX Domain stream socket program examples are at the end of these descriptive sections. You can refer to the example code as you work through the descriptions.

Preparing Address Variables

Before you begin to create a connection, establish the correct variables and collect the information that you need to request a connection.

Your server process needs to:

- · declare socket address variables;
- get the pathname (character string) for the service you want to provide.

Your client process needs to:

- · declare socket address variables;
- get the pathname (character string) for the service you want to use.

These activities are described next. Refer to the program example at the end of this chapter to see how these activities work together.

Declaring Socket Address Variables

You need to declare a variable of type struct sockaddr_un to use for socket addresses.

For example, the following declarations are used in the example client program:

```
struct sockaddr_un myaddr; /* for local socket address */
struct sockaddr_un peeraddr; /* for peer socket address */
```

Sockaddr_un is a special case of *sockaddr* and is used with the AF_UNIX address domain. The *sockaddr_un* address structure consists of the following fields:

short sun_family	Specifies the address family and should al- ways be set to AF_UNIX
u_char sun_path[108]	Specifies the pathname of the vnode to which the socket is bound or will be bound (e.g. /tmp/mysocket).

The server process only needs an address for its own socket. Your client process will not need an address for its own socket.

Writing the Server Process

This section discusses the calls your server process must make to connect with and serve a client process.

Creating a Socket

The server process must call socket to create a communication endpoint.

Socket and its parameters are described in the following table.

INCLUDE FILES: #include <sys/types.h> #include <sys/socket.h>

SYSTEM CALL:

s = socket(af, type, protocol)
int af, type, protocol;

Parameter	Description of Contents	INPUT Value
af	address family	AF_UNIX
type	socket type	SOCK_STREAM
protocol	underlying protocol to be us	ed 0 (default)
FUNCTION	Ċ	ocket number (HP-UX file lescriptor) -1 if failure occurs
EXAMPLE S	SYSTEM CALL: s	= socket (AF UNIX, SOCK STREAM, 0);

The socket number returned is the socket descriptor for the newly created socket. This number is an HP-UX file descriptor and can be used for reading, writing or any standard file system calls after a BSD IPC connection is established. A socket descriptor is treated like a file descriptor for an open file.

When to Create Sockets

Which Processes	When	
server process	before any other BSD IPC system calls	

Refer to the *socket(2)* entry in the *LAN Reference Pages* for more information on *socket*.

Binding a Socket Address to the Server Process's Socket

After your server process has created a socket, it must call *bind* to bind a socket address. Until an address is bound to the server socket, other processes have no way to reference it.

The server process must bind a specific pathname to this socket, which is used for listening. Otherwise, a client process would not know what pathname to connect to for the desired service.

Set up the address structure with a local address (as described in the "Preparing Address Variables" section) before you make a *bind* call. *Bind* and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>	
	<pre>#include <sys un.h=""></sys></pre>	
	<pre>#include <sys socket.h=""></sys></pre>	
SYSTEM CALL:	bind (s, addr, addrlen)	
	int s;	
	<pre>struct sockaddr_un *addr;</pre>	
	int addrlen;	

Parameter	Description of Contents	INPUT Value
S	socket descriptor of local socket	socket descriptor of socket to be bound
addr	socket address	pointer to address to be bound to s
addrlen	length of socket address	size of struct sockaddr_un
FUNCTION	RESULT:	0 if bind is successful -1 if failure occurs
EXAMPLE S	YSTEM CALL:	struct sockaddr_un myaddr;
		<pre>bind (ls, myaddr, sizeof(struct sockaddr_un));</pre>

When to Bind Socket Addresses

Which Processes

When

server process

after socket is created and before any other BSD IPC system calls

Refer to the *bind(2)* entry in the *LAN Reference Pages* for more information on *bind*.

Setting the Server Up to Wait for Connection Requests

Once your server process has an address bound to it, it must call listen to set up a queue that accepts incoming connection requests. The server process then monitors the queue for requests (using *select(2)* or *accept*, which is described in "Accepting a Connection"). The server process cannot respond to a connection request until it has executed *listen*.

Listen and its parameters are described in the following table.

INCLUDE FILES: none SYSTEM CALL: listen(s, backlog) int s, backlog;

Parameter	Descrip	tion of Contents	INPUT Value	
S	socket descriptor of local socket		server socket's descriptor	
backlog	maximum n connection queue at an	requests in the	size of queue (between 1 and 20)	
FUNCTION	RESULT:	0 if listen is su –1 if failure o		
EXAMPLE SYSTEM CALL:		listen (ls, 5);		

Backlog is the number of unaccepted incoming connections allowed at a given time. Further incoming connection requests are rejected.

When to Set Server Up to Listen

Which Processes	When		
server process	after socket is created and bound and before the server can respond to connection requests		

Refer to the *listen(2)* entry in the LAN Reference Pages for more information on *listen*.

Accepting a Connection

The server process can accept any connection requests that enter its queue after it executes *listen. Accept* creates a new socket for the connection and returns the socket descriptor for the new socket. The new socket:

- is created with the same properties as the old socket;
- has the same bound pathname as the old socket; and
- is connected to the client process' socket.

Accept blocks until there is a connection request from a client process in the queue.

Accept and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys un.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>s = accept(s,addr,addrlen)</pre>
	int s;
	struct sockaddr_un *addr;
	int *addrlen;

Parameter	Description of Co	ntents	INPUT Value	OUTPUT Value
S	socket descriptor of local socket	of	socket descriptor of server socket	unchanged
addr	socket address		pointer to address structure where address will be put	pointer to socket address of client socket that server's new socket is connected to
addrlen	length of address		pointer to the size of struct sockaddr_un	pointer to the actual length of address returned in addr
FUNCTIC	ON RESULT:	succes	t descriptor of new s ssful failure occurs	ocket if accept is
EXAMPL CALL:	E SYSTEM	 addrler	sockaddr_un peeraddr; n = sizeof(sockaddr_un); sept (ls, peeraddr, &addr	len);

There is no way for the server process to indicate which requests it can accept. It must accept all requests or none.

When to Accept a Connection

Which Processes	When	
server process	after executing the listen call	

Refer to the *accept(2)* entry in the *LAN Reference Pages* for more information on *accept*.

Writing the Client Process

This section discusses the calls your client process must make to connect with and be served by a server process.

Creating a Socket

The client process must call socket to create a communication endpoint.

Socket and its parameters are described in the following table.

INCLUDE FIL	ES:	<pre>#include <sys pre="" type<=""></sys></pre>	es.h>
		<pre>#include <sys pre="" sock<=""></sys></pre>	<pre>ket.h></pre>
SYSTEM CAL	L:	s = socket(af, typ int af, type, prot	
Parameter	Descripti	on of Contents	INPUT Value
af	address famil	ly	AF_UNIX
type	socket type		SOCK_STREAM
protocol	underlying pi used	rotocol to be	0 (default)
FUNCTION R	ESULT:	socket number -1 if failure oc	(HP-UX file descriptor) curs
EXAMPLE SY CALL:	STEM	s = socket (AF_UN	IX, SOCK_STREAM, 0);

The socket number returned is the socket descriptor for the newly created socket. This number is an HP-UX file descriptor and can be used for reading, writing or any standard file system calls after a BSD IPC connection is established. A socket descriptor is treated like a file descriptor for an open file.

When to Create Sockets

Which Processe

When

client process

before requesting a connection

Refer to the *socket(2)* entry in the *LAN Reference Pages* for more information on *socket*.

Requesting a Connection

Once the server process is listening for connection requests, the client process can request a connection with the *connect* call.

Connect and its parameters are described in the following table.

INCLUDE FII	LES:	<pre>#include <sys pre="" types.h<=""></sys></pre>	<۱>
		<pre>#include <sys un.h=""></sys></pre>	
		<pre>#include <sys pre="" socket.<=""></sys></pre>	h>
SYSTEM CAL	L:	connect(s, addr, addr	len)
		int s;	
		struct sockaddr_un *a	addr;
		int addrlen;	
Parameter	Descr	iption of Contents	INPUT Value
S	socket de	scriptor of local so	ocket descriptor of socket

S	socket descriptor of local socket	socket descriptor of socket request- ing a connection
addr	pointer to the socket address	pointer to the socket address of the socket to which client wants to connect
addrlen	length of addr	size of address structure pointed to by addr

0 if connect is successful -1 if failure occurs
struct sockaddr_un peeraddr;
<pre>connect (s, peeraddr, sizeof(struct sockaddr_un));</pre>

Connect initiates a connection. When the connection is ready, the client process completes its *connect* call and the server process can complete its *accept* call.

Note

The client process does not get feedback that the server process has completed the *accept* call. As soon as the *connect* call returns, the client process can send data.

When to Request a Connection

Which Processes	When	
client process	after socket is created and after server socket has a listening socket	

Refer to the *connect(2)* entry in the *LAN Reference Pages* for more information on *connect*.

Sending and Receiving Data

After the *connect* and *accept* calls are successfully executed, the connection is established and data can be sent and received between the two socket endpoints. Because the stream socket descriptors correspond to HP-UX file descriptors, you can use the *read* and *write* calls (in addition to *recv* and *send*) to pass data through a socket-terminated channel.

If you are considering the use of the *read* and *write* system calls instead of the *send* and *recv* calls described below, you should consider the following:

Advantage:	If you use <i>read</i> and <i>write</i> instead of <i>send</i> and <i>recv</i> , you can use a socket for <i>stdin</i> or <i>stdout</i> .
Disadvantage:	If you use <i>read</i> and <i>write</i> instead of <i>send</i> and <i>recv</i> , you cannot use the options specified with the <i>send</i> or <i>recv flags</i> parameter.

See the table called "Other System Calls," listed in the "Programming Hints" chapter for more information on which of these system calls are best for your application.

Sending Data

Send and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>count = send(s,msg,len,flags)</pre>
	int s;
	char *msg;
	int len, flags;

Parameter Descript		ion of Contents	INPUT Value	
S	socket descriptor of local sock- et pointer to data buffer size of data buffer		socket descriptor of socket sending data	
msg			pointer to data to be sent	
len			size of msg	
flags	settings for optional flags		0	
FUNCTION 1	RESULT:	number of byte -1 if failure occ		
EXAMPLE S CALL:	YSTEM	count = send (s, b	ouf, 10, 0);	

Send blocks until the specified number of bytes have been queued to be sent, unless you are using nonblocking I/O. (For information on nonblocking I/O, see the "Nonblocking I/O" section of the "Advanced Topics for Stream Sockets" chapter.)

When to Send Data

Which Processes	When		
server or client process	after connection is established		
Defends the set $l(2)$ entry in the	LAND Company for more informed		

Refer to the *send(2)* entry in the *LAN Reference Pages* for more information on *send*.

Receiving Data

Recv and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>count = recv(s,buf,len,flags)</pre>
	int s;
	char *buf;
	int len, flags;
121	

Parameter	Description of Contents socket descriptor of local socket pointer to data buffer maximum number of bytes that should be received		INPUT Value		
s			socket descriptor of socket receiving data		
buf			pointer to buffer that is to receive data		
len			size of data buffer		
flags	settings for optional flags		0		
FUNCTION	RESULT:	number of by -1 if failure of	tes actually received		
EXAMPLE S	YSTEM	<pre>count = recv(s,</pre>	buf, 10, 0);		

Recv blocks until there is at least 1 byte of data to be received, unless you are using nonblocking I/O. (For information on nonblocking I/O, see the "Nonblocking I/O" section of the "Advanced Topics for Stream Sockets" chapter.) The host does not wait for *len* bytes to be available; if less than *len* bytes are available, that number of bytes are received.

No more than *len* bytes of data are received. If there are more than *len* bytes of data on the socket, the remaining bytes are received on the next *recv*.

Flag Options

There are no *flags* options for UNIX Domain (AF_UNIX) sockets. The only supported value for this field is 0.

When to Receive Data

Which Processes	When		
server or client process	after connection is established		

Refer to the recv(2) entry in the LAN Reference Pages for more information on recv.

Closing a Socket

In most applications, you do not have to worry about cleaning up your sockets. When you exit your program and your process terminates, the sockets are closed for you.

If you need to close a socket while your program is still running, use the *close* system call. For example, you may have a daemon process that uses *fork* to create the server process. The daemon process creates the BSD IPC connection and then passes the socket descriptor to the server. You then have more than one process with the same socket descriptor. The daemon process should do a *close* of the socket descriptor to avoid keeping the socket open once the server is through with it. Because the server performs the work, the daemon does not use the socket after the *fork*.

Close decrements the file descriptor count and the calling process can no longer use that file descriptor.

When the last close is executed on a socket descriptor, any unsent data are sent before the socket is closed. Any unreceived data are lost.

Examples Using UNIX Domain Stream Sockets

```
/*
 *
        EXAMPLE PROGRAM
 *
 *
        CATCH - RECEIVE DATA FROM THE PITCHER
 *
 *
        Pitch and catch set up a simple UNIX Domain stream socket
 *
        client-server connection. The client (pitch) then sends data to
 *
        the server (catch), throughput is calculated, and the result is
 *
        printed to the client's stdout.
 */
#include <stdio.h>
#include <time.h>
#include <signal.h>
#include <errno.h>
#include <sys/types.h>
#include <sys/socket.h>
#include <sys/un.h>
#define SOCKNAME "p_n_c"
#define BUFSIZE
                         32*1024-1
char buffer[BUFSIZE]:
struct bullet {
        int bytes;
        int throughput;
        int magic;
} bullet = { 0, 0, 0 }:
send data(fd, buf, buflen)
        char *buf:
{
        int cc;
        while (buflen > 0) {
                cc = send(fd, buf, buflen, 0);
                 if (cc == -1) {
                         perror("send");
                         exit(0);
                 }
                 buf += cc;
                 buflen -= cc;
        }
}
recv data(fd, buf, buflen)
        char *buf;
```

```
{
        int cc;
        while (buflen > 0) {
                cc = recv(fd, buf, buflen, 0);
                 if (cc == -1) {
                        perror("recv");
                         exit(0);
                }
                buf += cc;
                buflen -= cc:
        }
}
main(argc, argv)
        int argc;
        char *argv[];
{
        int bufsize, bytes, cc, i, total, pid, counter_pid;
        float msec;
        struct timeval tp1, tp2;
        int s, ns, recvsize, secs, usec;
        struct timezone tzp;
        struct sockaddr_un sa;
        signal(SIGPIPE, SIG_IGN);
        signal(SIGCLD, SIG IGN);
        setbuf(stdout, 0);
        setbuf(stderr, 0);
        if (argc > 1) {
                 argv++;
                 counter pid = atoi(*argv++);
        } else
                counter pid = 0;
/*
 * Set up the socket variables - address family, socket name.
 * They'll be used later to bind() the name to the server socket.
 */
        sa.sun family = AF UNIX;
        strncpy(sa.sun path, SOCKNAME,
                                  (sizeof(struct sockaddr_un) - sizeof(short)));
 * Create the server socket
```

```
if ((s = socket( AF_UNIX, SOCK_STREAM, 0)) == -1) {
                perror("catch - socket failed"):
                exit(0):
        bufsize = BUFSIZE:
/*
 * Use setsockopt() to change the socket buffer size to improve throughput
 * for large data transfers
 */
        if ((setsockopt(s, SOL SOCKET, SO RCVBUF, &bufsize, sizeof(bufsize)))
                == -1) {
                         perror("catch - setsockopt failed");
                         exit(0);
        }
/*
 * Bind the server socket to its name
 */
        if ((bind(s, &sa, sizeof(struct sockaddr un))) == -1) {
                perror("catch - bind failed");
                exit(0):
        }
 * Call listen() to enable reception of connection requests
 * (listen() will silently change the given backlog, 0, to be 1 instead)
 */
        if ((listen(s, 0)) == -1) {
                perror("catch - listen failed");
                exit(0);
        }
next conn:
        i = sizeof(struct sockaddr un):
/*
 * Call accept() to accept the connection request. This call will block
 * until a connection request arrives.
 */
        if ((ns = accept(s, &sa, &i)) == -1) {
                if (errno == EINTR)
                         goto next conn;
                perror("catch - accept failed");
                exit(0):
        if ((pid = fork()) != 0) {
                close(ns);
                goto next conn;
        }
        close(s);
 * Receive the bullet to synchronize with the other side
 */
        recv data(ns, &bullet, sizeof(struct bullet));
```

