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IBM 4341 Processors Model Groups 9, 10, 11, and 12 Functional Characteristics



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Third Edition (October 1982)

This major revision obsoletes GA24-3797-1. This edition adds information for the 4341 Processor Model Groups 9 and 12. Technical changes and additions to the text and illustrations are indicated by a vertical line to the left of the change.

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Preface

This reference publication is for system analysts and programmers who require information about processor features, input/output characteristics, timings, machine instructions, and the functions of integrated I/O devices.

This document is presented in two sections. The main or front section describes the 4341 processors in general and contains those features and functions that are common to all models.

The second section is a set of appendixes; each appendix contains those features and characteristics of a specific model of the 4341. Refer to the Contents for the specific areas covered in each section.

The reader is assumed to have a working knowledge of the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode*, Order No. GA22-7070, and the *IBM System/370 Principles of Operation*, GA22-7000, and to have had programming experience with System/360, System/370, or other 4300 processors.

The main chapters cover:

- Introduction to the IBM 4341 Processors
- System Structure
- Input/Output Channel Characteristics
- Display Console
- Facility Descriptions.

The appendixes cover:

- Storage Characteristics
- Channel Configurations
- Instruction Timings
- IBM 4341 Processor Complex Configurators.
- Prerequisite Publications
- IBM 4300 Processors Principles of Operation for ECPS:VSE Mode, GA22-7070

- IBM System/370 Principles of Operation, GA22-7000
- IBM 4300 Processors Summary and Input/Output & Data Communications Configurator, GA33-1523.

Associated Publications

- IBM 4341 Processor Model Groups 1, 9, and 10 Channel Characteristics, GA24-3671
- IBM 4341 Processor Model Groups 2, 11, and 12 Channel Characteristics, GA24-3780
- IBM 4341 Processors Operator's Guide, GA24-3669
- IBM 4300 Processors Installation Manual-Physical Planning, GA24-3667
- IBM Input/Output Equipment Installation Manual-Physical Planning for System/360, System/370, and 4300 Processors, GA22-7064
- IBM 3270 Information Display System Component Description, GA27-2749
- Introduction to Programming the IBM 3270, GC27-6999
- IBM Disk Pack and Cartridge Handling Procedures, GA26-5756
- IBM Diskette–General Information Manual, GA21-9182
- *IBM 3268-2 Printer Planning and Site Preparation Guide*, GA24-3266
- IBM 3270 Information Display System Color and Programmed Symbols, GA33-3056
- Multiply Add Facility, GA22-7082.
- An Introduction to IBM 4341 Processor Problem Analysis, SA24-3938
- IBM 4341 Processor Problem Analysis Guide, SA24-3925.

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Figure 1. IBM 4341 Processor with Typical I/O Configuration

IBM 4341 Processors Functional Characteristics

This publication is intended as a reference for users of the 4341 processors; only items that are unique to the 4341 are discussed in detail. Effective use of this manual requires a comprehensive understanding of the information in the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode*, GA22-7070, as well as the *IBM System/370 Principles of Operation*, GA22-7000.

The 4341 Processor (Figure 1) is a high-availability data processing system that provides the reliability, performance, and convenience demanded by both business and scientific users. The 4341 Processor is compatible with other 4300 processors, and can run under current program operating systems.

Note: For specific Model Group highlights, refer to the appendix for that Model Group.

Highlights

The 4341 offers Virtual Storage, System Control Program (SCP) support, and System/370 compatibility, implemented by using Large Scale Integrated technology and large processor storage.

The 4341 Complex consists of the 4341 and the IBM 3278 Model 2A Display Console or the IBM 3279-2C Color Display Console. The processor provides arithmetic, logic and control functions, storage, channels, and diskette drive.

Other significant 4341 characteristics are:

- Ease of installation, with minimum disturbance of existing input/output configuration.
- DOS/VSE, OS/VS1, MVS/SP, and VM/370 program support.
- A standard high-speed buffer storage.
- Processor storage (Model Group dependent). (Some of this storage is required by the system, as described under "System Storage Requirements.")
- Improved reliability, availability, and serviceability (RAS), including instruction retry, Error Checking and Correction (ECC) to provide single-bit error correction and double-bit error detection in processor storage. Error recording by the hardware itself and the Remote Support Facility (RSF) for remote maintenance is also provided.
- Byte Multiplexer and Block Multiplexer Channels -

Byte-multiplexer channels have a single data path that can be fully utilized by one I/O device (burst mode), or shared by many I/O devices (multiplexer mode). Burst mode is defined as operation in which the device and the channel remain connected for a relatively long period of time in terms of system operation. Multiplexer mode is defined as byte-interleaved operation in which the channel and any one device remain connected for a relatively short period of time, typically long enough to transfer one byte or a small number of bytes.

Block-multiplexer channels can operate in either block-multiplexer mode or selector mode. The block-multiplexer channel in block-multiplexer mode allows concurrent operation of many devices. A block-multiplexer channel in selector mode, executing a command-chained channel program, is busy while the channel program is in operation, whether data transfer is occurring or not.

The channels may overlap operation from the Instruction Execution function of the 4341.

The block-multiplexer channels *appear* as selector channels to I/O devices that do not block multiplex.

An optional Channel-to-Channel Adapter is also available.