11-64 BSD IPC Using UNIX Domain Stream Sockets

```
if (bullet.magic != 12345) {
                 printf("catch: bad magic %d\n", bullet.magic);
                 exit(0);
        }
        bytes = bullet.bytes;
        recvsize = (bytes>BUFSIZE)?BUFSIZE:bytes;
1*
 * Send the bullet back to complete synchronization
 */
        send_data(ns, &bullet, sizeof(struct bullet));
        cc = 0;
        if (counter_pid)
                 kill(counter_pid, SIGUSR1);
        if (gettimeofday(&tp1, &tzp) == -1) {
                 perror("catch time of day failed");
                 exit(0);
        }
   Receive data from the client
 */
        total = 0;
        i = bytes;
        while (i > 0) {
                 cc = recvsize < i ? recvsize : i;</pre>
                 recv_data(ns, buffer, cc);
                 total += cc:
                 i -= cc;
        }
1*
 * Calculate throughput
 */
        if (gettimeofday(&tp2, &tzp) == -1) {
                 perror("catch time of day failed");
                 exit(0);
        }
        if (counter_pid)
                 kill(counter_pid, SIGUSR2);
        secs = tp2.tv_sec - tp1.tv_sec;
        usec = tp2.tv_usec - tp1.tv usec;
        if (usec < 0) {
                secs--;
                 usec += 1000000;
        }
        msec = 1000*(float)secs;
        msec += (float)usec/1000;
        bullet.throughput = bytes/msec;
```

```
EXAMPLE CLIENT PROGRAM
        PITCH - SEND DATA TO THE CATCHER
        Pitch and catch set up a simple UNIX Domain stream socket
 *
        client-server connection. The client (pitch) then sends data to
 *
        the server (catch), throughput is calculated, and the result is
 *
        printed to the client's stdout.
 */
#include <stdio.h>
#include <time.h>
#include <netdb.h>
#include <signal.h>
#include <sys/types.h>
#include <sys/socket.h>
#include <sys/un.h>
#define SOCKNAME "p_n_c"
#define BUFSIZE
                         32*1024-1
char buffer[BUFSIZE];
struct bullet {
        int bytes;
        int throughput;
        int magic;
} bullet = { 0, 0, 12345 };
send data(fd, buf, buflen)
        char *buf;
{
        int cc;
        while (buflen > 0) {
                cc = send(fd, buf, buflen, 0);
                 if (cc == -1) {
                         perror("send");
                         exit(0);
                 }
                buf += cc:
                buflen -= cc:
        }
}
recv data(fd, buf, buflen)
        char *buf;
```

```
{
        int cc;
        while (buflen > 0) {
                cc = recv(fd, buf, buflen, 0);
                 if (cc == -1) {
                         perror("recv");
                         exit(0);
                 }
                 buf += cc;
                 buflen -= cc:
        }
}
main( argc, argv)
        int argc;
        char *argv[]:
{
        int bufsize, bytes, cc, i, total, pid;
        float msec:
        struct timeval tpl, tp2;
        int s, sendsize, secs, usec;
        struct timezone tzp;
        struct sockaddr_un sa;
        signal(SIGPIPE, SIG IGN);
        setbuf(stdout, 0);
        setbuf(stderr, 0);
        if (argc < 2) {
                printf("usage: pitch Kbytes [pid]\n");
                 exit(0);
        }
        argv++;
/*
 * Set up the socket variables (address family; name of server socket)
 * (they'll be used later for the connect() call)
 */
        sa.sun_family = AF_UNIX;
        strncpy(sa.sun path, SOCKNAME,
                                 (sizeof(struct sockaddr_un) - sizeof(short)));
        bullet.bytes = bytes = 1024*atoi(*argv++);
        if (argc > 2)
                pid = atoi(*argv++);
        else
                 pid = 0;
        sendsize = (bytes < BUFSIZE) ? bytes : BUFSIZE;</pre>
```

11-68 BSD IPC Using UNIX Domain Stream Sockets

```
/*
   Create the client socket
 */
        if ((s = socket( AF UNIX, SOCK STREAM, 0)) == -1) {
                perror("pitch - socket failed");
                exit(0);
        }
        bufsize = BUFSIZE;
/*
   Change the default buffer size to improve throughput for
   large data transfers
 */
        if ((setsockopt(s, SOL_SOCKET, SO_SNDBUF, &bufsize, sizeof(bufsize)))
                == -1) {
                         perror("pitch - setsockopt failed");
                         exit(0);
        }
/*
 * Connect to the server
 */
        if ((connect(s, &sa, sizeof(struct sockaddr un))) == -1) {
                perror("pitch - connect failed");
                exit(0);
        }
/*
 * send and receive the bullet to synchronize both sides
 */
        send data(s, &bullet, sizeof(struct bullet));
        recv data(s, &bullet, sizeof(struct bullet));
        cc = 0;
        if (pid)
                 kill(pid,SIGUSR1);
        if (gettimeofday(&tpl, &tzp) == -1) {
                perror("pitch time of day failed");
                exit(0);
        }
        i = bytes:
        total = 0;
/*
 * Send the data
 */
        while (i > 0) {
                cc = sendsize < i ? sendsize : i;
                send data(s, buffer, cc);
                 i -= cc:
                total += cc;
        }
```

```
/*
 * Receive the bullet to calculate throughput
 */
         recv data(s, &bullet, sizeof(struct bullet));
         if (gettimeofday(&tp2, &tzp) == -1) {
                  perror("pitch time of day failed");
                  exit(0);
         }
         if (pid)
                  kill(pid, SIGUSR2);
/*
 * Close the socket
 */
         close(s);
         secs = tp2.tv sec - tp1.tv sec;
         usec = tp2.tv_usec - tp1.tv_usec;
         if (usec < 0) {
                  secs--;
                  usec += 1000000;
         }
         msec = 1000*(float)secs;
         msec += (float)usec/1000;
         printf("PITCH: %d Kbytes/sec\n", (int)(bytes/msec));
printf("CATCH: %d Kbytes/sec\n", bullet.throughput);
         printf("AVG: %d Kbytes/sec\n", ((int)(bytes/msec)+bullet.throughput)/2);
}
```

Advanced Topics for Stream Sockets

Socket Options

The operation of sockets is controlled by socket level options. The following options are supported for Internet stream sockets:

- SO_REUSEADDR
- SO_KEEPALIVE
- SO_DONTROUTE
- SO_SNDBUF
- SO_RCVBUF
- SO_LINGER
- SO_DONTLINGER

The following options are supported for UNIX Domain stream sockets:

- SO_SNDBUF
- SO_RCVBUF

In addition, the SO_DEBUG option is supported for compatibility only; it has no functionality.

The next section discusses how to set socket options and get the current value of a socket option. Following those discussions is a description of each available option.

Getting and Setting Socket Options

The socket options are defined in the $\langle sys / socket.h \rangle$ file. You can get the current status of an option with the *getsockopt* call, and you can set the value of an option with the setsockopt call.

Setsockopt and its parameters are described in the following table:

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>setsockopt(s, level, optname, optval, optlen)</pre>
	int s, level, optname;
	char *optval;
	int optlen:

Parameter	Parameter Description		INPUT Value
s	socket descriptor		socket descriptor for which options are to be set
level	protocol lev	vel	SOL_SOCKET
optname	name of option		supported option name
optval	pointer to option input value		0, or if optname = SO_LINGER, a pointer to the linger interval value of SO_LINGER, or if optname = SO_SNDBUF or SO_RCVBUF, a pointer to an integer containing the new buffer size value.
optlen	length of op	otval	0 or size of optval
FUNCTION RESULT:		0 if setsockopt –1 if failure oc	
EXAMPLE SYSTEM CALL:		See the descrip	ption of the DDR option for an example.

For options that do not take an input value, *optval* and *optlen* should both be 0.

Refer to the getsockopt(2) entry in the ARPA/Berkeley Services Reference Pages for more information on setsockopt.

Getsockopt and its parameters are described in the following table:

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>		
	<pre>#include <sys socket.h=""></sys></pre>		
SYSTEM CALL:	<pre>getsockopt(s, level, optname, optval, optlen) int s, level, optname; char *optval; int *optlen;</pre>		

Parameter	Description of Contents socket descriptor		INPUT Value OUTPUT Value		
\$			socket descriptor for which option values are to be returned	unchanged	
level	protocol level		SOL_SOCKET	unchanged	
optname	name of option		supported option name	unchanged	
optval	pointer to current value of option		pointer to buffer where option's cur- rent value is to be returned	pointer to buffer that contains current op- tion value	
optlen	pointer to length of optval		pointer to maximum number of bytes to be returned by optval	pointer to actual size of optval returned	
FUNCTIO	ON RESULT:	0 if th	ne option is set		
			th errno $=$ ENOPRO cified option is not set		
			th errno = some oth ure occurs	er value	
EXAMPLE SYSTEM CALL:		getsoc	getsockopt(s, SOL_SOCKET, SO_REUSEADDR, 0, 0);		

For options that do not take an input value, *optval* and *optlen* should both be 0.

Refer to the getsockopt(2) entry in the ARPA/Berkeley Services Reference Pages section for more information on getsockopt.

SO_REUSEADDR

Note that this option is not supported for UNIX Domain sockets.

SO_REUSEADDR enables you to restart a daemon which was killed or terminated.

This option modifies the rules used by *bind* to validate local addresses, but it does not violate the uniqueness requirements of an association. SO_REUSEADDR modifies the bind rules only when a wildcard IP address is used in combination with a particular protocol port. The host still checks at connection time to be sure any other sockets with the same local address and local port do not have the same remote address and remote port. *Connect* fails if the uniqueness requirement is violated.

The following example shows the SO_REUSEADDR option's use:

Suppose that a network daemon server is listening on a specific port: port 2000. If you executed netstat -an, part of the output would resemble:

Active of	connectio	ons (including servers)		
Proto Re	ecv-Q Se	nd-Q	Local Address	Foreign Address	(state)
tcp	0	0	*.2000	*.*	LISTEN

Network Daemon Server Listening at Port 2000

When the network daemon accepts a connection request, the accepted socket will bind to port 2000 and to the Internet Protocol address where the daemon is running (e.g. 192.6.250.100).

If you then executed netstat -an, the output would resemble:

Active (connectio	ons (including servers)		
Proto Re	ecv-Q Ser	id-Q	Local Address	Foreign Address	(state)
tcp	0	0	192.6.250.100.2000	192.6.250.101.4000	ESTABLISHED
tcp	0	0	*.2000	*.*	LISTEN

New Connection Established, Daemon Server Still Listening

Here the network daemon has established a connection to the client (192.6.250.101.4000) with a new server socket. The original network daemon server continues to listen for more connection requests.

If the listening network daemon process is killed, attempts to restart the daemon fail if SO_REUSEADDR is not set. The restart fails because the daemon attempts to bind to port 2000 and a wildcard Internet Protocol address (e.g. *.2000). The wildcard Internet Protocol address matches the Internet Protocol address of the established connection (192.6.250.100), so the bind aborts to avoid duplicate socket naming.

When SO_REUSEADDR is set, *bind* ignores the wildcard match, so the network daemon can be restarted.

SO_REUSEADDR cannot be cleared once you set it.

An example usage of this option is:

```
setsockopt (s, SOL_SOCKET, SO_REUSEADDR, (char *)0, 0);
bind (s, &sin, sizeof(sin));
```

SO_KEEPALIVE

Note that this option is not supported for UNIX Domain sockets.

This option enables the periodic transmission of messages on a connected socket. This occurs at the transport level and does not require any work in your application programs.

If the peer socket does not respond to these messages, the connection is considered broken. The next time one of your processes attempts to use a connection that is considered broken, the process is notified (with a SIG-PIPE signal if you are trying to send, or an end-of-file condition if you are trying to receive) that the connection is broken.

SO_KEEPALIVE cannot be cleared once you set it.

SO_DONTROUTE

Note that this option is not supported for UNIX Domain sockets.

SO_DONTROUTE indicates that outgoing messages should bypass the standard routing facilities. Instead, messages are directed to the appropriate network interface according to the network portion of the destination address.

SO_SNDBUF

SO_SNDBUF changes the send socket buffer size. Increasing the send socket buffer size allows a user to send more data before the user's application will block, waiting for more buffer space.

Note

Increasing buffer size to send larger portions of data before the application blocks **may** increase throughput, but the best method of tuning performance is to experiment with various buffer sizes.

You can increase a stream socket's buffer size at any time but decrease it only prior to establishing a connection.

The maximum buffer size for stream sockets is 65535 bytes.

Example:

```
int result;
int buffsize = 10,000;
result = setsockopt(s, SOL_SOCKET, SO_SNDBUF, &buffsize, sizeof(buffsize));
```

SO_RCVBUF

SO_RCVBUF changes the receive socket buffer size.

You can increase a stream socket's buffer size at any time but decrease it only prior to establishing a connection.

The maximum buffer size for stream sockets is 65535 bytes.

Example:

```
int result;
int buffsize = 10,000;
result = setsockopt(s, SOL SOCKET, SO RCVBUF, &buffsize, sizeof(buffsize));
```

Summary Information for Changing Socket Buffer Size

Socket Type (Protocol)	When Buffer Size Increase Allowed	When Buffer Size Decrease Allowed	Maximum Buffer Size
stream (TCP)	at any time	only prior to estab- lishing a connection	65535 bytes

SO_LINGER

Note that this option is not supported for UNIX Domain sockets.

SO_LINGER controls the actions taken when a *close* is executed on a socket that has unsent data.

This option can be cleared by setting SO_DONTLINGER. The default is SO_DONTLINGER.

The linger timeout interval is set with a parameter in the *setsockopt* call. The only useful values are zero and nonzero:

- If SO_LINGER is set with a nonzero timeout interval, the host blocks the *close* call until it is able to transmit the remaining data or until the protocol itself (TCP) expires. This is called a graceful disconnect.
- If SO_LINGER is set with a zero timeout interval, *close* is not blocked even if queued data exist. This is called a hard close, because it closes the socket immediately, whether data need to be sent or not. All unsent data are immediately lost.

Example:

```
int result;
int linger = 1;
result = setsockopt(s, SOL_SOCKET, SO_LINGER, &linger, sizeof(linger));
```

SO_DONTLINGER

This option is the default. It can be overridden by setting SO_LINGER.

SO_DONTLINGER controls the actions taken when a *close* is executed on a socket. If SO_DONTLINGER is set on a stream socket with unsent data, the host allows the close call to return immediately, but it tells TCP to wait. Queued data are sent if possible, until TCP times out. This is also called a graceful disconnect.

Socket Option	Linger Interval	Graceful Close	Hard Close	Waits for Close	Does Not Wait for Close
SO_DONTLINGER SO_LINGER	don't care zero	х	х		X X
SO_LINGER	nonzero	X		X	

Summary of Linger Options on Close

Synchronous I/O Multiplexing with Select

The *select* system call can be used with sockets to provide a synchronous multiplexing mechanism. The system call has several parameters which govern its behavior. If you specify a zero pointer for the **timout** parameter, *select* will block until one or more of the specified socket descriptors are ready. If timeout is a non-zero pointer, it specifies a maximum interval to wait for the selection to complete.

A select of a socket descriptor for reading is useful on:

- a connected socket, because it determines when data has arrived and is ready to be read without blocking; use the FIONREAD parameter to the *ioctl* system call to determine exactly how much data is available.
- a listening socket, because it determines when you can accept a connection without blocking.

A select of a socket descriptor for writing is useful on:

- a connecting socket, because it determines when the connection is complete.
- a connected socket, because it determines when more data can be sent without blocking. This implies that at least one byte can be sent; there is no way, however, to determine exactly how many bytes can be sent.

Selecting for exceptional conditions is currently meaningless for Berkeley sockets. *Select* will always return true for sockets that are no longer capable of being used (e.g. if a *close* or *shutdown* system call has been executed against them).

Select is used in the same way as in other applications. Refer to the select(2) entry in the HP-UX Reference manual for information on how to use select. For an asynchronous alternative to select, see the next section, "Sending and Receiving Data Asynchronously."