For further information on channels, refer to the Channel Characteristics manual noted in the preface. For further information on a particular processor's channel capability, refer to the appendixes in this manual.

- Two modes of operation selectable at Initial Microcode Load (IML) time:
 - Extended Control Program Support VSE (ECPS:VSE) Mode – allows operation of an appropriately generated DOS/VSE system with enhanced performance.
 - System/370 Mode allows operation of any program written for the System/360 or System/370 that does not violate the exceptions noted in the "Compatibility" section below. For those System Control Programs (SCPs) that contain 4341 support, see "Programming Support."

In this mode, three mutually exclusive options are available:

ECPS:VS1 Assist – provides a hardware assist that reduces the processor time needed to execute certain frequently used supervisor functions in VS1, Release 7 or later. In this mode, other supported SCPs operate but without enhanced performance.

ECPS:VM Assist – reduces the processor time needed to execute certain frequently used supervisor functions in VM/370, Release 6 or later. In this mode, other supported SCPs operate but without enhanced performance.

ECPS:MVS – allows the 4341 processor to operate with compatibility support of the System/370 Extended Facility while running in S/370 mode. This selection provides the System/370 facilities that are a prerequisite for operation in the MVS/SP environment.

Note: ECPS:VM and ECPS:MVS can be run concurrently on a Model Group 12 processor that has the control storage expansion feature.

Engineering Scientific Assist - this assist improves the performance of certain mathematical computations. The assist consists of one instruction, MULTIPLY AND ADD, that may be used in either System/370 or in ECPS:VSE mode.

The MULTIPLY AND ADD instruction performs a combination of vector multiplication and addition operations which may replace the inner loop of common matrix computations and reduce processor calculation time.

The IBM 3278 Model 2A Display Console or 3279-2C Color Display Console required for interaction with the 4341 for both operation and maintenance. The Operator Console Keyboard with its Operator Control Panel (OCP) is used for turning power on and off, for Initial Microcode Load (IML), and for starting and stopping processor operations.

Both the Display mode and the Printer-Keyboard Emulation mode are supported. In Display mode, the keyboard is used for input, and the display is used for the output of up to 20 lines of up to 80 characters each.

In Printer-Keyboard mode, the keyboard is used for input. The display and a recommended 3268 Printer Model 2 or 3287 Printer Model 1, 2, 1C, or 2C are used for output. The 3278-2A Display Console and the 3287 Printer appear to the system as a console printer-keyboard. This allows using an operating system that has been generated for a System/360 with a 1052 Printer-Keyboard or a System/370 with a 3210 or 3215 Console Printer/Keyboard. An optional alternate console (with one display/keyboard and one 3287 Printer) can also be configured.

The console also provides for *normal* versus *instruction step* processing, for address compare stopping, for altering certain registers and storage areas, and for displaying processor status.

For maintenance and service support, the console can display and store the status of the 4341 complex and other pertinent servicing information. It also allows for the use of diagnostic tools.

Up to three optional 3278 Model 2A Display Consoles or 3279-2C Color Display Consoles, or 3268-2 or 3287 Model 1 or 2 Printers can be configured (for a total of four). The optional printer has a separate address in display mode and requires Multiple Console Support (MCS).

Note: The procedures for configuring 3268 Model 2 or 3287 Printers depend upon the operating system being used. For OS/VS1, for example, the 3287 is supported by specifying either a 3286 or 3210 Printer.

The 3278-2A and 3279-2C features other than those basic to the primary display console are not supported.

- The Support Processor for automatic analysis of failure symptoms. The result of this "self diagnosis" is a processor-generated *reference code* that contains information to guide the service representative to the failing unit. This reference code is logged on the system diskette, and displayed to alert the operator of possible machine malfunction.
- The System Diskette Drive for both IML of microcode and recording of errors for later diagnosis. The removable diskettes provide all of the microcode required for initializing basic processor features (and optional features, when ordered), as well as diagnostics for the service representative.
- The Remote Support Facility (RSF) for use (when installed and authorized by the customer) to enhance hardware maintenance.
- The Remote Operator Console Facility (ROCF) for assisting the operation of a 4341 in a Distributed Data Processing (DDP) environment.

Modes of Operation

The 4341 executes all of the processing and input/output functions described in the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode,* as well as those functions described in the *IBM System/370 Principles of Operation.* The major difference between the mutually exclusive System/370 and ECPS:VSE modes is in the handling of virtual addresses.

- ECPS:VSE mode uses internal address translation for both processor and channel addresses. All storage addresses are virtual addresses.
- System/370 mode uses segment and page tables for processor Dynamic Address Translation (DAT). Channel addresses are real addresses that are translated by the System Control Program (assisted by the Channel Indirect Data Addressing facility).

Programming Support

Programming support for the 4341 in ECPS:VSE mode is provided by DOS/VSE. In System/370 mode, programming support is provided by DOS/VSE, OS/VS1, MVS/SP, and VM/370.

Note: MVS/SP is not supported on the one-megabyte Model Group 9 processor. It is supported, but not recommended on a two-megabyte processor.

Brief descriptions of these program support packages (and references to the publications that describe them in detail) are available from your IBM representative. Additional information about 4341 processing and input/output functions, and Basic Control (BC) and Extended Control (EC) modes, is presented in the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode* or, for System/370 mode, in the *IBM System/370 Principles of Operation*.

Remote Support Facility (RSF)

This facility (when installed and with customer authorization) allows remote control of the 4341 from an IBM RETAIN/370 site, and allows the on-site service representative to access the IBM RETAIN data bank for the latest service aids and information.

While in this mode, the IBM Remote Support personnel can perform online diagnosis as though he were at the customer's site. Logout data stored on the system diskette drive can be saved in RETAIN during the Data Link (DL) operation for later offline analysis. Microcode patches may also be applied remotely.

The remote connection is via a customer-supplied Data Access Arrangement (DAA). For connection information, refer to the *IBM 4300 Processors Installation Manual–Physical Planning*, Order No. GA24-3667. Remote console operation from any IBM RETAIN terminal is through the Data Link facility of RETAIN.

The 3278-2A is used to monitor RSF data transmission. The DISC key can be used to terminate data transmission at any time.

In customer installations where the IBM RETAIN facilities cannot be used, remote control is possible via an IBM 3275 Display Station (using a dial-up, 1200-baud, bisynchronous, switched line).

Remote Operator Console Facility (ROCF)

The Remote Operator Console Facility (optional feature) is an extension of the Remote Support Facility (RSF). When installed and enabled, the Remote Operator Console Facility is active in the support processor when the 4341 is powered on. (When ROCF is active, the RSF functions described above are inactive.)

In a Distributed Data Processing (DDP) environment, Remote Operator Console Facility allows personnel at the host site to dial-up the remote 4341 and control the remote system from the host site. This control is accomplished through system operation functions such as IML/IPL, Reset, Restart, Compare/Trace, and Display/Alter.

To use the Remote Operator Console Facility, the remote 4341 system must be equipped with a customer supplied auto-answer modem and Data Access Arrangement (DAA). For details, refer to the IBM 4300 Processors Installation Manual- Physical Planning. Communication with the Remote Operator Console Facility is via an IBM 3275 Display Station or by programming support provided by IBM Program Products designed to provide remote console communication capability.

In Remote Operator Console Facility mode, the optional security keylock feature on the 4341 system console (if installed) allows the host site to control the remote 4341 without interference from unauthorized personnel at the remote site. A password verification function is also part of the Remote Operator Console Facility and serves to protect against unauthorized access to the remote 4341. If a higher level of data security is required, an external encryption device may be attached to the dial-up link.

After the remote 4341 system is successfully initialized, normal transfer of data and control information between the host and the remote system should be handled through a conventional communication network (e.g. 270X or 370X communication controllers).

Minimum Configuration for Hardware System Maintenance

The following minimum configuration is required for hardware maintenance. The individual System Control Programs (SCPs) have their own minimum requirements depending upon the SCP type and release level.

Minimum Configuration with Demountable Direct Access Storage

- 4341 Processor
- 3278-2A Display Console or 3279-2C Color Display Console
- Access to one of the following groups of devices:
 - 1 Card Image I/O device* and
 - 2 Direct Access devices** and
 - 1 Hard-Copy Output device,

or:

- 1 Card Image I/O device* and
- 1 Direct Access device** and
- 2 Magnetic Tape devices*** and
- 1 Hard-Copy Output device,

or:

- 1 Card Image I/O device* and
- 3 Magnetic Tape devices*** and
- 1 Hard-Copy Output device.
- * Card Image is defined as:
 - Any supported Card Reader, or
 - An addressable diskette input/output unit (such as a 3540) and key-to-diskette capability, or
 - A magnetic tape drive and provisions for entering card-image formatted records onto magnetic tape, or
 - Capability provided by the customer through his operating system facilities to create card-image format on either tape or diskette. The customer must supply an operator to key the card images at the direction of the service representative.

- ** Must be demountable Direct Access Storage Device (DASD).
- *** If 2400 Series, seven-track, magnetic tapes are used, Data Conversion features (No. 3228 and 3236) must be installed on the 2803 or 2804 Tape Control unit.

Minimum Configuration with Nondemountable Direct Access Storage

For configurations with nonremovable direct access storage devices (DASD), the following devices constitute the minimum configuration for hardware maintenance, if the first forty cylinders on a nonremovable drive (other than the System Residence drive) are made available for the generation and maintenance of service programs. This space must be allocated for initial installation, for modifications to the configuration, and for the application of maintenance facility updates.

- IBM 4341 Processor
- IBM 3278-2A Display Console or 3279-2C Color Display Console
- Card Image I/O Device (See * on previous page)
- Nonremovable DASD: IBM 3350 — The first 40 cylinders of a drive dedicated when required.

Note: After use of the 3350 by the service representative, this drive may need to be reformatted by the customer for customer use.

- Magnetic Tape Device
- Hard-Copy Output Device

Additional Requirements for Installation and Operational Maintainability

In all configurations, each processor must use IBM programs (or equivalent) that provide for error recording, with elements for handling machine check interruptions and for recording status of the processor when a failure is detected. Routines for error recording are contained in some releases of DOS/VSE, OS/VS1, MVS/SP, and VM/370. IBM's ability to service configurations that do not meet the above requirements may be impaired, and may affect system availability.

To further enhance maintainability and availability, Remote Support Facility (RSF) should be provided.