Example:

The following example illustrates the select system call. Since it is possible for a process to have more than 32 open file descriptors, the bit masks used by *select* are interpreted as arrays of intergers. The following useful macros can be used to manipulate bit masks of this form.

```
#define BPI 32
                        /* bits per int */
#define FD ZERO(p)
                     bzero((char *) (p), sizeof(*(p)))
#define FD SET(n, p) ((p)-fdm bits[(n)/BPI] = (1 < ((n) % BPI)))</pre>
#define FD CLR(n, p) ((p)-fdm bits[(n)/BPI] &= ~(1 < ((n) % BPI)))</pre>
#define FD ISSET(n, p) ((p)-fdm bits[(n)/BPI] & (1 < ((n) % BPI)))</pre>
struct fd mask {
        u long fdm bits[NOFILE/BPI+1] /* NOFILE max # of fd's per process */
};
do select(s)
                        /* socket to select on. initialized */
int s;
ł
        struct fd set read mask, write mask;
                                                /* bit masks */
        int nfds:
                                         /* number to select on */
        int nfd:
                                         /* number found */
        for (;;) {
                                         /* for example... */
                FD ZERO(&read_mask);
                                        /* select will overwrite on return */
                FD ZERO(&write mask);
                                        /* we care only about the socket */
                FD SET(s, &read mask);
                FD SET(s, &write mask);
                nfds = s;
                                         /* select descriptors 0 through s */
                nfd = select(nfds, &read_mask, &write_mask, (int *) 0,
                                   (struct timeval *) 0); /* will block*/
                if (nfd == -1) {
                        perror( "select: unexpected condition" );
                        exit(1);
                if (FD ISSET(s, &read mask))
                        do read(s);
                                        /* something to read on socket s */
                                         /* fall through as maybe more to do */
                if (FD ISSET(s, &write mask))
                        do write(s);
                                       /* space to write on socket s */
        }
}
```

11-80 Advanced Topics for Stream Sockets

Sending and Receiving Data Asynchronously

Asynchronous sockets allow a user program to receive a SIGIO signal when the socket's state changes. This state change can occur, for example, when new data arrives. Currently the user must issue a *select* system call to determine if data are available. If other processing is required of the user program, the need to call *select* can complicate an application by forcing the user to implement some form of polling, whereby all sockets are checked periodically. Asynchronous sockets would allow the user to separate socket processing from other processing, eliminating polling altogether. *Select* may still be required to determine exactly why the signal is being delivered or to which socket the signal applies.

Generation of the SIGIO signal is protocol dependent. It mimics the semantics of *select* in the sense that the signal is generated whenever *select* would return true. It is generally accepted that connectionless protocols deliver the signal whenever a new packet arrives. For connection oriented protocols, the signal is also delivered when connections are established or broken, as well as when additional outgoing buffer space becomes available. Be aware that these assertions are guidelines only; any signal handler should be robust enough to handle signals in unexpected situations.

The delivery of SIGIO signal is dependent upon two things. First, the socket state must be set as asynchronous; this is done using the FIOASYNC flag of the *ioctl* system call. Secondly, the process group (pgrp) associated with the socket must be set; this is done using the SIOCSPGRP flag of *ioctl*. The sign value of the pgrp can lead to various signals being delivered. Specifically, if the pgrp is negative, this implies that a signal should be delivered to the process whose PID is the absolute value of the pgrp. If the pgrp is positive, a signal should be delivered to the process group identified by the absolute value of the pgrp.

Any application that chooses to use asynchronous sockets must explicitly activate the described mechanism. The SIGIO signal is a "safe" signal in the sense that if a process is unprepared to handle it, the default action is to ignore it. Thus, any existing applications are immune to spurious signal delivery. Notification that out of band data has been received is also done asynchronously; for more details, see the section in this chapter, "Sending and Receiving Out of Band Data." **Example:**

The following example sets up a listen SOCK_STREAM socket as asynchronous. This is typical of an application that needs to be notified when connection requests arrive.

```
/* listen SOCK_STREAM socket */
int ls:
                                   /* for ioctl, to turn on async */
         int flag = 1;
         int iohndlr();
                                   /* the function which handles the SIGIO */
         signal( SIGIO, iohndlr ); /* set up the handler */
         if( ioctl( ls, FIOASYNC, &flag ) == -1) {
                  perror( "can't set async on socket" );
             exit(1);
         }
                                  /* process group negative == deliver to process */
         flag = -getpid();
         if( ioctl( ls, SIOCSPGRP, &flag ) == -1 ) {
                  perror( "can't get pgrp" );
           exit(1);
         }
                    /* signal can come any time now */
```

The following example illustrates the use of process group notification. Note that the real utility of this feature is to allow multiple processes to receive the signal, which is not illustrated here. For example, the socket type could be SOCK_DGRAM; a signal here can be interpreted as the arrival of a service-request packet. Multiple identical servers could be set up, and the first available one could receive and process the packet.

```
/* ioctl to turn on async */
int flag = 1;
    int iohndlr();
    signal( SIGIO, iohndlr );
                                          /* set my processes' process group */
    setpgrp();
    if( ioctl( s, FIOASYNC, &flag ) == -1) {
              perror( "can't set async on socket" );
        exit(1);
}
    flag = getpid();
                                          /* process group + == deliver to every
                                             process in group */
    if( ioctl( s, SIOCSPGRP, &flag ) == -1 ) {
              perror( "can't set pgrp" );
           exit(1);
    }
                                 /* signal can come any time now */
```

Nonblocking I/O

Sockets are created in blocking mode I/O by default. You can specify that a socket be put in nonblocking mode by using the *ioctl* system call with the FIOSNBIO request.

An example usage of this call is:

```
#include <sys/ioctl.h>
...
ioctl(s, FIOSNBIO, &arg);
```

Arg is a pointer to int:

- When int equals 0, the socket is changed to blocking mode.
- When int equals 1, the socket is changed to nonblocking mode.

If a socket is in nonblocking mode, the following calls are affected:

accept	If no connection requests are present, <i>accept</i> returns immediately with the EWOULDBLOCK error.
connect	If the connection cannot be completed immediately, <i>connect</i> returns with the EINPROGRESS error.
recv	If no data are available to be received, recv returns the value -1 and the EWOULDBLOCK error. This is also true for read.
send	If there is no available buffer space for the data to be transmitted, <i>send</i> returns the value -1 and the EWOULDBLOCK error. This is also true for <i>write</i> .

The O_NDELAY flag for *fcntl(2)* is also supported. If you use this flag and there are no data available to be received on a *recv*, *recvfrom*, or *read* call, the call returns immediately with the value of 0. This is the same as returning an end-of-file condition. This is also true for *send*, *sendto* and *write* if there is not enough buffer space to complete the send.

Note

The O_NDELAY flag has precedence over the FIOSNBIO flag.

Using Shutdown

When your program is done reading or writing on a particular socket connection, you can use *shutdown* to bring down a part of the connection. (See the example programs for stream sockets.)

When one process uses *shutdown* on a socket descriptor, all other processes with the same socket descriptor are affected. *Shutdown* causes all or part of a full-duplex connection on the specified socket to be disabled. When *shutdown* is executed, the specified socket is marked unable to send or receive, according to the value of *how*:

- If *how* = 0, the specified socket can no longer receive data. The connection is not completely down until both sides have done a *shutdown* or a *close*.
- If *how* = 1, *shutdown* starts a graceful disconnect by attempting to send any unsent data before blocking further sending. *Shutdown* sends an end-of-file condition to the peer, indicating that there are no more data to be sent.

Once both shutdown(s, 0) and shutdown(s, 1) have been executed on the same socket descriptor, the only valid operation on the socket at this point is a *close*. • If *how* = 2, the specified socket can no longer send or receive data. The only valid operation on the socket is a *close*. This has the same effect as executing shutdown(s,0) and shutdown(s,1) on the same socket descriptor.

If you use *close* on a socket, *close* pays attention to the SO_LINGER option, but *shutdown(s, 2)* does not. With *close*, the socket descriptor is deal-located and the last process using the socket destroys it.

Shutdown and its parameters are described in the following table.

INCLUDE FILES: none

SYSTEM CALL:

shutdown(s,how)
int s, how;

Parameter	Description of Contents	INPUT Value
s	socket descriptor	socket descriptor of socket to be shut down
how	number that indicates the type of shutdown	0, 1 or 2

FUNCTION RESULT:	0 if shutdown is successful -1 if failure occurs
EXAMPLE SYSTEM CALL:	shutdown (s, 1);

When to Shut Down a Socket

Which Processes	When	
server or client process	(optionally) after the process has sent all messages and wants to indi- cate that it is done sending.	

Refer to the *shutdown(2)* entry in the *ARPA/Berkeley Services Reference Pages* section for more information on *shutdown*.

Using Read and Write to Make Stream Sockets Transparent

An example application of *read* and *write* with stream sockets is to fork a command with a socket descriptor as *stdout*. The peer process can *read* input from the command. The command can be any command and does not have to know that *stdout* is a socket. It might use *printf*, which results in the use of *write*. Thus, the stream sockets are transparent.

Sending and Receiving Out of Band Data

Note that this option is not supported for UNIX Domain (AF_UNIX) sockets.

If an abnormal condition occurs when a process is in the middle of sending a long stream of data, it is useful to be able to alert the other process with an urgent message. The TCP stream socket implementation includes an out of band data facility. Out of band data uses a **logically** independent transmission channel associated with a pair of connected stream sockets. TCP supports the reliable delivery of only one out of band message at a time. The message can be a maximum of one byte long.

Out of band data arrives at the destination node in sequence and in stream, but is delivered independently of normal data; the out of band data receiver is notified with the SIGURG signal. The receiving process can read the out of band message and take the appropriate action based on the message contents. A logical mark is placed in the normal data stream to indicate the point at which the out of band data was sent, so that data before the message can be handled differently (if necessary) from data following the message.



Data Stream with Out of Band Marker

For a program to know when out of band data is available to be received, you may arrange the program to catch the SIGURG signal as follows:

11-86 Advanced Topics for Stream Sockets

```
struct sigvec vec;
int onurg();
int pid, s;
/*
 ** arrange for onurg() to be called when SIGURG is received:
 */
vec.sv_handler = onurg;
vec.sv_mask = 0;
vec.sv_onstack = 0;
if (sigvector(SIGURG, &vec, (struct sigvec *) 0) < 0) {
    perror("sigvector(SIGURG)");
}
```

Onurg() is a routine that handles out of band data in the client program.

In addition, the socket's process group must be set, as shown below. The kernel will not send the signal to the process (or process group) unless this is done, even though the signal handler has been enabled.

```
/*
** arrange for the current process to receive SIGURG
** when the socket s has urgent data:
*/
pid = -getpid();
if (ioctl(s, SIOCSPGRP, (char *) &pid) < 0) {
    perror("ioctl(SIOCSPGRP)");
}</pre>
```

Refer to the *socket(7)* entry in the *ARPA/Berkeley Services Reference Pages* for more details.

If the server process is sending data to the client process, and a problem occurs, the server can send an out of band data byte by executing a *send* with the MSG_OOB flag set. This sends the out of band data and a SIGURG signal to the receiving process.

send(sd, &msg, 1, MSG_00B)

When a SIGURG signal is received, *onurg* is called. *Onurg* receives the out of band data byte with the MSG_OOB flag set on a *recv* call.

It is possible that the out of band byte has not arrived when the SIGURG signal arrives. *recv* never blocks on a receive of out of band data, so the client may need to repeat the *recv* call until the out of band byte arrives. *Recv* will return EINVAL if the out of band data is not available.

The out of band data byte is stored independently from the normal data stream. You cannot read **past** the out of band pointer location in one *recv* call. If you request more data than the amount queued on the socket before the out of band pointer, then *recv* will return only the data up to the out of band pointer. However, once you read past the out of band pointer location with subsequent *recv* calls, the out of band byte can no longer be read.

Usually the out of band data message indicates that all data currently in the stream can be flushed. This involves moving the stream pointer with successive *recv* calls, to the location of the out of band data pointer.

The *ioctl* request SIOCATMARK informs you, as you receive data from the stream, when the stream pointer has reached the out of band pointer. If *ioctl* returns a 0, the next *recv* provides data sent by the server prior to transmission of the out of band data. *Ioctl* returns a 1 when the stream pointer reaches the out of band byte pointer. The next *recv* provides data sent by the server after the out of band message.

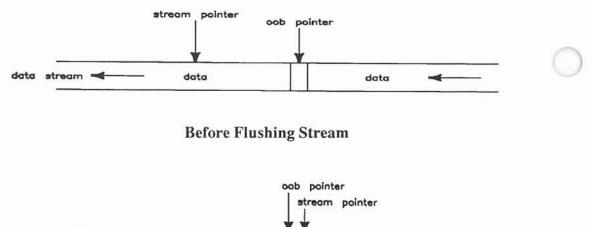
The following code segment illustrates how the SIOCATMARK request can be used in a SIGURG interrupt handler. The example also shows a buffer being flushed.

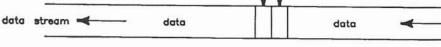
```
/* s is the socket with urgent data */
```

```
onurg()
{
    int atmark;
    char mark;
    char flush [100];
    while (1) {
        /*
        ** check whether we have read the stream
        ** up to the OOB mark yet
        */
```

```
if (ioctl(s, SIOCATMARK, &atmark) < 0) {</pre>
        /* if the ioctl failed */
        perror("ioct1(SIOCATMARK)");
        return;
    }
    if (atmark) {
        /* we have read the stream up to the OOB mark */
        break;
    }
    1*
    ** read the stream data preceding the mark,
    ** only to throw it away
    */
    if (read(s, flush, sizeof(flush)) <= 0) {</pre>
        /* if the read failed */
        return;
    }
}
1*
** receive the OOB byte
*/
recv(s, &mark, 1, MSG_00B);
printf("received %c 00B\n", mark);
return;
```

}







Note

This completes the discussion of stream sockets. If you do not plan to use datagram sockets, skip to the "Programming Hints" section.

IPC Using Internet Datagram Sockets

As discussed in the "Protocols" section, Internet UDP datagram sockets provide bidirectional flow of data with record boundaries preserved. However, there is no guarantee that messages are reliably delivered. If a message is delivered, there is no guarantee that it is in sequence and unduplicated, but the data in the message are guaranteed to be intact. Datagram sockets are not supported for UNIX sockets.

Datagram sockets allow you to send and receive messages without establishing a connection. Each message includes a destination address. Processes involved in data transfer are not required to have a server-client relationship; the processes can be symmetrical.

Unlike stream sockets, datagram sockets allow you to send to many destinations from one socket, and receive from many sources with one socket. There is no two-process model, although a two-process model is the simplest case of a more general multiprocess model. The terms server and client are used in this section only in the application sense. There is no difference in the calls that must be made by the processes involved in the data transfer.

For example, you might have a name server process that receives host names from clients all over a network. That server process can send host name and internet address combinations back to the clients. This can all be done with one UDP socket.

The simplest two-process case is used in this chapter to describe IPC using datagram sockets.

The following table lists the steps required to exchange data between datagram sockets. Each step is described in more detail in the sections that follow the table.

Client Process Activity	System Call Used	Server Process Activity	System Call Used
create a socket	socket()	create a socket	socket()
bind a socket address	bind()	bind a socket address	s <i>bind()</i>
send message	sendto()	receive message send message	recvfrom() sendto()
receive message	recvfrom()		

Setting Up for Data Transfer Using Datagram Sockets

The following sections discuss each of the activities mentioned in the previous table. The description of each activity specifies a system call and includes:

- what happens when the system call is used;
- when to make the system call;
- what the parameters do;
- how the call interacts with other IPC system calls; and
- where to find details on the system call.

The datagram socket program examples are at the end of these descriptive sections. You can refer to them as you work through the descriptions.

Preparing Address Variables

Before your client process can make a request of the server process, you must establish the correct variables and collect the information that you need about the server process and the service provided.

The server process needs to:

- declare socket addresss variables;
- · assign a wildcard address; and
- get the port address of the service that you want to provide.

The client process needs to:

- declare socket address variables;
- get the remote server's internet address; and
- get the port address for the service that you want to use.

These activities are described next. In addition, refer to the program example at the end of the "IPC Using Datagram Sockets" section to see how these activities work together.

Declaring Socket Address Variables

You need to declare a variable of type struct *sockaddr_in* to use for the local socket address for both processes.

For example, the following declarations are used in the example client program:

```
struct sockaddr_in myaddr; /* for local socket address */
struct sockaddr_in servaddr; /* for server socket address */
```

Sockaddr_in is a special case of sockaddr and is used with the AF_INET addressing domain. Both types are shown in this chapter, but sockaddr_in makes it easier to manipulate the internet and port addresses. Some of the IPC system calls are declared using a pointer to sockaddr, but it can also be a pointer to sockaddr_in.

The sockaddr_in address structure consists of the following fields:

short sin_family	Specifies the address family and should al- ways be set to AF_INET.
u_short <i>sin_port</i>	Specifies the port address. Assign this field when you bind the port address for the socket or when you get a port address for a specific service.
struct in_addr sin_addr	Specifies the internet address. Assign this field when you get the internet address for the remote host.

The server process must bind the port address of the service to its own socket and establish an address structure to store the clients' addresses when they are received with *recvfrom*.

The client process does not have to bind a port address for its local socket; the host binds one automatically if one is not already bound.

Refer to the *inet*(7F) entry in the ARPA/Berkeley Services Reference Pages for more information on sockaddr_in.

11-94 IPC Using Internet Datagram Sockets

Getting the Remote Host's Network Address

The client process can use *gethostbyname* to obtain the internet address of the host and the length of that address (as the size of struct *inaddr*) from */etc/hosts*.

Gethostbyname and its parameters are described in the following table.

INCLUDE FILES: #include <netdb.h> SYSTEM CALL: struct hostent *gethostbyname(name) char *name;

Parameter	Descrip	tion of Contents	INPUT Value
name	pointer to a (null-termin	valid node name nated string)	host name
FUNCTION	RESULT:	address	r (0) if failure occurs
EXAMPLE S	YSTEM		

EXAMPLE SYSTEM CALL:

#include <netdb.h>
struct hostent *hp; /* point to host info for name server host */
...
servaddr.sin_family = AF_INET;
hp = gethostbyname (argv[1]);
servaddr.sin_addr.s_addr = ((struct in_addr *)(hp->h_addr))->s_addr;

The *argv[1]* parameter is the host name specified in the client program command line.

Refer to the gethostent(3N) entry in the ARPA/Berkeley Services Reference Pages for more information on gethostbyname.

Getting the Port Address for the Desired Service

When a client process needs to use a service that is offered by some server process, it must send a message to the server's socket. The client process must know the port address for that socket. If the service is not in *letc/services*, you must add it.

Getservbyname obtains the port address of the specified service from *letc/services*.

Getservbyname and its parameters are described in the following table.

INCLUDE FIL	LES:	#include <netdb.h></netdb.h>	ě
SYSTEM CAL	L:	struct servent *ge char *name, *proto	etservbyname(name, proto) ;;
Parameter	Descript	ion of Contents	INPUT Value
name	pointer to a	valid service name	service name
proto	pointer to th used	ne protocol to be	"udp" or 0 if UDP is the only protocol for the service
FUNCTION R	ESULT:	dress	ct servent containing port ad- (0) if failure occurs
EXAMPLE SY CALL:	STEM		o; /* pointer to service info */ e ("example", "udp");

When to Get Server's Socket Address

Which Processes	When
server process	before binding
client process	before client requests the service from the host

Refer to the getservent(3N) entry in the ARPA/Berkeley Services Reference Pages for more information on getservbyname.

Using a Wildcard Local Address

Wildcard addressing simplifies local address binding. When an address is assigned the value of INADDR_ANY, the host interprets the address as any valid address.

This means that the server process can receive on a wildcard address and does not have to look up its own internet address. For example, to bind a specific port address to a socket, but leave the local internet address unspecified, the following source code could be used:

```
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
...
struct sockaddr_in sin;
...
s = socket(AF_INET, SOCK_DGRAM, 0);
sin.sin_family = AF_INET;
sin.sin_addr.s_addr = INADDR_ANY;
sin.sin_port = MYPORT;
bind (s, &sin, sizeof(sin));
```

Writing the Server and Client Processes

This section discusses the calls your server and client processes must make.

Creating Sockets

Both processes must call socket to create communication endpoints.

Socket and its parameters are described in the following table.

INCLUDE FIL	LES:	<pre>#include <sys #include="" <sys="" pre="" sock<="" type=""></sys></pre>	
SYSTEM CAL	L:	s = socket(af, typ int s, af, type, p	
Parameter	Descripti	on of Contents	INPUT Value
af	address fami	ly	AF_INET
type	socket type		SOCK_DGRAM
protocol	underlying p	rotocol to be used	0 (default) or value returned by getprotobyname
FUNCTION R	ESULT:	socket number –1 if failure occ	(HP-UX file descriptor)
EXAMPLE SY CALL:	STEM	ls = socket (AF_IN	ET, SOCK_DGRAM, 0);

The socket number returned is the socket descriptor for the newly created socket. This number is an HP-UX file descriptor and can be used for reading, writing or any standard file system calls. A socket descriptor is treated like a file descriptor for an open file.

Note

To use *write(2)* with a datagram socket, you must declare a default address. See the "Advanced Topics for Datagram Sockets: Specifying a Default Socket Address" section for instructions.

When to Create Sockets

Which Processes	When	

server or client process

before any other IPC system calls

Refer to the *socket(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *socket*.

Binding Socket Addresses to Datagram Sockets

After each process has created a socket, it must call *bind* to bind a socket address. Until an address is bound, other processes have no way to reference it.

The server process must bind a specific port address to its socket. Otherwise, a client process would not know what port to send requests to for the desired service.

The client process can let the local host bind its local port address. The client does not need to know its own port address, and if the server process needs to send a reply to the client's request, the server can find out the client's port address when it receives with *recvfrom*.

Set up the address structure with a local address (as described in the "Preparing Address Variables" section) before you make a *bind* call. Use the wildcard address so your processes do not have to look up their own Internet addresses.

Bind and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <netinet in.h=""></netinet></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	bind (s, addr, addrlen) int s;
	struct sockaddr *addr;
	int addrlen;

Parameter	Descrip	tion of Contents	INPUT Value
s	socket desc socket	riptor of local	socket descriptor of socket to be bound
addr	socket addr	ress	pointer to address to be bound to s
addrlen	length of so	ocket address	size of struct sockaddr_in address
FUNCTION	RESULT:	0 if bind is su -1 if failure o	
EXAMPLE S	SYSTEM	struct sockaddr_ 	in myaddr;
		bind (s, myaddr,	<pre>sizeof(struct sockaddr_in));</pre>

When to Bind Socket Addresses

Which Processes	When
client and server process	after socket is created and before any other IPC system calls

Refer to the *bind(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *bind*.

Sending and Receiving Messages

The *sendto* and *recvfrom* system calls are usually used to transmit and receive messages. They are described in the next sections.

Sending Messages

Use sendto to send messages.

If you have declared a default address (as described in the "Advanced Topics for Datagram Sockets: Specifying a Default Socket Address" section) you can use *send* or *sendto* to send messages. If you use *sendto* in this special case, be sure you specify 0 as the address value, or an error will occur.

Send is described in the "IPC Using Stream Sockets: Sending Data" section of this chapter and in the send(2) entry in the ARPA/Berkeley Services Reference Pages. Sendto and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <netinet in.h=""></netinet></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>count = sendto(s,msg,len,flags,to,tolen)</pre>
	int s;
	char *msg;
	int len, flags;
	struct sockaddr *to;
	int tolen;

Parameter	Description of	INPUT Value
S	socket descriptor of socket	cal socket descriptor of socket sending message
msg	pointer to data buff	pointer to data to be sent
len	size of data buffer	size of msg
flags	settings for optiona	ags 0 (no options are currently sup- ported)
to	address of recipient socket	pointer to the socket address that message should be sent to
tolen	size of to	length of address structure that to points to
FUNCTION		er of bytes actually sent he event of an error

EXAMPLE SYSTEM CALL:

count = sendto(s,argv[2],strlen(argv[2]),0,servaddr,sizeof(struct sockaddr_in));

If the message is too long to send as a single packet (largest size is 2860 bytes for this implementation), an error occurs.

You should not count on receiving error messages when using datagram sockets. The protocol is unreliable, meaning that messages may or may not reach their destination. However, if a message reaches its destination, the contents of the message are guaranteed to be intact.

If you need reliable message transfer, you must build it into your application programs or resend a message if the expected response does not occur.

When to Send Data

Which Processes	When	_
client or server process	after sockets are bound	

Refer to the *send(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *sendto*.

Receiving Messages

Use recvfrom to receive messages.

Recv can also be used if you do not need to know what socket sent the message. However, if you want to send a response to the message, you must know where it came from. Except for the extra information returned by *recvfrom*, the two calls are identical.

Recv is described in the "IPC Using Stream Sockets: Receiving Data" section of this chapter and in the *recv(2)* entry in the *ARPA/Berkeley Services Reference Pages*.

Recvfrom and its parameters are described in the following table.

INCLUDE FILES:	<pre>#include <sys types.h=""></sys></pre>
	<pre>#include <netinet in.h=""></netinet></pre>
	<pre>#include <sys socket.h=""></sys></pre>
SYSTEM CALL:	<pre>count = recvfrom(s,buf,len,flags,from,fromlen)</pre>
	int s;
	char *buf;
	int len, flags;
	struct sockaddr *from;

int *fromlen;

Parameter	Description of Contents	INPUT Value	OUTPUT Value
S	socket descriptor of local socket	socket descriptor of socket receiving message	unchanged
buf	pointer to data buffer	pointer to buffer that is to receive data	pointer to received data
len	maximum number of bytes that should be received	size of data buffer	unchanged
flags	settings for optional flags	0 or MSG_PEEK	unchanged
from	address of socket that sent message	pointer to address structure, not used for input	pointer to socket ad- dress of socket that sent the message
fromlen	pointer to the size of from	pointer to size of from	pointer to the actual size of address returned

FUNCTION RESULT: Number of bytes actually received -1 if an error occurs

EXAMPLE SYSTEM CALL:

```
addrlen = sizeof(sockaddr_in);
...
count = recvfrom(s, buffer, BUFFERSIZE, 0, clientaddr, &addrlen);
```

Recvfrom blocks until there is a message to be received.

11-104 IPC Using Internet Datagram Sockets

No more than *len* bytes of data are returned. The entire message is read in one *recvfrom*, *recv* or *read* operation. If the message is too long for the allocated buffer, the excess data are discarded. Because only one message can be returned in a *recvfrom* call, if a second message is in the queue, it is not affected. Therefore, the best technique is to receive as much as possible on each call.

The host does not wait for *len* bytes to be available; if less than *len* bytes are available, that number of bytes are returned.

Flag Options

The *flags* options are:

- 0 for no options or
- MSG_PEEK for a nondestructive read.

Use the MSG_PEEK option to preview an incoming message. If this option is set on a *recvfrom*, any message returned remains in the data buffer as though it had not been read yet. The next *recvfrom* returns the **same message**.

When to Receive Data

Which Processes	When	
client or server process	after sockets are bound	

Refer to the *recv(2)* entry in the *ARPA/Berkeley Services Reference Pages* for more information on *recvfrom*.

Closing a Socket

In most applications, you do not have to worry about cleaning up your sockets. When you exit your program and your process terminates, the sockets are closed for you.

If you need to close a socket while your program is still running, use the *close* HP-UX file system call.

You may have more than one process with the same socket descriptor if the process with the socket descriptor executes a *fork*. *Close* decrements the file descriptor count and the calling process can no longer use that file descriptor.

When the last *close* is executed on a socket, any unsent messages are sent and the socket is closed. Then the socket is destroyed and can no longer be used.

For syntax and details on *close*, refer to the *close(2)* entry in the *HP-UX Reference* manual.

Example Using Datagram Sockets

These program examples demonstrate how to set up and use datagram sockets. The client program is intended to run in conjunction with the server program.

This example implements a simple name server. The server process receives requests from the client process. It determines the Internet address of the specified host and sends that address to the client process. If the specified host's Internet address is unknown, the server process returns an address of all 1's.

The client process requests the Internet address of a host and receives the results from the server process.

Before you run the example programs:

- make the following entry in the two hosts' /etc/services files: example 22375/udp
- compile with the -lbsdipc option.

The source code for these two programs follows. It is also located in the directory */usr/netdemo/socket*.

/* * SERV. UDP * * This is an example program that demonstrates the use of * datagram sockets as an IPC mechanism. This contains the server, * and is intended to operate in conjunction with the client * program found in client.udp. Together, these two programs * demonstrate many of the features of sockets, as well as good * conventions for using these features. * * This program provides a service called "example". It is an * example of a simple name server. In order for * it to function, an entry for it needs to exist in the * /etc/services file. The port address for this service can be * any port number that is likely to be unused, such as 22375, * for example. The host on which the client will be running * must also have the same entry (same port number) in its * /etc/services file. * */ #include <sys/types.h> #include <sys/socket.h> #include <netinet/in.h> #include <stdio.h> #include <netdb.h> int s: /* socket descriptor */ #define BUFFERSIZE 1024 /* maximum size of packets to be received */ /* contains the number of bytes read */ int cc: char buffer[BUFFERSIZE]; /* buffer for packets to be read into */ struct hostent *hp; /* pointer to host info for requested host */ struct servent *sp: /* pointer to service information */ /* for local socket address */ struct sockaddr in myaddr in; struct sockaddr in clientaddr in; /* for client's socket address */ struct in addr regaddr; /* for requested host's address */ #define ADDRNOTFOUND 0xffffffff /* return address for unfound host */ /* * MAIN * * This routine starts the server. It forks, leaving the child * to do all the work, so it does not have to be run in the * background. It sets up the socket, and for each incoming * request, it returns an answer. Each request consists of a * host name for which the requester desires to know the internet address. The server will look up the name in its

11-108 IPC Using Internet Datagram Sockets

```
*
        /etc/hosts file. and return the internet address to the
 *
        client. An internet address value of all ones will be returned
 *
        if the host name is not found.
 *
*/
main(argc, argv)
int argc;
char *argv∏:
        int addrlen:
                /* clear out address structures */
        memset ((char *)&myaddr in, 0, sizeof(struct sockaddr in));
        memset ((char *)&clientaddr in, 0, sizeof(struct sockaddr in));
                /* Set up address structure for the socket. */
        myaddr_in.sin_family = AF_INET;
                /* The server should receive on the wildcard address,
                 * rather than its own internet address. This is
                 * generally good practice for servers, because on
                 * systems which are connected to more than one
                 * network at once will be able to have one server
                 * listening on all networks at once. Even when the
                 * host is connected to only one network, this is good
                 * practice, because it makes the server program more
                 * portable.
                 */
        myaddr in.sin addr.s addr = INADDR ANY;
                /* Find the information for the "example" server
                 * in order to get the needed port number.
                 */
        sp = getservbyname ("example", "udp");
        if (sp == NULL) {
                printf("%s: example not found in /etc/services\n",
                                argv[0]);
                exit(1);
        }
        myaddr_in.sin_port = sp->s_port;
                /* Create the socket. */
        s = socket (AF INET, SOCK DGRAM, 0);
        if (s == -1) {
                perror(argv[0]);
                printf("%s: unable to create socket\n", argv[0]);
                exit(1);
        }
```

```
/* Bind the server's address to the socket. */
if (bind(s, &myaddr in, sizeof(struct sockaddr in)) == -1) {
        perror(argv[0]):
        printf("%s: unable to bind address\n", argv[0]);
        exit(1):
}
        /* Now, all the initialization of the server is
         * complete, and any user errors will have already
        * been detected. Now we can fork the daemon and
         * return to the user. We need to do a setporp
         * so that the daemon will no longer be associated
         * with the user's control terminal. This is done
         * before the fork. so that the child will not be
         * a process group leader. Otherwise, if the child
         * were to open a terminal, it would become associated
         * with that terminal as its control terminal. It is
         * always best for the parent to do the setporp.
         */
setpgrp();
switch (fork()) {
case -1:
                        /* Unable to fork. for some reason. */
        perror(argv[0]):
        printf("%s: unable to fork daemon\n", argv[0]);
        exit(1):
                        /* The child process (daemon) comes here. */
case 0:
                /* Close stdin, stdout, and stderr so that they will
                 * not be kept open. From now on, the daemon will
                 * not report any error messages. This daemon
                 * will loop forever, waiting for requests and
                 * responding to them.
                 */
        close(stdin):
        close(stdout):
        close(stderr):
                /* This will open the /etc/hosts file and keep
                 * it open. This will make accesses to it faster.
                 */
        sethostent(1);
        for(::) {
                        /* Note that addrlen is passed as a pointer
                         * so that the recyfrom call can return the
                         * size of the returned address.
                         */
                addrlen = sizeof(struct sockaddr in);
```

```
/* This call will block until a new
                                 * request arrives. Then, it will
                                 * return the address of the client.
                                 * and a buffer containing its request.
                                 * BUFFERSIZE - 1 bytes are read so that
                                 * room is left at the end of the buffer
                                 * for a null character.
                                 */
                        cc = recvfrom(s, buffer, BUFFERSIZE - 1, 0,
                                                &clientaddr in, &addrlen);
                        if ( cc == -1) exit(1);
                                /* Make sure the message received is
                                 * null terminated.
                                 */
                        buffer[cc]='\0';
                                /* Treat the message as a string containing
                                 * a hostname. Search for the name in
                                 * /etc/hosts.
                                 *1
                        hp = gethostbyname (buffer);
                        if (hp == NULL) {
                                        /* Name was not found. Return a
                                         * special value signifying the
                                         * error.
                                         */
                                regaddr.s addr = ADDRNOTFOUND;
                        } else {
                                        /* Copy address of host into the
                                         * return buffer.
                                         */
                                regaddr.s addr =
                                       ((struct in addr *)(hp->h addr))->s addr;
                        }
                                /* Send the response back to the
                                 * requesting client. The address
                                 * is sent in network byte order. Note that
                                 * all errors are ignored. The client
                                 * will retry if it does not receive
                                 * the response.
                                 */
                        sendto (s, &regaddr, sizeof(struct in addr),
                                        0, &clientaddr in, addrlen);
                }
        default:
                                /* Parent process comes here. */
                exit(0):
        }
}
```