System Residence and Storage Requirements

Optimum performance and maximum availability are

obtained when a disk-storage facility is provided. The DOS/VSE, VS1, MVS/SP, and VM/370 operating systems *require* a disk storage facility. These storage requirements are assumed to be attached through a block-multiplexer channel.

A portion of processor storage is required for dynamic tables. This reduces the amount of processor storage available for user programming.

Compatibility with System/360, System/370, and other 4300 Processors

An important difference between the System/370 and the 4300 processors, when operated in ECPS:VSE mode, is the concept of virtual storage being mapped to real storage under hardware and microcode control. *Real storage* is the amount of storage that is physically installed. The apparent storage (called *virtual storage*) can be any amount of storage that an application requires, up to 16,777,216 bytes.

Any program written for IBM System/370 can operate on the 4341 Processor in System/370 mode, if it:

- Is not time-dependent.
- Does not depend on system facilities (storage size, I/O equipment, optional features, etc.) being present when the facilities are not included in the configuration.
- Does not depend on system facilities (interruptions, operation codes, etc.) being absent when the facilities are included in the 4341.
- Does not depend on results or functions that are defined in the *Principles of Operation* to be unpredictable or model-dependent.

Any program written for the 4300 processors in ECPS:VSE mode operates on the 4341 Processor, if it follows the above rules.

Any program written for the System/360 will operate on the 4341 if it follows the above rules and does not depend on functions that differ between System/360 and System/370. The System/370 functions that differ from System/360 functions are described in an appendix of the *IBM System/370 Principles of Operation*.

For additional information about compatibility, see *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode*, GA22-7070.

An important aspect of compatibility is the disk data format. System/360 and System/370 uses the Count-Key-Data (CKD) architecture. The 4341 supports disk units with both the CKD format and Fixed-Block Architecture (FBA) formats. Existing disk volumes can be mapped onto system disk devices.

Data Representation

The 4341 is both character- and word-oriented. The basic addressable unit is an eight-bit byte (a character, two decimal digits, or eight bits). This allows efficient use of storage; and highly-effective input/output rates for decimal data, variable field lengths, broad and flexible code conversion, and decimal arithmetic. This data representation provides for 32-bit words and 16-bit half-words for fixed-point arithmetic, 32-bit words and 64-bit doublewords for floating-point arithmetic. Functions such as translate and edit can be implemented using this data representation.

Permanently Assigned Storage Areas

All byte locations of processor storage are available for programming functions except the permanently assigned processor storage areas (storage locations 0 through 511, as described in the *IBM System/370 Principles of Operation* and the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode*).

Control Storage

Reloadable control storage, not available to the user, accommodates the microcode for controlling standard and installed optional features. A control storage expansion feature is available on the Model Group 12.

High-Speed Buffer Storage

The high-speed buffer storage functions as a smaller and faster subset of processor storage to improve performance. This high-speed buffer is not part of user-addressable storage. The buffer size (capacity) varies with processor model.

Local Storage

This 1K-byte area contains the general registers, floating-point registers, channel unit control word (UCW) directory, etc.

Channel Buffer Storage

A data buffer of 256 bytes per channel facilitates data transfer operations between processor storage and the channel I/O interfaces.

Arithmetic and Logic Unit (ALU)

The ALU has an access width of eight bytes and provides for binary and decimal addition and subtraction operations, as well as for logical AND, OR, and EXCLUSIVE OR operations.

Support Processor

The support processor (SP) controls processor initialization (initial microcode loading), instruction retry, error analysis and logging, and supports the display console and related equipment.

Display Console

The IBM 3278 Model 2A Display Console or 3279-2C Color Display Console serves as the online input/output device for operator-system communications. The display console is used to enter or display program-system control parameters, responses to system messages, and the display or alteration of application data, general and floating-point registers, etc.

Operating the System

The 4341 is operated, monitored, and controlled from the display console. For detailed information about the basic display console, optional display consoles and printer configurations, and system operation, refer to "Display Console."

Operator Controls

The operator is made aware of processor status via the displays described below and can invoke the described functions from the keyboard.

Normally, the display/keyboard communicates with an application program. Twenty lines of the screen are reserved for this purpose. The bottom lines on the screen are reserved to show processor status (operating, wait, manual, test, and load, and the address and data fields) and the setting of the rate control, check control, and address compare controls. The functions that can be invoked by pressing a function key are:

Start Stop Mode Selection Change Display Copy Page External Interruption Cursor Controls (left, right, insert, delete, erase EOF, erase IPT) Request (PA1) Cancel (PA2)

When the Mode or Change Display keys are pressed, the operator is able to set controls or invoke functions via the keyboard, either with or without menu prompting, as follows:

Alter/Display Address Compare Interval Timer Enable/Disable Check Control Program Load Operation Rate Restart System Reset (Program or Clear) Store Status (in System/370 mode) or Save (in ECPS:VSE mode) Time-of-Day Clock Enable-Set

System and Control Status

The indicator fields displayed are:

Operating Indicator is on when the processor is not in stopped state.

Wait Indicator is displayed when the processor is in wait state.

Manual Indicator is on when the processor is in the stopped state.

Test Indicator is on when a control is not in its normal state or when a maintenance function is being performed.