/* *	CLIE	N T . U D P
* *		ram that demonstrates the use of datagram
*	sockets as an IPC mechan	nism. This contains the client, and is conjunction with the server program found
*	in serv.udp. Together,	these two programs demonstrate many of
*	the features of sockets, these features.	, as well as good conventions for using
*	This program requests a	service called "example". In order for
*	it to function, an entry	y for it needs to exist in the
*		e port address for this service can be likely to be unused, such as 22375, for
*	example. The host on wh	nich the server will be running must also ne port number) in its /etc/services file.
*		
*	application. The host t	s an example of a simple name server that is to provide this service is
*		etc/hosts file. Also, the host providing knows the internet addresses of many
*	hosts which the local ho	ost does not. Therefore, this program
*	the serving host. The s	et address of a target host by name from serving host will return the requested
*	internet address as a response, and will return an address of all ones if it does not recognize the host name.	
* */		
	e <sys types.h=""> e <sys socket.h=""></sys></sys>	
	e <sys errno.h=""></sys>	
	e <netinet in.h=""> e <stdio.h></stdio.h></netinet>	
	e <signal.h></signal.h>	
#includ	e <netdb.h></netdb.h>	
extern	int errno;	
int s;		/* socket descriptor */
	hostent *hp; servent *sp;	/* pointer to host info for nameserver host */ /* pointer to service information */
struct	<pre>sockaddr_in myaddr_in; sockaddr_in servaddr_in; in_addr reqaddr;</pre>	/* for local socket address */ /* for server socket address */ /* for returned internet address */

11-112 IPC Using Internet Datagram Sockets

```
#define ADDRNOTFOUND
                        0xffffffff
                                        /* value returned for unknown host */
#define RETRIES 5
                                /* number of times to retry before giving up */
                        HANDLER
 *
 *
        This routine is the signal handler for the alarm signal.
 *
        It simply re-installs itself as the handler and returns.
*/
handler()
{
        signal(SIGALRM, handler);
}
 *
                        MAIN
* * * * *
        This routine is the client which requests service from the remote
        "example server". It will send a message to the remote nameserver
        requesting the internet address corresponding to a given hostname.
        The server will look up the name, and return its internet address.
 *
        The returned address will be written to stdout.
 *
 *
        The name of the system to which the requests will be sent is given
 *
        as the first parameter to the command. The second parameter should
 *
        be the the name of the target host for which the internet address
 *
        is sought.
*/
main(argc, argv)
int argc:
char *argv[];
{
        int i:
        int retry = RETRIES;
                                       /* holds the retry count */
        char *inet ntoa();
        if (argc != 3) {
                fprintf(stderr, "Usage: %s <nameserver> <target>\n", argv[0]);
                exit(1):
        }
                /* clear out address structures */
        memset ((char *)&myaddr_in, 0, sizeof(struct sockaddr_in));
        memset ((char *)&servaddr in, 0, sizeof(struct sockaddr_in));
```

```
/* Set up the server address. */
servaddr in.sin family = AF INET:
        /* Get the host information for the server's hostname that the
         * user passed in.
         */
hp = gethostbyname (argv[1]);
if (hp == NULL) {
        fprintf(stderr, "%s: %s not found in /etc/hosts\n",
                        argv[0], argv[1]);
        exit(1):
servaddr in.sin addr.s addr = ((struct in addr *)(hp->h addr))->s addr;
        /* Find the information for the "example" server
         * in order to get the needed port number.
         */
sp = getservbyname ("example", "udp");
if (sp == NULL) {
        fprintf(stderr, "%s: example not found in /etc/services\n",
                        argv[0]);
        exit(1);
servaddr in.sin port = sp->s port;
        /* Create the socket. */
s = socket (AF INET, SOCK DGRAM, 0);
if (s == -1) {
        perror(argv[0]);
        fprintf(stderr, "%s: unable to create socket\n", argv[0]);
        exit(1);
}
        /* Bind socket to some local address so that the
         * server can send the reply back. A port number
         * of zero will be used so that the system will
         * assign any available port number. An address
         * of INADDR ANY will be used so we do not have to
         * look up the internet address of the local host.
         */
myaddr in.sin family = AF INET;
myaddr in.sin port = 0;
myaddr in.sin addr.s addr = INADDR ANY;
if (bind(s, &myaddr_in, sizeof(struct sockaddr_in)) == -1) {
        perror(argv[0]);
        fprintf(stderr, "%s: unable to bind socket\n", argv[0]);
        exit(1):
}
        /* Set up alarm signal handler. */
signal(SIGALRM, handler);
```

```
/* Send the request to the nameserver. */
again: if (sendto (s, argv[2], strlen(argv[2]), 0, &servaddr in,
                                sizeof(struct sockaddr in)) == -1) {
                perror(argv[0]);
                fprintf(stderr, "%s: unable to send request\n", argv[0]);
                exit(1):
        }
                /* Set up a timeout so I don't hang in case the packet
                 * gets lost. After all, UDP does not guarantee
                 * delivery.
                 */
        alarm(5);
                /* Wait for the reply to come in. We assume that
                 * no messages will come from any other source,
                 * so that we do not need to do a recvfrom nor
                 * check the responder's address.
                 */
        if (recv (s, &regaddr, sizeof(struct in addr), 0) == -1) {
                if (errno == EINTR) {
                                /* Alarm went off and aborted the receive.
                                 * Need to retry the request if we have
                                 * not already exceeded the retry limit.
                                 */
                        if (--retry) {
                                goto again;
                        } else {
                                printf("Unable to get response from");
                                printf(" %s after %d attempts.\n",
                                                 argv[1]. RETRIES);
                                exit(1);
                } else {
                        perror(argv[0]):
                        fprintf(stderr, "%s: unable to receive response\n",
                                                                 argv[0]):
                        exit(1);
                }
        }
        a larm(0);
                /* Print out response. */
        if (regaddr.s addr == ADDRNOTFOUND) {
                printf("Host %s unknown by nameserver %s.\n", argv[2],
                                                                 argv[1]);
                exit(1);
        } else {
                printf("Address for %s is %s.\n", argv[2],
                        inet_ntoa(regaddr));
        }
}
```

Advanced Topics for Internet Datagram Sockets

Specifying a Default Socket Address

It is possible (but not required) to specify a default address for a remote datagram socket.

This allows you to send messages without specifying the remote address each time. In fact, if you use *sendto*, an error occurs if you enter any value other than 0 for the socket address after the default address has been recorded. You can use *send* or *write* instead of *sendto* once you have specified the default address.

Use *recv* for receiving messages. Although *recvfrom* can be used, it is not necessary, because you already know that the message came from the default remote socket. (Messages from sockets other than the default socket are discarded without notice.) *Read(2)* can also be used, but does not allow you to use the MSG_PEEK flag.

Specify the default address with the connect system call.

When a datagram socket descriptor is specified in a *connect* call, *connect* associates the specified socket with a particular remote socket address. *Connect* returns immediately because it only records the peer's socket address. After *connect* records the default address, any message sent from that socket is automatically addressed to the peer process and only messages from that peer are delivered to the socket.

Connect can be called any number of times to change the associated destination address.

Note

This call does not behave the same as a *connect* for stream sockets. There is no connection, just a default destination. The remote host that you specify as the default may or may not use *connect* to specify your local host as its default remote host. The default remote host is **not** notified if your local socket is destroyed.

Connect and its parameters are described in the following table.

INCLUDE FILES:

#include <sys/types.h>
#include <netinet/in.h>

#include <sys/socket.h>

SYSTEM CALL:

connect(s, addr, addrlen)
int s;
struct sockaddr *addr;
int addrlen;

Parameter	Descrip	tion of Contents	INPUT Value	
S	socket desc socket	riptor of local	socket descriptor of socket request- ing a default peer address	
addr	pointer to the socket address		pointer to socket address of the socket to be the peer	
addrlen	length of address		length of address pointed to by addr	
FUNCTION	RESULT:	0 if connect is -1 if failure of		

When to Specify a Default Socket Address

Which Processes	When
client or server process	after sockets are bound

Synchronous I/O Multiplexing with Select

The *select* system call can be used with sockets to provide a synchronous multiplexing mechanism. The system call has several parameters which govern its behavior. If you specify a zero pointer for the **timout** parameter, *select* will block until one or more of the specified socket descriptors are ready. If timeout is a non-zero pointer, it specifies a maximum interval to wait for the selection to complete.

Select is useful for datagram socket descriptors to determine when data has arrived and is ready to be read without blocking; use the FION-READ parameter to the *ioctl* system call to determine exactly how much data is available.

Selecting for exceptional conditions is currently meaningless for Berkeley sockets. *Select* will always return true for sockets that are no longer capable of being used (e.g. if a *close* or *shutdown* system call has been executed against them).

Select is used the same way as in other applications. Refer to the select(2) entry in the *HP-UX Reference* manual for information on how to use select.

Sending and Receiving Data Asynchronously

Asynchronous sockets allow a user program to receive a SIGIO signal when the state of the socket changes. This state change can occur, for example, when new data arrives. A complete description of SIGIO can be found in the "Advanced Topics for Stream Sockets" section of this manual.

Nonblocking I/O

Sockets are created in blocking mode I/O by default. You can specify that a socket be put in nonblocking mode by using the *ioctl* system call with the FIOSNBIO request.

An example usage of this call is:

```
#include <sys/ioctl.h>
...
ioctl(s, FIOSNBIO, &arg);
```

Arg is a pointer to int:

- When *int* equals 0, the socket is changed to blocking mode.
- When int equals 1, the socket is changed to nonblocking mode.

If a socket is in nonblocking mode, the following calls are affected:

recvfrom	If no messages are available to be received, recvfrom returns the value -1 and the EWOULDBLOCK error. This is also true for <i>recv</i> and <i>read</i> .
-	

sendto If there is no available message space for the message to be transmitted, sendto returns the value -1 and the EWOULDBLOCK error. The O_NDELAY flag for *fcntl(2)* is also supported. If you use this flag and there is no message available to be received on a *recv*, *recvfrom*, or *read* call, the call returns immediately with the value of 0. This is the same as returning an end-of-file condition. This is also true for *send*, *sendto*, and *write* if there is not enough buffer space to complete the send.

Note

The O_NDELAY flag has precedence over the FIOSNBIO flag.

Using Broadcast Addresses

In place of a unique internet address or the wildcard address, you can also specify a broadcast address. The broadcast address is a local address portion of the internet address equal to all 1's. You must be the superuser to use a broadcast address.

If you use broadcast addressing, be careful not to overload your network.

Programming Hints

Note

Refer to the "Portability Issues" appendix for information about the differences between 4.2 BSD and the HP-UX implementation of IPC.

Troubleshooting

You can avoid many problems by using good programming and debugging techniques. Your programs should check for a returned error after each system call and print any that occur. For example, the following program lines print an error message for *read*:

```
cc=read(sock,buffer,1000);
if (cc<0) {
    perror ("reading message")
    exit(1)
}
```

Refer to the *HP-UX Reference* manual for information about *perror(3C)*. Refer to the *ARPA/Berkeley Services Reference Pages* for information about errors returned by the IPC system calls such as *read*.

You can also compile your program with the debugging option (-g) and use one of the debuggers (e.g. cdb or xdb) to help debug the programs.

Port Addresses

The following port values are reserved for the super-user: 1 - 1023, 1260, 1536, 1542 and 4672. These ports are for:

Port Addresses	Used By
1 - 1023	ARPA/Berkeley services
1260	NS daemon rlbdaemon
1536	NS daemon nftdaemon
1542	NS service Remote Process Management (Series 500 only)
4672	NS daemon rfadaemon

It is possible that you could assign one of these ports and cause a service to fail. For example, if the *nftdaemon* is not running, and you assign its port, the *nftdaemon* will fail when you try to start it.

Using Diagnostic Utilities as Troubleshooting Tools

You can use the following diagnostic utilities to help debug your programs. It is helpful if you have multiple access to the system so you can obtain information about the program while it is running.

ping	Use <i>ping</i> to verify the physical connection with the destination node.
netstat	Use the <i>netstat</i> displays of sockets and associations to help you troubleshoot problems in your applica- tion programs. Use <i>netstat</i> to determine if your program has successfully created a connection. If you are using stream sockets (TCP protocol), <i>netstat</i> can provide the TCP state of the connection. To check the status of a connection at any point in the program, use the <i>sleep</i> (seconds) statement in your program to pause the program. While the program is paused, execute <i>netstat -a</i> from another terminal.
Network Tracing	<i>Network Tracing</i> can be used to trace packets. For the trace information to be useful, you must have a working knowledge of network protocols.
Network Event Logging	<i>Network Event Logging</i> is an error logging mechanism. Use it in conjunction with other diagnostic tools.

These utilities are described in detail in the Installing and Maintaining NS-ARPA Services manual.

Adding a Server Process to the Internet Daemon

This section contains example IPC programs that use the internet daemon, called *inetd*. For more information on *inetd*, refer to the "Configuration and Maintenance" chapter of the *Installing and Maintaining NS-ARPA Services* manual and the *inetd(1M)* entry in the *ARPA/Berkeley Services Reference Pages*.

You can invoke the example server programs from *inetd* if you have **super-user** capabilities and you make the following configuration modifications:

• Add the following lines to the /etc/inetd.conf file:

example stream tcp nowait root <path>/server.tcp server.tcp example dgram udp wait root <path>/server.udp server.udp

where <path> is the path to the files on your host. (For detailed information on this file, refer to the "Configuration and Maintenance" chapter of the *Installing and Maintaining NS-ARPA Services* manual or to the *inetd.conf(4)* entry in the *ARPA/Berkeley Services Reference Pages*.)

• Add the following lines to the /etc/services file:

example 22375/tcp example 22375/udp

• If *inetd* is already running, execute the following command so that *inetd* recognizes the changes:

/etc/inetd -c

These example programs do the same thing as the previous example servers do, but they are designed to be called from *inetd*. They do not have daemon loops or listen for incoming connection requests, because *inetd* does that. The source code for the two example servers follows.

SERVER. TCP This is a variation of the example program called serv.tcp. This one performs the same function, except that it is designed to be called from /etc/inetd. Hence, this version does not contain a daemon loop, and does not listen for incoming connections on the socket. /etc/inetd does these functions. This server simply assumes that the socket to receive the messages from and send the responses to is file descriptor 0 when the program is started. It also assumes that the client's connection is already established to the socket. For the sake of simplicity, the activity logging functions of serv.tcp have also been removed. */ MAIN This is the actual server routine that the /etc/inetd forks to handle each individual connection. Its purpose is to receive the request packets from the remote client, process them, and return the results to the client. */ main() char buf[10]; /* This example uses 10 byte messages. */ int len, len1; /* Go into a loop, receiving requests from the remote * client. After the client has sent the last request. * it will do a shutdown for sending, which will cause * an end-of-file condition to appear on this end of the * connection. After all of the client's requests have * been received, the next recv call will return zero * bytes, signalling an end-of-file condition. This is * how the server will know that no more requests will * follow, and the loop will be exited. */ while (len = recv(0, buf, 10, 0)) { if (len == -1) { exit (1); /* error from recv */ } /* The reason this while loop exists is that there * is a remote possibility of the above recv returning * less than 10 bytes. This is because a recy returns * as soon as there is some data, and will not wait for * all of the requested data to arrive. Since 10 bytes * is relatively small compared to the allowed TCP

1* *

*

*

*

*

*

*

*

*

*

*

1* *

> * *

> > *

*

*

*

{

* packet sizes, a partial receive is unlikely. If

```
* this example had used 2048 bytes requests instead,
                 * a partial receive would be far more likely.
                 * This loop will keep receiving until all 10 bytes
                 * have been received, thus guaranteeing that the
                 * next recv at the top of the loop will start at
                 * the beginning of the next request.
                 */
        while (len < 10) {
                len1 = recv(0, &buf[len], 10-len, 0);
                if (len1 == -1) {
                        exit (1):
                len += len1:
        }
                /* This sleep simulates the processing of the
                 * request that a real server might do.
                 */
        sleep(1);
                /* Send a response back to the client. */
        if (send(0, buf, 10, 0) != 10) {
                exit (1);
        }
}
        /* The loop has terminated, because there are no
         * more requests to be serviced.
         */
exit (0);
```

```
}
```

```
SERVER, UDP
        This is a variation of the example program called serv.udp.
        This one performs the same function, except that it is
        designed to be called from /etc/inetd. Hence, this version
        does not contain a daemon loop, and does not wait for requests
 *
        to arrive on a socket. /etc/inetd does these functions. This
 *
        server simply assumes that the socket to receive the message
 *
        from and send the response to is file descriptor 0 when
 *
        the program is started. It also assumes that the client's
 *
        request is already ready to be received from the socket.
 *
*/
#include <sys/types.h>
#include <netinet/in.h>
#include <stdin h>
#include <netdb.h>
#define BUFFERSIZE
                       1024
                                /* maximum size of packets to be received */
int cc:
                                /* contains the number of bytes read */
                               /* buffer for packets to be read into */
char buffer[BUFFERSIZE]:
struct hostent *hp;
                               /* pointer to host info for requested host */
                                       /* for client's socket address */
struct sockaddr in clientaddr in;
                                /* for requested host's address */
struct in addr regaddr;
#define ADDRNOTFOUND
                        Oxffffffff
                                        /* return address for unfound host */
                        MAIN
        This routine receives the request and returns an answer.
 *
        Each request consists of a
 *
        host name for which the requester desires to know the
 *
        internet address. The server will look up the name in its
 *
        /etc/hosts file, and return the internet address to the
 *
        client. An internet address value of all ones will be returned
 *
        if the host name is not found.
 *
 */
main()
{
        int addrlen;
                /* clear out address structure */
        memset ((char *)&clientaddr in, 0, sizeof(struct sockaddr in));
```

```
/* Note that addrlen is passed as a pointer
         * so that the recvfrom call can return the
        * size of the returned address.
         */
addrlen = sizeof(struct sockaddr in):
        /* This call will
        * return the address of the client.
        * and a buffer containing its request.
         * BUFFERSIZE - 1 bytes are read so that
         * room is left at the end of the buffer
         * for a null character.
         */
cc = recvfrom(0, buffer, BUFFERSIZE - 1, 0, &clientaddr in, &addrlen);
if ( cc == -1) exit(1);
        /* Make sure the message received is
         * null terminated.
         */
buffer[cc]='\0';
        /* Treat the message as a string containing
         * a hostname. Search for the name in
         * /etc/hosts.
         */
hp = gethostbyname (buffer);
if (hp == NULL) {
                /* Name was not found. Return a
                 * special value signifying the
                 * error.
                 */
        reqaddr.s_addr = ADDRNOTFOUND;
} else {
                /* Copy address of host into the
                 * return buffer.
                 */
        reqaddr.s_addr =
               ((struct in addr *)(hp->h addr))->s addr;
}
        /* Send the response back to the
         * requesting client. The address
         * is sent in network byte order. Note that
         * all errors are ignored. The client
         * will retry if it does not receive
         * the response.
         */
sendto (0, &regaddr, sizeof(struct in addr), 0,
                        &clientaddr in, addrlen);
exit(0);
```

```
11-128 Adding a Server Process to the Internet Daemon
```

}

Summary Tables for System and Library Calls

The following table contains a summary of the IPC system calls.

IPC System Calls

System Call	Description
socket	Creates a socket, or communication endpoint for the calling process.
bind	Assigns a socket address to the socket specified by the calling process.
listen	Sets up a queue for incoming connection requests. (Stream sockets only.)
connect	For stream sockets, requests and creates a connec- tion between the remote socket (specified by ad- dress) and the socket (specified by descriptor) of the calling process.
	For datagram sockets, permanently specifies the remote peer socket.
accept	Receives a connection between the socket of the call- ing process and the socket specified in the associated connect call. (Stream sockets only.)
send, sendto	Sends data from the specified socket.
recv, recvfrom	Receives data at the specified socket.
shutdown	Disconnects the specified socket.
getsockname	Gets the socket address of the specified socket.

System Call	Description
getsockopt, setsockopt	Gets, or sets, the options associated with a socket.
getpeername	Gets the name of the peer socket connected to the specified socket.

11-130 Summary Tables for System and Library Calls

The following table contains a summary of the other system calls that can be used with IPC.

Other System Calls

System Call Description

- *read* Can be used to read data at stream or datagram sockets just like *recv* or *recvfrom*, without the benefit of the *recv* flags. *Read* offers implementation independence; the descriptor can be for a file, a socket or any other object.
- *write* Can be used to write data from stream sockets (and datagram sockets if you declare a default remote socket address) just like *send*. *Write* offers implementation independence; the descriptor can be for a file, a socket or any other object.
- *close* Deallocates socket descriptors. The last *close* can be used to destroy a socket. *Close* does a graceful disconnect or a hard close, depending on the LINGER option. Refer to the "Closing a Socket" sections of this chapter.
- select Can be used to improve efficiency for a process that accesses multiple sockets or other I/O devices simultaneously. Refer to the "I/O Multiplexing with Select" sections of this chapter.
- *ioctl* Can be used for finding the number of receivable bytes with FIONREAD and for setting the nonblocking I/O flag. Can also be used for setting a socket to receive asynchronous signals with FIOASYNC.
- *fcntl* Can be used for duplicating a socket descriptor and for setting the O_NDELAY flag.

IPC attempts to isolate host-specific information from applications by providing library calls that return the necessary information.

The following table contains a summary of the library calls used with IPC. The library calls are in the common "c" library named *libc.a.* Therefore, there is no need to specify any library name on the cc command line to use these library calls - *libc.a* is used automatically.

Library Call	Description	
htonl htons ntohl ntohs	convert values between host and network byte order (for portability to DEC VAX hosts)	
inet_addr inet_lnaof inet_makeaddr inet_netof inet_network	internet address manipulation routines	
setservent endservent getservbyname getservbyport getservent	get or set service entry	
setprotoent endprotoent getprotobyname getprotobynumber getprotoent	get or set protocol entry	

Library Calls

Library Call

Description

setnetent endnetent getnetbyaddr getnetbyname getnetent	get or set network entry
sethostent endhostent gethostbyaddr gethostbyname	get or set host entry

gethostent

Interprocess Communication 11-133

11-134 Summary Tables for System and Library Calls

Portability Issues

This appendix describes implementation differences between 4.2 BSD IPC and HP-UX IPC. It contains porting issues for:

- IPC functions and library calls; and for
- other functions and library calls typically used by IPC programs.

Because HP-UX IPC is based on 4.2 BSD IPC (it is a subset of 4.2 BSD), programs should port easily between HP-UX and 4.2 BSD systems. If you need to have portable applications, keep the information in this appendix in mind when you write your IPC programs.

Porting Issues for IPC Functions and Library Calls

The following is a list of differences in IPC functions and library calls to be aware of if you want to port your IPC applications between HP-UX and 4.2 BSD systems.

Shutdown

When *shutdown* has been used on a datagram socket on an HP-UX system, the local port number bound to that socket remains unavailable for use until that socket has been destroyed by *close*.

Some other systems free that port number for use immediately after the *shutdown*. In general, sockets should be destroyed by *close* (or by terminating the process) when they are no longer needed. This allows you to avoid unnecessary delay in deallocating local port numbers.

Address Conversion Functions for DEC VAX Hosts

The functions *htonl*, *htons*, *ntonl* and *ntons* are not required on HP-UX systems. They are included for porting to a DEC VAX host. You can use these functions in your HP-UX programs for portability; they are defined as null macros on HP-UX systems, and are found in <*netinet/in.h*>.

FIONREAD Return Values

For HP-UX systems, the FIONREAD *ioctl* request on a datagram socket returns a number that may be larger than the number of bytes actually readable. Previously, HP-UX systems returned the maximum number of bytes that a subsequent *recv* would be able to return.

Listen's Backlog Parameter

HP-UX treats the *listen(2)* backlog value as the actual size of the queue for pending connections. Some implementations set their queue size to 3/2 * B + 1, where B is the backlog value.

Pending Connections

There is no guarantee as to which pending connection on a listening socket will be returned by *accept*. HP-UX systems return the newest pending connection. Applications should be written such that they do not depend upon connections being returned by *accept* on a first-come, firstserved basis.

Errno Values

HP-UX IPC system calls have some *errno* values that are different from other implementations. These are listed in the following table.

System Call	Error	HP-UX Implementation	Other Implementations
connect	socket is a listening socket	EINVAL	ETIMEDOUT
socket	invalid socket type	EPROTONOSUPPORT and EPROTOTYPE	EPROTOTYPE
socket	invalid protocol	EPROTONOSUPPORT	EPROTOTYPE

Losing a TCP Connection

On a stream socket connection, if the connection has been lost due to some error, HP-UX systems return the same *errno* value for each subsequent *recv*. Some other implementations only return the error on the first *recv* after the connection is lost, and then return the end-of-file condition on subsequent *recv* calls.

Unsupported IPC Features

The following is a list of 4.2 BSD IPC features which are not supported on HP-UX systems.

The HP-UX implementation does not support:

- AF_UNIX or other addressing domains (only AF_INET is supported);
- the use of readv(2) and writev(2) on sockets;
- the sendmsg(2) and recvmsg(2) system calls; or
- the SOCK_RAW socket type.

Porting Issues for Other Functions and Library Calls Typically Used by IPC

The following is a list of differences in functions and library calls to be aware of when you port your IPC applications between HP-UX and 4.2 BSD systems.

loctl and Fcntl Calls

4.2 BSD terminal *ioctl* calls are incompatible with the HP-UX implementation. These calls are typically used in virtual terminal applications. The HP-UX implementation uses UNIX System V compatible calls.

Pty Location

Look for the *pty* masters in */dev/ptym/ptyp?* and for the *pty* slaves in */dev/pty/ttyp?*. An alternative location to check is */dev*.

Dup2

You must use the -IBSD compile option to use the 4.1 or 4.2 BSD version of the dup2(2) system call on an HP-UX system.

Size Limit for Send

For Series 300 HP-UX systems, the maximum size message that can be sent on a datagram socket or on a nonblocking stream socket is 9216 bytes. For Series 800 HP-UX systems, the maximum size message is 2048 bytes. For 4.2 BSD systems, the maximum size message is 2048 bytes.

Utmp

The 4.2 BSD /*etc*/*utmp* file format is incompatible with the HP-UX implementation. The HP-UX implementation uses UNIX System V compatible calls. Refer to the utmp(5) entry in the HP-UX Reference manual for details.

Library Equivalencies

Certain commonly used library calls in 4.2 BSD are not present in HP-UX systems, but they do have HP-UX equivalents. To make code porting easier, use the following equivalent library calls. You can do this by putting them in an include file, or by adding the define statements (listed in the following table) to your code.

	4.2 BSD Library	HP-UX Library	
#define	index(a,b)	strchr(a,b)	
#define	rindex(a,b)	strrchr(a,b)	
#define	bcmp(a,b,c)	memcmp(a,b,c)	
#define	bcopy(a,b,c)	memcpy(b,a,c)	
#define	bzero(a,b)	memset(a,0,b)	
#define	getwd(a)	getcwd(a,MAXPATHLEN)	

Definition of Library Equivalents

Note

Include *< string.h >* before using *strchr* and *strrchr*. Include *< sys/param.h >* before using *getcwd*.

Signal Calls

Normal HP-UX *signal* calls are different from 4.2 BSD signals. See the *sigvector(2)* entry in the *HP-UX Reference* manual for information on signal implementation. Note the following signal mapping.

Definitions of Signal Equivalents

4.2 BSD Signal	is mapped to	HP-UX Signal
SIGCHLD		SIGCLD

Sprintf Return Value

For 4.2 BSD, *sprintf* returns a pointer to a string. For HP-UX systems, *sprintf* returns a count of the number of characters in the buffer.

A-8 Porting Issues for Other Functions and Library Calls Typically Used by IPC

Glossary

Account name:	A synonym for user name or login name.
Address family:	The address format used to interpret addresses specified in socket operations. The internet address family (AF_INET) is supported.
Address:	An Interprocess Communication term that refers to the means of labeling a socket so that it is distin- guishable from other sockets, and routes to that sock- et are able to be determined.
Advanced Re- search Projects Agency:	A U.S. government research agency that was in- strumental in developing and using the original ARPA Services on the ARPANET.
Alias:	A term used to refer to alternate names for net- works, hosts and protocols. This is also an internet- work mailing term that refers an alternate name for a recipient or list of recipients (a mailing list).
ARPA:	See "Advanced Research Projects Agency."
ARPA/Berkeley Services:	The set of services originally developed for use on the ARPANET (i.e., $telnet(1)$) or distributed with the Berkeley Software Distribution of UNIX, version 4.2 (i.e., $rlogin(1)$).
ARPANET:	The Advanced Research Projects Agency Network.

Association:	An Interprocess Communication connection (e.g., a socket) is defined by an association. An association contains the (protocol, local address, local port, remote address, remote port)-tuple. Associations must be unique ; duplicate associations on the same system may not exist.
Asynchronous Sockets:	Sockets set up via <i>ioctl</i> with the FIOASYNC option to be notified with a SIGIO signal whenever a change on a socket occurs. It is primarily used for sending and receiving data without blocking.
Berkeley Software Distribution:	A version of UNIX software released by the Univer- sity of California at Berkeley.
Binding:	Establishing the address of a socket which allows other sockets to connect to it or to send data to it.
BSD:	See "Berkeley Software Distribution."
Channel:	A communication path created by establishing a con- nection between sockets.
Client:	A process that is requesting some service from another process.
Client host:	The host on which a client process is running.
Communication domain:	A set of properties that describes the characteristics of processes communicating through sockets. Only the Internet domain is supported.
Connection:	A communications path to send and receive data. A connection is uniquely identified by the pair of sockets at either end of the connection. See also, "Association."
Daemon:	A software process that runs continuously and provides services on request.

	DARPA:	See "Defense Advanced Research Projects Agency."
0	Datagram sockets:	A socket that maintains record boundaries and treats data as individual messages rather than a stream of bytes. Messages may be sent to and received from many other datagram sockets. Datagram sockets do not support the concept of a connection. Messages could be lost or duplicated and may not arrive in the same sequence sent. Datagram sockets use the User Datagram Protocol.
	Defense Advanced Re- search Projects Agency:	The military arm of the Advanced Research Projects Agency. DARPA is instrumental in defining stand- ards for ARPA services.
	Domain:	A set of allowable names or values. See also, "Com- munication domain."
\bigcirc	Equivalent account:	An account (or user name) specified in the <i>\$HOME/.rhosts</i> file that allows the specified remote users to access the local user's account without requiring a password.
	Equivalent host:	A remote host that is considered by your local host as an "equivalent computer." Users from equivalent hosts can bypass password validation if they have the same account name on both hosts.
	Equivalent user:	See "Equivalent account."
	File Transfer Protocol:	The file transfer protocol that is traditionally used in ARPA networks. The <i>ftp</i> command uses the FTP protocol.
\bigcirc	Forwarding:	The process of forwarding a mail message to another destination (i.e., another user name, host name or network).
	4.2 BSD:	See "Berkeley Software Distribution."

Frame:	See "Packet."
FTP:	See "File Transfer Protocol."
Gateway:	A node that connects two or more networks together and routes packets between those networks.
Host:	A node that has primary functions other than switch- ing data for the network.
International Standards Organization:	Called "ISO," this organization created a network model that identifies the seven commonly-used protocol levels for networking.
Internet:	All ARPA networks that are registered with the Net- work Information Center.
Internet address:	A four-byte quantity that is distinct from a link-level address and is the network address of a computer node. This address identifies both which network is on the Internet and which host is on the network.
Internetwork:	A term used to mean "among different physical net- works."
Interprocess Communication:	A facility that allows a process to communicate with another process on the same host or on a remote host. IPC provides system calls that access sockets. This facility is distinct from Bell System V IPC. See also, "Sockets."
IPC:	See "Interprocess Communication."
ISO:	See "International Standards Organization."
Link-level address:	A six-byte quantity that is distinct from the internet address and is the unique address of the LAN inter- face card on each LAN.