Load Indicator is on during program load. It goes off when the new PSW is successfully loaded.

Save Indicator is turned on upon successful completion of Machine Save (in ECPS:VSE mode) or Store Status (in System/370 mode).

Addressing Indicator displays the address of the next

instruction to be executed.

V-ADDR indicates that the machine is in System/370 mode and that it is operating with virtual addresses.

R-ADDR indicates that the machine is in System/370 mode and that it is operating with real addresses.

ADDR indicates that the machine is in ECPS:VSE mode and that it is operating with virtual addresses.

Depending on the dynamic address translation (DAT) setting, either 'V' or 'R' can appear in System/370 mode.

Data Indicator displays the contents of the next address to be executed when the machine is in stopped state. While the machine is operating the display is blank.

Timer Indicator is displayed in both System/370 mode and ECPS:VSE mode and shows the setting of the Interval Timer control.

Rate Indicator is displayed when the Rate control is set to Instruction Step.

Check Control Indicator is displayed when the Check Control control is set to Stop After Log-Switch, Channel Log, Hardstop, No Retry, or Disable.

Address Compare Indicator:

Control is displayed when the Address Compare control is set to Stop, Trace Stop, or Trace Wrap.

Type displays the setting of the way in which an address compare is made: any reference, a data store reference, an I/O reference, or a reference to fetch an instruction.

Addr/data displays the address for the compare or the data contents of that address if Address Compare is on Data.

Check Stop Indicator is on when the machine is in check-stopped state after an error.

Manual Controls

System Reset:

Program clears all equipment check indications and sets the processor so that operation may be resumed. An I/O reset is also performed.

Clear issues Program Reset and clears processor storage, the general purpose registers, control registers, and the floating-point registers.

Note: These functions can be automatic if selected from the program load screen.

Start key causes the processor to enter the operating state.

Stop key causes the processor to enter the stopped state when the current processing step is completed and any pending interruptions are handled.

Address Compare: The processor enters the stopped state, or traces when an address set by the address compare function matches an address used in a storage reference. An address compare stop/trace can occur on any reference, a data store reference, an I/O reference, or a reference to fetch an instruction.

Data Compare: The processor enters the stopped state, or traces when data established by the data compare function for an address matches data in a storage reference at that address. A stop can occur on any reference, a data store reference, an I/O reference, or a reference to fetch an instruction.

Interval Timer: When the interval timer is activated, processor storage location 80 is decremented every 1/300 of a second. When the location goes negative, an interruption is generated. If Interval Timer is off, location 80 is available for other uses.

Check Control:

Normal: On an error, the processor may perform retry. If successful, system recovery is reported. If unsuccessful, instruction processing damage or system damage is reported.

Hardstop: No machine check handling is permitted, and no reference code is generated.

No Retry: Retry is disabled.

Disable: On a CPU or channel error, the processor ignores the error and attempts to continue processing.

Stop After Log - Switch: On a CPU or channel error, the processor performs a logout and check-stops, or performs a retry operation and, if unsuccessful, performs a logout and enters the check stop state.

Program Load: After an Initial Program reset or a Clear reset (a user choice), the processor loads 24 bytes from the load unit addressed, continues until the Initial Program Load (IPL) chaining sequence is complete, and starts the processor under control of the new PSW.

Operation Rate:

Normal operates the machine at normal speed when start control is activated.

Instruction Step performs one unit of execution, takes all pending allowed interruptions, and returns to the stopped state.

Restart stores the old PSW at location 8, fetches the new PSW from location 0, and enters the operating state.

Interrupt Key raises an interruption to the processor. The pending request is cleared when it causes an interruption or by a program reset.

Time-of-Day Enable/Disable allows the operator to change the status of the machine to allow, or disallow, Set Clock instructions (TOD Clock Enabled).

The delay from the time that the control is activated to the time the clock is enabled is approximately 25 ms. TOD secure selection delay is normally 3 seconds.

Basic Functions

Alter / Display (System)

Several facilities can be displayed, some of which can be altered. Any of the selections described below cause the processor to enter the manual state. A return to the operating state is required for processor operations to continue.

Display

Registers, processor storage, and certain other facilities can be displayed.

These facilities can be displayed in two ways. One way is to go through the mode selection frame to the alter/display frame and then use the menu to fill in the required information. The other way is to use *fast selection* that is, keying in a command indicating the function desired, including any required information, such as storage address.

Alter

Registers, current PSW, and processor storage can be altered in two ways: menu selection or fast selection (see "Display" above).

During menu selection, the facility is displayed as it exists in the processor. The operator can then change part or all of it by moving the cursor to the characters to be changed and entering the new data. After the change is entered, the display console shows the change. For fast selection, data is entered with the command line. The facility is only displayed after the change.

Facilities

The facilities displayed are:

- G General Registers
- C Control Registers
- F Floating-Point Registers
- P Current PSW
- K Storage Key (in System/370 mode), or Block Description (in ECPS:VSE mode)
- V Processor Storage-Virtual
- M Processor Storage-Real (available only in System/370 mode)
- T Trace Area
- U UCW/Device Directory

General Registers, Control Registers, and Floating-Point Registers can be both displayed and altered. For an alter operation, if the new data is included in the command, any byte of a register can be addressed. Several registers can be altered at a time. When the cursor is used, any hexadecimal digit can be changed.