Message:	In IPC, the data sent in one UDP packet. When using <i>sendmail</i> a message is the information unit transferred by mail.
Node:	A computer system that is attached to or is part of a computer network.
Node manager:	The person who is responsible for managing the net- working services on a specific node or host.
Official host name:	The first host name in each entry in the <i>/etc/hosts</i> file. The official host name cannot be an alias.
Packet:	A data unit that is transmitted between processes. Also called a "frame."
Peer:	An Interprocess Communication socket at the other end of a connection.
Port:	An address within a host that is used to differentiate between multiple sockets with the same internet ad- dress.
Protocol:	A set of conventions for transferring information be- tween computers on a network (e.g., UDP or TCP).
Remote host:	A computer that is accessible through the network or via a gateway.
Reserved port:	A port number between 1 and 1023 that is only for super-user use.
Server:	A process or host that performs operations that local or remote client hosts request.
Service:	A facility that uses Interprocess Communication to perform remote functions for a user (e.g., <i>rlogin(1)</i> or <i>telnet(1)</i>).

Simple Mail Transfer Protocol:	A standard protocol for transporting mail reliably and efficiently through LANs or the ARPANET. <i>Sendmail</i> uses the SMTP protocol.		
SMTP:	See "Simple Mail Transfer Protocol."		
Socket:	Addressable entities that are at either end of an In- terprocess Communication connection. A socket is identified by a socket descriptor. A program can write data to and read data from a socket, just as it writes and reads data to and from files.		
Socket address:	The internet address, port address and address fami- ly of a socket. The port and internet address com- bination allows the network to locate a socket.		
Socket descriptor:	An HP-UX file descriptor accessed for reading, writ- ing or any standard file system calls after an Inter- process Communication connection is established. All Interprocess Communication system calls use socket descriptors as arguments.		
Stream socket:	A socket that, when connected to another stream socket, passes data as a byte stream (with no record boundaries). Data is guaranteed to arrive in the se- quence sent. Stream sockets use the TCP protocol.		
TCP:	See "Transmission Control Protocol."		
Telnet:	A virtual terminal protocol traditionally used on ARPA networks that allows a user to log into a remote host. The <i>telnet</i> command uses the Telnet protocol.		
Transmission Control Protocol:	A protocol that provides the underlying communica- tion support for AF_INET stream sockets. TCP is used to implement reliable, sequenced, flow-control- led two-way communication based on a stream of bytes similar to pipes.		

UDP:	See "User Datagram Protocol."	
User Datagram Protocol:	A protocol that provides the underlying communica- tion support for datagram sockets. UDP is an unreli- able protocol. A process receiving messages on a datagram socket could find that messages are dupli- cated, out-of-sequence or missing. Messages retain their record boundaries and are sent as individually addressed packets. There is no concept of a connec- tion between the communicating sockets.	
Virtual Terminal Protocol:	A protocol that provides terminal access to interac- tive services on remote hosts (e.g., <i>telnet(1)</i>).	
UNIX Domain Address:	A character string continuing the UNIX pathname to a UNIX Domain socket.	
UNIX Domain Protocol:	A protocol providing fast communication between processes executing on the same node and using the AF_UNIX socket address family.	

Index

I

! command ftp, 8–10 telnet, 6-17 #, 8-36 \$HOME directory, 7-8 \$HOME/.cshrc, 9-7, 10-17 \$HOME/.login, 7-14, 7-26, 9-7, 10 - 17\$HOME/.netrc, 8-75, 8-77 creation of local, 8-75 protection of local, 8-76 \$HOME/.profile, 7-14, 7-26, 9-7, 10 - 17\$HOME/.rhosts, Glossary 3, 2-8, 7-7, 7-22 to 7-23, 9-3, 9-16, 9-30, 10-2, 10-7 creation of local, 7-24, 9-31, 10-8 creation of remote, 7-8, 9-3, 10-3 protection of, 7–9 protection of local, 7-25, 9-32, 10-9 protection of remote, 9-5, 10-4 .cshrc file, 9-7, 10-17 *login* file, 7–14, 7–26, 9–7, 10–17. .netrc file, 2-7, 8-75, 8-77 creation of local, 8-75 protection of local, 8-76 .profile file, 7-14, 7-26, 9-7, 10-17

.rhosts file, Glossary 3, 2-8, 7-7, 7-22 to 7-23, 9-3, 9-16, 9-30, 10-2, 10-7 creation of local, 7-24, 9-31, 10 - 8creation of remote, 7-8, 9-3, 10 - 3protection of, 7-9 protection of local, 7-25, 9-32, 10 - 9protection of remote, 9-5, 10-4 /dev/null, 10–15 /etc/csh.login, 7-14 /etc/hosts, Glossary 5, 2-7, 6-7 /etc/hosts.equiv, 2-7, 7-7, 9-3, 9-16, 10-2/etc/networks, 2–7 /etc/profile, 7-14 /etc/protocols, 2–8 /etc/services, 2-8 /usr/bin/remsh, 7–26 /usr/hosts, 7-26, 10-17 ? command ftp, 8–9 telnet, 6-14 ~, 7-2, 7-10, 7-12

Α

accept, A-3, 11-10, 11-19, 11-25, 11-45, 11-51, 11-57, 11-83, 11-129 account, 8-78 account name, Glossary 1 account prompt, ftp, 8-6 address, Glossary 1 address conversion call, 11-30 address family, Glossary 1 addressing domain, 11-12, 11-47 Advanced Research Projects Agency, Glossary 1, 2–1, 2–3 Advanced Research Projects Agency Network, Glossary 1 Advanced Topics for Datagram Sockets, 11-116 AF INET, Glossary 1, 11–10, 11–94 alias, Glossary 1 anonymous ftp account, 8-81 anonymous *ftp* account, login to, 8 - 82append command, ftp, 8-51, 8-55, 8 - 66ARPA, Glossary 1, 2–1, 2–3 ARPA Services, 2-1, 2-3 ARPA/Berkeley Services, Glossary 1 **ARPANET**, Glossary 1 ascii command, ftp, 8-35 ascii file transfer type in *ftp*, 8–33 association, Glossary 2 asynchronous sockets, Glossary 2, 11-81, 11-118

В

bcmp, A-6 bcopy, A–6 bell command, ftp, 8-36 bell sound for ftp file transfer completion, 8-36, 8-74 Berkeley Services, 2-1, 2-3 Berkeley Software Distribution, Glossary 1 to Glossary 2, 2–1 binary command, ftp, 8–34 binary file transfer type in ftp, 8-33 to 8-34 bind, 11-10, 11-45, 11-74, 11-92, 11-99, 11-129 binding, Glossary 2, 11-3 blocking mode, 11–119 break, 7–11, 7–14 broadcast address, 11-120 BSD, Glossary 1 to Glossary 2, 2 - 1bye command, ftp, 8–7 bzero, A-6

С

C programming language, 1–6 carriage return behavior in *telnet*, 6–10 *cd* command, *ftp*, 8–16 channel, Glossary 2 characters, terminal configuration, 6–5, 7–2 client, Glossary 2 client host, Glossary 2 client-server model, 11–4

close, A-2, 11-10, 11-29, 11-45, 11-61, 11-78, 11-85, 11-106, 11-131 close command, telnet, 6-12, 6-24 combination copies *rcp* local and remote to local, 9–22 *rcp* local and remote to remote, 9–26 rcp local to remote, 9-10 rcp remote to local, 9-14 rcp remote to remote, 9–18 command search path, 7-26 command state, telnet 6-2, 6-13 command, remsh execution of remote, 10-5 communication domain, Glossary 2 concepts, 1-1 configuration, 1-6 connect, A-3, 11-23, 11-55, 11-74, 11-116, 11-129 connection, Glossary 2 connectivity, 1-1 consumer, 9-2 control character entry as *rlogin* escape character, 7-10, 7-12 entry as telnet escape character, 6-6 conventions, 1-5 crmod command, telnet, 6-10 csh.login file, 7-14 CTRL-], 6-3, 6-9

D

daemon, Glossary 2 DARPA, Glossary 3 datagram socket, Glossary 3 datagram sockets, 2–8, 11–9, 11–92, 11–107, 11–116, 11–129 Defense Advanced Research Projects Agency, Glossary 3 delete command, ftp, 8-30, 8-68 *dir* command, *ftp*, 8–17, 8–20 directories, ftp listing of multiple remote, 8-23 directory changing ftp local working, 8-15 changing *ftp* remote working, 8 - 16ftp creation of remote, 8–29 ftp deletion of remote, 8-30 ftp listing of remote, 8–20 listing *ftp* remote working, 8–17 directory copies rcp local and remote to local, 9 - 22*rcp* local and remote to remote, 9 - 26rcp local to remote, 9-10 rcp remote to local, 9-14 rcp remote to remote, 9-18 directory name ftp change of remote, 8-31 ftp display of remote working, 8 - 28documentation map, 1-1 domain, Glossary 2 to Glossary 3 dup2, A-5

E

endhostent, 11–133 endnetent, 11–133 endprotoent, 11–132 endservent, 11–132 equivalent account, Glossary 3 equivalent host, Glossary 3, 2–7 equivalent user, Glossary 3, 2–8 errno, A–3 escape character *rlogin*, 7–1 to 7–2, 7–10, 7–12, 7–15 to 7–16, 7–18 to 7–19 *telnet*, 6–2 to 6–4, 6–6, 6–9, 6–20 *escape* command, *telnet*, 6–4

F

fcntl, 11-84, 11-120, 11-131 file use of *ftp* to append local to remote, 8-51, 8-55 use of *ftp* to append text to remote, 8-65 use of *ftp* to change name of remote, 8-72 use of *ftp* to create remote, 8-64 use of *ftp* to delete remote, 8–68 use of *ftp* to display remote, 8–62 use of *ftp* to move remote, 8–72 use of *ftp* to delete multiple remote, 8-69 file attributes rcp's effect on, 9-28 file copies *rcp* local and remote to local, 9–20 *rcp* local and remote to remote, 9–24 rcp local to remote, 9-8 rcp remote to local, 9-12 rcp remote to remote, 9-16 file operations in *ftp*, remote, 8–62 file transfer *ftp* local to remote multiple, 8 - 57ftp local to remote single,

8-49, 8-53 ftp remote to local multiple, 8-44 ftp remote to local single, 8-40, 8-42 ftp selective, 8–37 file transfer environment in *ftp*, 8-33 file transfer options in ftp, 8-39 File Transfer Protocol, Glossary 3 file transfer type in ftp, 8-33 file transfer type changing to ftp ascii, 8-35 changing to ftp binary, 8-34 displaying current ftp, 8-34 FIOASYNC, 11-81, 11-131 FIONREAD, A-2, 11-79, 11-118, 11-131 FIOSNBIO, 11-83, 11-120 font conventions, 1–5 forwarding, Glossary 3 frame, Glossary 5 ftp, Glossary 3, 2-6, 8-1 ! command, 8-10 -g option, 8–83 -i option, 8–83 **-n** option, 8–77, 8–83 -v option, 8-83 ? command, 8-9 account prompt, 8-6 anonymous account, 8-81 anonymous account, login to, 8 - 82append command, 8-51, 8-55, 8-66 ascii command, 8-35 ascii file transfer type, 8-33 automatic remote login, 8–75

automatic remote login, disabling, 8-77, 8-83 bell command, 8-36 bell sound for file transfer completion, 8-36, 8-74 binary command, 8–34 binary file transfer type, 8-33 to 8-34 bye command, 8–7 cd command, 8-16 command descriptions, 8-9 command list, 8–9 connection to remote host, 8-4, 8-74, 8-83 delete command, 8-30, 8-68 *dir* command, 8–17, 8–20 directories, listing multiple remote, 8 - 23directory changing local working, 8-15 changing name of remote, 8-31 changing remote working, 8-16 creating remote, 8–29 deleting remote, 8–30 displaying name of remote working, 8–28 listing remote, 8-20 listing remote working, 8-17 directory operations, 8-14 disconnection from remote host, 8-7 to 8-8 display of remote responses, 8-2 execution, 8-2 exit from, 8–7 file appending local to remote, 8-51, 8-55 appending text to remote, 8-65 changing name of remote, 8–72 creating remote, 8-64

deleting multiple remote, 8–69 deleting remote, 8–68 displaying remote, 8-62 moving remote, 8–72 file transfer local to remote multiple, 8–57 local to remote single, 8–49, 8-53 remote to local multiple, 8-44 remote to local single, 8-40, 8-42 selective, 8-37 file transfer environment, 8-33 file transfer options, 8–39 file transfer progress, monitoring, 8-36 file transfer type, 8–33, 8–74 file transfer type changing to ascii, 8-35 changing to binary, 8-34 displaying current, 8-34 get command, 8-40, 8-42, 8-62 glob command, 8-12 globbing, 8-12, 8-74, 8-83 globbing behavior of commands in, 8-13 to 8-14 guest account, 8-81 guest account, login to, 8-82 hash command, 8-36 hash sign file transfer progress indicator, 8–36 hash sign for file transfer progress, 8-74 help command, 8–9 interactive mode, 8-24, 8-26, 8-37, 8-44, 8-47, 8-57, 8-60, 8-69, 8-71, 8-74, 8-83 *lcd* command, 8–15 local work within, 8–10

ls command, 8–17, 8–20 mdelete command, 8-69, 8-71 mdir command, 8-23, 8-26 metacharacter expansion in, 8–13 to 8–14 metacharacter use within, 8–12 mget command, 8-44, 8-47 *mkdir* command, 8–29 mls command, 8–23, 8–26 *mput* command, 8–57, 8–60 open command, 8-4 prompt command, 8–37 public account, 8–81 public account, login to, 8-82 public directory structure, 8–81 put command, 8-49, 8-53, 8-64 pwd command, 8–28 quit command, 8–7 recv command, 8-40, 8-42, 8-62 *rename* command, 8–31, 8–72 *rmdir* command, 8–30 security, 8-5 to 8-6 send command, 8-49, 8-53, 8-64 status command, 8-74 status display, 8-74 type command, 8–34 user command, 8–78 to 8–79 verbose command, 8–3 verbose mode, 8-2, 8-83 wild card character use within, 8–12

G

gateway, Glossary 4 get command, ftp, 8–40, 8–42, 8–62 getcwd, A–6 gethostbyaddr, 11–21, 11–133 gethostbyent, 11–133

gethostbyname, 11–95, 11–133 gethostent, 11-13, 11-95 getnetbyaddr, 11–133 getnetbyent, 11–133 getnetbyname, 11–133 getpeerbyname, 11-131 getprotobyent, 11–132 getprotobyname, 11-16, 11-98, 11 - 132getprotobynumber, 11–132 getservbyent, 11-132 getservbyname, 11–15, 11–96, 11 - 132getservbyport, 11-132 getservent, 11–97 getsockbyname, 11–129 getsockname, 11-30 getsockopt, 11–72, 11–131 getwd, A-6 glob command, ftp, 8–12 globbing, 8–12, 8–74, 8–83 globbing behavior of commands in *ftp*, 8–13 to 8–14 guest ftp account, 8–81 guest *ftp* account, login to, 8–82

Η

hash command, ftp, 8–36
hash sign for ftp file transfer progress, 8–36, 8–74
help command, ftp, 8–9
help, telnet, 6–14
home directory, 7–8
host, Glossary 4, 2–2
host alias, 6–7
host internet address, 6–7
host load, 4–1 host name, 2–7, 4–1, 5–1, 6–7 host name, official, Glossary 5 host status, 4–1 host, remote, Glossary 5 *hosts* file, Glossary 5, 2–7, 6–7 *hosts.equiv* file, 2–7, 7–7, 9–3, 9–16, 10–2 HP-UX operating system, 1–6 *htonl*, A–2, 11–132 *htons*, A–2, 11–132

I

idle user, 4-2, 5-1 index, A-6 inet, 11-95 inet addr, 11-132 inet lnaof, 11-132 inet makeaddr, 11-132 inet_netof, 11-132 inet network, 11-132 inetd, 11-124 initiator, 9–2 input state, 6-2, 6-9, 6-13 installation, 1-6 interactive mode, 8-24, 8-26, 8-37, 8-44, 8-47, 8-57, 8-60, 8-69, 8-71, 8-74, 8-83 interface card, LAN, Glossary 4 International Standards Organization, Glossary 4 Internet, Glossary 4 internet address, Glossary 4, 2-7, 11-7 internet daemon, 11-124 Internet domain, Glossary 2 internetwork, Glossary 4

Internetwork communication socket address, Glossary 6 Interprocess communication, Glossary 4, 2–8, 11–1 accepting a connection, 11-20, 11 - 52adding server process to the Internet daemon, 11–124 address, Glossary 1 address conversion, A-2 address conversion call, 11-30 address family, 11-7, 11-12, 11 - 94addressing domain, 11-12, 11-47 advanced topics for stream sockets, 11-71 AF INET, 11-12 AF_UNIX, 11-9 association, Glossary 2, 11-8 binding, Glossary 2, 11-3, 11-8 binding a socket address to server process's, 11-17, 11-49 binding socket addresses to datagram sockets, 11–99 BSD IPC, A-1, A-5 to A-7 BSD IPC connections, 11-44 channel, Glossary 2, 11-7 client, Glossary 2, 11-4 client-server model, 11-4 closing a socket, 11–29, 11–61, 11 - 106communication domain, Glossary 2, 11–7 compile option, A-5 connection, Glossary 2 creating a socket, 11-16, 11-22, 11-47, 11-54 creating sockets, 11-98

datagram sockets, Glossary 3, 11-9, 11 - 91declaring socket address variables, 11-11, 11-46, 11-94 errno values, A-3 example using stream sockets, 11 - 30examples using datagram sockets, 11-107 FIONREAD, A-2 FIOSBNIO, 11–83 flag Options, 11-28, 11-60, 11-105 frame, Glossary 5 getting and setting socket options, 11 - 72getting the port address for the desired server, 11–96 getting the remote host's Internet address, 11-13 getting the remote host's network address, 11-95 graceful close, 11-32 graceful disconnect, 11-78 hard close, 11-78 I/O multiplexing with select, 11–118 INADDR_ANY, 11–97 incoming connection requests, 11 - 129internet address, 11-7, 11-11, 11-91 ioctl, A-2, A-5, 11-88 IPC connections, 11–2, 11–10 IPC system calls, 11–129 IPC using datagram sockets, 11–91 library calls, A-1 to A-2, A-6, 11-132 library equivalencies, A-6 library routines, 11–6 LINGER options, 11-29 *listen*'s backlog parameter, A-3

message, 0-5, 11-7 MSG OOB, 11–27, 11–87 MSG_PEEK, 11-27, 11-105, 11 - 116nonblocking I/O, 11–24, 11–119 nondestructive read, 11-28 other system calls, 11–131 out of band data, 11-28, 11-86 packet, Glossary 5, 11-7 pathname, 11-46 peer, Glossary 5, 11-7 pending connections, A-3 port, 11-8 port address, 11-11, 11-93 portability issues, A-1 preparing address variables, 11-11, 11-46, 11-93 preview an incoming message, 11 - 105preview incoming data, 11-28 programming hints, 11-121 protocols, 11-9 pty location, A-5 receiving data, 11-27, 11-59 receiving messages, 11-103 requesting a connection, 11-23, 11-55 reserved port addresses, 11-122 send size limit, A-5 sending and receiving data, 11-25, 11-57 sending and receiving messages, 11 - 101sending and receiving out of band data, 11-86 sending data, 11-26, 11-58 sending messages, 11–101 server, 11-4

setting the server up to wait for connection, 11-19, 11-51 signal calls, A-7 SIOCATMARK, 11-88 SO DEBUG, 11–71 SO_DONTLINGER, 11-71, 11-78 SO DONTROUTE, 11-71, 11-76 SO KEEPALIVE, 11–71, 11–75 SO LINGER, 11-77, 11-85 SO RCVBUF, 11-71, 11-77 SO_REUSEADDR, 11-71, 11-74 SO_SNDBUF, 11-71, 11-76 sockaddr, 11-12, 11-47, 11-94 sockaddr_in, 11-12, 11-47, 11-94 socket, Glossary 6 socket address, 11-8, 11-11, 11-46 socket descriptor, Glossary 6, 11–3, 11-9, 11-16, 11-48 sockets, 11-1 specifying a default socket address, 11-116 sprintf return value, A-7 stream sockets, Glossary 6, 11-9 summary tables for system and Library calls, 11-129 synchronous I/O Multiplexing with select, 11–79 TCP, 11–9 troubleshooting, 11-121 UDP, 11–9 unsupported IPC features, A-4 using a wildcard local address, 11–15, 11–97 using broadcast addresses, 11 - 121using diagnostic utilities as troubleshooting, 11-123 using read/write to make stream sockets trans, 11-86

using shutdown, 11-84 wildcard address, 11-11 wildcard addressing, 11–15, 11-97 writing the client process, 11-22, 11-54 writing the server and client processes, 11-98 writing the server process, 11-16, 11-47 interprocessing communication addressing domain, 11-94 ioctl, 11-81, 11-83, 11-119, 11 - 131IPC, Glossary 4, 2–8, 11–1 IPC connections, 11–2, 11–10, 11–16, 11–44, 11–48 ISO, Glossary 4

L

lcd command, *ftp*, 8–15 library calls, A–1 to A–2, A–6 library functions, 2–7 LINGER, 11–29 link-level address, Glossary 4 *listen*, A–3, 11–10, 11–19, 11–45, 11–51, 11–129 local, 2–2 login name, 2–7, 5–1 login, remote, 6–1 losing a TCP connection, A–4 *ls* command, *ftp*, 8–17, 8–20

Μ

mail, 1-2 mail destination local file, 3–2 local user, 3-5 remote user, 3-5 mail errors, 3-5 mail message aliasing, 3-1 mail message forwarding, 3-1 mail program, 3-2 mail transaction transcript, 3-5 mailq, 3-6 mailstats, 3-6 mailx, 3-2mdelete command, ftp, 8-69, 8-71 mdir command, ftp, 8-23, 8-26 memcmp, A-6 memcpy, A-6 memset, A-6 message, Glossary 5 metacharacter expansion in ftp, 8-13 to 8-14 metacharacters, 8-12, 9-29, 10-12 to 10-13 mget command, ftp, 8-44, 8-47 mkdir command, ftp, 8–29 mls command, ftp, 8-23, 8-26 mput command, *ftp*, 8–57, 8–60 MSG OOB, 11–27 MSG PEEK, 11-27, 11-59

Ν

netstat, 11–123 network alias, 2–7 network event logging, 11–123

Network Information Center, Glossary 4 network name, 2-7 network number, 2–7 Network Services, 1–6 network tracing, 11-123 networking, 1-6 networks file, 2-7 nftdaemon, 11–122 node, Glossary 5 node manager, Glossary 5 nonblocking I/O, 11-24, 11-83, 11 - 119NS, 1-6 ntohl, 11–132 ntohs, 11-30, 11-132 ntonl, A-2 ntons, A-2

0

O_NDELAY, 11–84, 11–120, 11–131 open command ftp, 8–4 telnet, 6–7, 6–11 to 6–12 out of band data, 11–86

Ρ

packet, Glossary 5, 11–7 password, 2–7 pathname, 11–49 peer, Glossary 5, 11–7 permission for local login from remote, 7–24 *perror*, 11–121 *ping*, 11–123

port, Glossary 5, 11–8 port address, 11-11, 11-14, 11-17, 11-93, 11-99 port number, 2-8 port, reserved, Glossary 5 praliases, 3-6 producer, 9-2 profile file, 7–14 prompt command, ftp, 8-37 protocol, Glossary 5 protocol alias, 2-8 protocol name, 2-8 protocol number, 2–8 protocols file, 2-8 pty, A-5 public ftp account, 8-81 public ftp account, login to, 8-82 public *ftp* directory structure, 8–81 put command, ftp, 8-49, 8-53, 8-64 pwd command, ftp, 8-28

Q

quit command ftp, 8–7 telnet, 6–13

R

rcp, 2–6, 9–1 -r option, 9–10, 9–14, 9–18, 9–22, 9–26, 9–29 combination copies local and remote to local, 9–22 local and remote to remote, 9–26 local to remote, 9–10 remote to local, 9–14

remote to remote, 9–18 copy as someone else, 9-30 directory copies local and remote to local, 9–22 local and remote to remote, 9 - 26local to remote, 9–10 remote to local, 9-14 remote to remote, 9-18 directory, remote working, 9-6, 9 - 30errors, 9–7 file attributes, effect on, 9-28 file copies local and remote to local, 9-20 local and remote to remote, 9 - 24local to remote, 9-8 remote to local, 9-12 remote to remote, 9-16 link to file, treatment of, 9-8, 9 - 12metacharacters, use of, 9-29 permission for local copy from remote, 9-31 result of copy with, 9-7 sources, allowed copy, 9-6 special files, treatment of, 9-6 wild card characters, use of, 9 - 29read, 11-9, 11-25, 11-57, 11-83, 11-105, 11-116, 11-120, 11 - 131readv, A-4 recv, A-2, A-4, 11-83, 11-105, 11-117, 11-119, 11-129, 11-131 recv command, ftp, 8-40, 8-42, 8 - 62

recvfrom, 11-7, 11-10, 11-25, 11-45, 11-84, 11-92, 11-99, 11-104, 11-116, 11-119, 11-129, 11-131 recvmsg, A-4 references, 1-6 remote, 2-2 remote host, Glossary 5 remote login, 6-1 remsh, 2-6, 7-26, 10-1 **-n** option, 10–15 command s execution as someone else, 10–7 execution of multiple remote, 10-10 execution of remote, 10-5 execution problems, 10-2, 10-14 process attributes of remote, 10-10 search paths, 10-6 use of interactive, 10-2 without input, 10–14 to 10–15 directory, remote working, 10-5, 10-7 hangup signal, treatment of, 10 - 5interrupt signal, treatment of, 10-5 metacharacters, use of, 10-12 to 10 - 13permission for local command execution, 10-8 quit signal, treatment of, 10–5 shorthand syntax, 10-17 signals, treatment of, 10-5 stdin, 10–14 stdin, stdout, and stderr, 10-12 to 10-13 terminate signal, treatment of, 10-5 rename command, ftp, 8-31, 8-72 reserved port, Glossary 5

return key behavior in telnet, 6 - 10rexec, 2-6 rfadaemon, 11-122 rindex, A-6 rlbdaemon, 11-122 rlogin, 2-5, 7-1, 10-2, 10-6 -7 option, 7–10, 7–12, 7–22 -e option, 7-10, 7-12, 7-22 -l option, 7–12, 7–22 automatic login, 7-7, 7-10 character size, 7-4, 7-10, 7-12 to 7–13 conditions requiring seven-bit characters, 7-6 escape character, 7-1 to 7-2, 7–10, 7–12, 7–15 to 7–16, 7–18 to 7–19 exit from (logout), 7–15 local work within, 7-16 login as someone else, 7-22 manual login, 7–12 requirements for sending eight-bit characters, 7-4 shorthand syntax, 7-26 rmail, 3–5 rmdir command, ftp, 8-30 rsh, 10–1 *ruptime*, 2–4, 4–1 -a option, 4-2, 4-4 to 4-11 -l option, 4–10 to 4–11 **-r** option, 4–5, 4–7, 4–9, 4–11 -t option, 4-6 to 4-7 **-u** option, 4–8 to 4–9 rwho, 2-5, 5-1 -a option, 5-2, 5-4

S

search path, command, 7-26, 10-17 select, 11-19, 11-51, 11-79 to 11-81, 11-118, 11-131 send, 11-10, 11-25, 11-45, 11-57, 11-83, 11-101, 11-116, 11-120, 11 - 129send command, ftp, 8-49, 8-53, 8-64 sendmail, Glossary 6, 2-4, 3-1 configuration file, 3-3 executing, 3-2 message, Glossary 5 message collection, 3–3 message routing, 3-3 production system, 3-3 sendmsg, A-4 sendto, 11-84, 11-92, 11-102, 11-116, 11–119, 11–129 server, Glossary 5 service, Glossary 5 service name, 2-8 services file, 2-8 sethostent, 11-133 setnetent, 11-133 setprotoent, 11–132 setservent, 11-132 setsockopt, 11-73, 11-75, 11 - 130shell escape ftp, 8–10 rlogin, 7-16 telnet, 6-17 shutdown, A-2, 11-30, 11-85, 11-129 SIGCHLD, A-7 SIGCLD, A-7 SIGIO signal, 11-81, 11-118 signal, A-7, 11-86 sigvector, A-7

Simple Mail Transfer Protocol, Glossary 6, 2–4 SIOCSPGRP, 11-81 SMTP, Glossary 6, 2-4 SMTP delivery module, 3-5 SO_DONTLINGER, 11-78 SO KEEPALIVE, 11-75 SO LINGER, 11-71, 11-77 SO RCVBUF, 11-77 SO REUSEADDR, 11-75 SO_SNDBUF, 11-76 SOCK DGRAM, 11-82 SOCK_STREAM, 11-82 socket, Glossary 6, 11-10, 11-45 socket address, Glossary 6, 11 - 93socket descriptor, Glossary 6, 11 - 3sockets, 11-3 special file, 9-6 sprintf, A-7 status ftp, 8–74 telnet, 6-20 status command ftp, 8–74 telnet, 6-20 status display in *ftp*, 8–74 strchr, A-6 stream sockets, Glossary 6, 2-8, 11-27, 11-59, 11-71, 11-129 strrchr, A-6

Т

TCP, Glossary 6, 2–8, 11–9 telnet, Glossary 6, 2-5, 6-1 ! command, 6-17 ? command, 6-14 *close* command, 6–12, 6–24 command descriptions, 6–14 to 6–15 command execution, 6–9 command list, 6-14 to 6-15 command state, 6–13 connection to remote host, 6–7, 6-20, 6-24 crmod command, 6–10 disconnection from remote host, 6-11 to 6-12 escape character, 6-2 to 6-4, 6-6, 6-9, 6-20 escape command, 6-4 execution, 6-2 exit from, 6-11, 6-13 help, 6–14 input state, 6–9, 6–13 local work within, 6-17 open command, 6-7, 6-11 to 6-12 quit command, 6-13 state, 6-2 status, 6-20 status command, 6-20 TELNET protocol, Glossary 6 terminal configuration characters, 6-5, 7-2 terminal line, 5-1 terms, 1-1 tilde, 7–2, 7–10, 7–12 Transmission Control Protocol, Glossary 6, 2-8 troubleshooting, 1-6 type command, ftp, 8-34

U

UCB, 2-1, 2-3 UDP, Glossary 3, Glossary 8, 11 - 9University of California at Berkeley, 2–1, 2–3 **UNIX** Domain address, Glossary 8 protocol, Glossary 8 user command, ftp, 8–78 to 8-79 User Datagram Protocol, Glossary 3, Glossary 8 user status, 4-1 utmp, A-6 uupath, 3–6 uux, 3–5

V

verbose command, ftp, 8–3 verbose mode, 8–2, 8–83 virtual terminal, 6–1 virtual terminal protocol, Glossary 8

W

wild card characters, 8–12, 9–29 wildcard address, 11–11, 11–15, 11–30, 11–93, 11–99, 11–120 write, 11–9, 11–25, 11–57, 11–83, 11–99, 11–116, 11–120, 11–131 writev, A–4

Reader Comment Card

HP 9000 Series 300 Using ARPA Services 50952-90001, E0189

We welcome your evaluation of this manual. Your comments and suggestions will help us improve our publications. Please tear this card out and mail it in. Use and attach additional pages if necessary.

Please circle the following Yes or No:

 Is this manual well organized? 	Yes	No
• Is the information technically accurate?	Yes	No
• Are instructions complete?	Yes	No
• Are concepts and wording easy to understand?	Yes	No
• Are examples and pictures helpful?	Yes	No
• Are there enough examples and pictures?	Yes	No

Comments:

N	2	m	le	•
1.4	u	**	10	•

Title:

Company:_____

Address:

City & State:



Zip:

BUSINESS REPLY MAIL

FIRST CLASS PERMIT NO. 37 LOVELAND, CO

POSTAGE WILL BE PAID BY ADDRESSEE

Hewlett-Packard Company Colorado Networks Division 3404 East Harmony Road Fort Collins, CO 80525

ATTN: User Information Development Department

Fold Here

NO POSTAGE NECESSARY IF MAILED IN THE UNITED STATES



Printed in U.S.A. 50952-90001 E0189