Current PSW can be both displayed and altered. When displayed, the entire PSW is shown in hexadecimal. Certain fields in the PSW are also formatted.

Storage Key can be both displayed and altered. When altered, the entire key must be entered.

The display is in the same format that is required for the input. Certain fields in the Storage Key are also formatted.

Processor Storage-Virtual can be displayed or altered. A page ability is included. If the Dynamic Address Translation (DAT) bit is off (System/370 mode), the message ADDRESS NOT AVAILABLE is issued.

Processor Storage-Real/Machine functions like Processor Storage-Virtual.

Trace Area displays the addresses in the trace area.

UCW/Device Directory contains the configurations of all devices on the channels.

Page

The page function displays processor storage. The appropriate display routine fills in the next required address to the previous command. The two page keys allow for either page forward or backward. These keys can be used any time a facility, such as processor storage, cannot be fully displayed on one screen.

Mode

These keys and controls are described in more detail in "Display Console."

Mode Sel Key

When the MODE SEL key is pressed, the general selection frame shows on the screen, and the operator can select an additional frame.

Chg Dply Key

The CHG DPLY key changes the screen between display mode and manual control mode.

Copy Key

The Copy key is only operational when in manual control mode and a 3268-2 or 3287 Printer installed. Pressing the Copy key causes the screen image to be printed on the printer that is attached to the adapter.

This can be used to save displays, etc.

Checking the System

The 4341 can be checked by using the System Test programs that are available from IBM on magnetic tape or external disk storage. These programs test and report the condition of the processor and I/O devices made available.

Problem Analysis

The 4341 has an enhanced maintenance capability called Problem Analysis (PA). PA is an automatic problem determination routine. It is part of the basic microcode in the 4341 Processor and includes a separate Processing Unit Analysis Diskette (PUAD). PA is designed to help resolve problems faster and get jobs running again.

The operator can start PA from the General Selection screen when suspecting a problem. The *IBM 4341 Processor Problem Analysis Guide* (SA24-3925) provides the detailed operating instructions.

System Structure

This section describes the major 4341 components to provide an overview of machine and program interrelationships.

Data Flow

In Figure 2, the simplified data flow of the processor describes the 4341 operation. Microinstructions from the system diskette drive are loaded into control storage during the initial microcode load (IML) procedure. The microcode controls data transfer operations to or from both processor storage and control storage. Work storage is used as a source or destination for this data. During I/O operations, data is transferred between storage and the I/O interface via the channel data buffer. The dotted lines represent optional features.

For an understanding of logical and physical channel relationships, refer to "Input/Output Channel Characteristics."

Processor

The processor contains processor storage, control storage, the system control panel, and other facilities necessary to perform arithmetic functions and logical processing of data. The processor also contains the input/output channel circuitry for control of data transfers between the processor and I/O devices.

The manner in which the various data-flow elements interact depends upon the microcode loaded into control storage. The processor executes the instructions defined in *IBM 4300 Processors Principles* of Operation for ECPS:VSE Mode and the *IBM*



Figure 2. Simplified 4341 Data Flow

System/370 Principles of Operation, including input/output instructions and commands.

Storage Concept

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The storage concept (Figure 3) consists of real storage and virtual storage.



Figure 3. Storage Concept

Error Handling

The error handling function provides both automatic recovery from many hardware malfunctions and reporting by means of a machine check interruption to assist in program damage assessment and recovery.

System Retry Overview

The Retry function makes intermittent, processor detected, hardware failures logically transparent to processing. Retry is, in general, done on a machine instruction basis. The data in certain machine facilities is saved during instruction execution. To perform retry when a malfunction occurs, this data is restored in those machine facilities, and the instruction is re-executed.

Because operand addressability is not verified by pretesting, the retry mechanism can also be invoked when an operand crosses a storage page boundary and an access exception occurs. This has the effect of logically backing up the processor to the state preceding the instruction so that a program interruption can be taken as if the instruction had been nullified.

No retry is performed on malfunctions affecting channel operations (these result in channel control checks and/or machine checks). For both retriable and unretriable errors, the support processor performs an internal logout of hardware latches for analysis. Several logouts are retained on the system diskette.

Machine Check Handling

The 4341 error recovery facilities and machine check interruption procedures comply with the general definitions in the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode* and the *IBM System/370 Principles of Operation.*

The error handling function can perform any of the following:

- 1. Re-execute an instruction or interruption that failed because of an intermittent hardware malfunction
- 2. Post a machine check interruption for uncorrectable malfunctions and alert conditions
- 3. Enter check stop state when a malfunction makes it undesirable or impossible to continue operation.

If application-program errors occur, the operating system attempts to handle the exception and to provide any operator messages. Refer to the applicable programming publications for the operating system you use.

If a failure occurs within the processor or an I/O unit, provisions have been made to retry the failing operation. Facilities are incorporated to record any such failures. (This is in addition to any provisions made by the operating system for error retry and error logging.)

Instruction retry, limited channel logout, storage validation (error checking and correction–ECC) for processor storage, and other error-detection and error-handling provisions are standard.

Error Checking and Correction (ECC)

Error checking and correction on processor storage provides automatic, single-bit error detection and correction. ECC also detects all double-bit errors and most multiple-bit storage errors, but does not correct them. Parity checking is used to verify other data that is not contained in processor storage.

Channel Command Retry

Channel command retry is a control-unit initiated procedure between the channel and the control unit. (Not all control units have this capability.) No I/O interruption is required. The number of retries is device-dependent.

Internal Logout

Each machine check normally leads to an internal logout. Because the data flow area and microcode are used for the control of channels and integrated adapters as well as for processor control, the logout buffer contents can relate to either a machine check condition or a channel control check. Thus, the interruption may relate to either a processor error or a channel error.

For a hardware error other than corrected single-bit errors, the processor generates an internal logout as a preliminary action to error analysis.

Machine Check Interruption Bits

The machine check interruption code (MCIC) bits are described in the IBM 4300 Processors Principles of Operation for ECPS:VSE Mode and the IBM System/370 Principles of Operation.

Check Stops

The processor enters the check stop condition if any of the following situations occurs:

- Invalid microcode is executed.
- Certain severe hardware failures.
- A storage error in the system storage (aux) area which cannot be corrected.
- A malfunction occurs during IML/IPL.
- PSW bit 13 is set to 1, and an error occurs during a machine check interruption that is attempting to report system damage or instruction processing damage.
- PSW bit 13 is set to 0, and an error occurs that causes the system damage bit or the instruction processing damage bit to be set to 1 in the MCIC buffer.
- An error occurs, a successful retry was not performed, and the Stop After Log switch is active.

System Operation

Operation is initiated when the operator presses the Power On key. After the power-on sequence is completed, the microcode is loaded (under operator control) from the system diskette drive into the control storage of the processor. Normally, the operating system is then loaded from the system disk by a program-load command from the display console. The customer's application programs are read into processor storage from an appropriate input device.

Program execution follows the normal System/370 pattern. The processor executes the machine instructions in sequence under control of the current PSW and the mask bits in the control registers. The normal sequential execution of instructions can be changed by branching.

During ECPS:VSE mode operation, processor storage is always addressed by designating one of 16-megabyte locations. Locations that are *addressable* according to the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode* are automatically mapped to their corresponding physical storage locations.

Interruptions of six different classes (machine check, supervisor call, program, external, input/output, and restart) can occur during operation. An interruption involves storing the current PSW as the old PSW, storing a code identifying the cause of the interruption, and fetching a new PSW. Processing is resumed according to the new PSW.

Support Processor

The support processor (SP) maximizes the total system availability and provides for rapid isolation and repair of failures, whenever possible, by automating and simplifying failure diagnosis.

The support processor provides the services necessary for initialization, monitoring, and maintenance. Integrated adapters and a logic element communicate with console I/O devices and other elements of the processor to provide for microcode loading, messages, and routine checking facilities. The system diskette drive is used for loading microcode that controls processor and channel operations, and for residence areas for error log and analysis data.

The support processor is connected to the 4341 through the local channel adapter (LCA) and the support bus adapter (SBA).

The common communications adapter (CCA) and the remote support facility (RSF) included in the support processor are used by IBM personnel for maintenance.

System Diskette Drive

The system diskette drive is the initial microcode load (IML) device and provides, by means of removable diskettes, the microcode for initialization, system operations, diagnostics, etc. Data loaded into control storage is not accessible to the programmer.

Diskette Wear

Diskettes provide for ease of media handling and storage, etc. Note, however, that during recording and reading, the read/write head is in contact with the media causing wear over time. Care in storage, use, and handling can also affect diskette life. (See guidelines in *IBM Diskette–General Information Manual*, GA21-9182.)

Initial Microcode Loading (IML)

Initial microcode loading is executed in two phases: the support processor, and the main processor. When the appropriate controls on the console are activated, disk reading starts, and the load routine transfers microcode from the system diskette drive into the support processor. Then, optionally, the main processor microcode is loaded into control storage. When loading is completed, the system-reset microcode is executed. After system reset, the processor stops until an external action is taken, such as initial program load (IPL).

Input/Output Channel Characteristics

The 4341 I/O channels, with a few variations, are identical to those of System/370. When in ECPS:VSE mode, virtual addressing is used. When in System/370 mode, channel indirect data addressing can be used. The description of I/O channels that follows describes System/370 mode. For additional information about channel operations, refer to the IBM 4300 Processors Principles of Operation for ECPS:VSE Mode and the IBM System/370 Principles of Operation.

For information about I/O devices attaching to the 4341 Processor interface, refer to the *IBM 4300 Processors Summary and Input/Output & Data Communications configurator.*

This section covers the basic characteristics and defines the limitations of the I/O channels.

The following channel configurations are available:

• A byte-multiplexer channel, logical address 0.

The support processor is internally attached to this channel. The 3278-2A or 3279-2C (and any optional 3268-2 or 3287 printers or additional 3278-2A or 3279-2C devices) are internally connected to the support processor, and are addressed via channel 0.

Addresses 0F0 through 0FF are reserved for natively attached I/O devices and the support processor, and are not available for I/O devices attached to channel 0.

- At least two standard block-multiplexer channels, logical addresses 1 and 2.
- Optionally, a number of block-multiplexer channels, logical addresses 3, 4, 5. (Service personnel can connect an additional channel, usually channel 4, as a second byte-multiplexer channel, if so selected.)

For further information on a particular processor model's channel capability, refer to the appendix for that model.

The byte-multiplexer channels can operate in either byte mode or burst mode. This is determined by the characteristics of the device operating on the channel. When a device forces burst mode, no other device can operate with that channel until the burst-mode operation is completed. In byte mode, the single data path of the byte-multiplexer channel can be shared by a number of low-speed I/O devices operating simultaneously. The channel multiplexes data to or from these devices (one device at a time) in groups of bytes as required by the I/O device being serviced.

The block-multiplexer channel is optimized for relatively high-speed burst operations and can multiplex complete blocks of data. This channel is particularly suited to buffered or cyclic devices with high data rates (such as disk-storage devices). The multiplexing facility of the block-multiplexer channel allows the interleaved execution of several channel programs by the same channel.

Note: The block-multiplexer channels *appear* as selector channels to I/O devices that do not block multiplex.

Block-multiplexing control (defined in *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode* and the *IBM System/370 Principles of Operation*) allows operation of the block-multiplexer channel as a selector channel.

The byte-multiplexer and block-multiplexer channels operate from the same I/O instruction and command formats used for System/370.

Standard Input/Output Interface

The standard System/370 I/O interface is used to connect the channels to I/O devices or control units.

High-Speed Transfer

This facility enables data transfer to take place (on block-multiplexer channels) faster than the data rates that can be obtained with service-in and service-out alone. Because of the higher-speed data transfer, this facility can also be used to permit locating a control unit at a greater distance from the channel than would otherwise be possible. This facility includes two additional tag lines: data-in and data-out.

Data-in can be alternated with service-in tag line to enable transfer of data at a higher rate than is possible if service-in alone is used. In this case, data-out is alternated with service-out as the response to data-in. Data-out is the response to data-in as service-out is the response to service-in.

Data Stream Mode

Data Stream mode is an extension of the hardware protocol for I/O data transfer which permits faster data transfer rates and increased channel to control unit cable length. Data Stream mode can operate on any block-multiplexer channel up to the maximum data rate specified in the appendixes for each model group.

I/O Error Alert

An additional selection (formerly tag) line called disconnect-in provides (on block-multiplexer and byte-multiplexer channels) control units with the ability to alert the processor of a malfunction that prevents the control unit from signaling correctly over the I/O interface.

Disconnect-in can be activated by a control unit only when it is connected to the channel (has operational-in up). The channel performs a selective reset in response to disconnect-in and indicates to the operating system the occurrence of disconnect-in by causing an I/Ointerruption and posting an interface control check in the channel status word (CSW).

Channel Command Retry

Command retry is a combined channel-control unit procedure that can cause a command to be retried without requiring an I/O interruption.

Based on such factors as whether operator intervention or program reorientation is required before retry, the control unit determines if the last command can be retried.

Command retry applies only to block-multiplexer channels and requires an additional interface line called *mark 0-in*.

Channel Data Rates

Each channel operates on a time sharing basis. Each channel is assigned a certain time slot to minimize the impact on each channel's throughput in relation to other channel loading. Each channel's data rate is not affected by the other channels.

Channel configurations and maximum data rates are shown in the appendixes.

Input/Output Interruptions

The 4341 input/output interruptions are described in the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode.*

Byte-Multiplexer Channel(s)

For byte-multiplexer channel operations, the status, data, or control communication with a device is coordinated by the standard-interface signal sequences between the device and the byte-multiplexer channel.

Every device attached to the byte-multiplexer channel must have a unit control word (UCW) assigned to it. When a byte-multiplexer channel operation requires using the information in a UCW, that UCW is read from system (auxiliary) storage into channel local storage. The operation specified by the UCW is then performed, and the UCW is updated and returned to auxiliary storage when the operation is completed. The UCW carries a dynamic record of the operation for the I/O device assigned.

Note: I/O devices, such as magnetic tape units, that operate in burst mode and are subject to data overrun (that is, the possibility of data loss) are not supported on the byte-multiplexer channel.

Block-Multiplexer Channels

Block-multiplexer channels are available on the 4341. The block-multiplexing capability allows concurrent operation of many I/O devices on the channel's single data path.

Although only one device may actually be transmitting data at any given instant, multiple channel programs can be concurrently active.

Block multiplexing involves temporarily disconnecting an operation in a sequence of chained channel commands. This frees the channel during nondata transfer activity of the device, and allows other devices access to the channel during this time.

The multiplexing facility of the block-multiplexer channel allows interleaved execution of several channel programs by the same channel.

I/O equipment with high data rates is normally attached to block-multiplexer channels. Non-multiplexing devices attached to these channels operate in selector-channel mode.

Subchannels and UCWs

A maximum of 256 I/O unit addresses is available on each channel. An exception is channel 0, which can have a maximum of 240 I/O unit addresses because addresses 0F0-0FF are reserved. However, the maximum number of UCWs is 1,024, which are assigned to I/O devices as required. Each control unit or device requires a subchannel (UCW). With a shared control unit, several different I/O devices may share one UCW. Examples are:

- A single control unit that controls one I/O unit. For example, the 1443 Printer Model N1 requires one subchannel (UCW).
- A single unit that contains several control units. For example, the 2821 Control Unit that handles functions for each 1403 Printer, the 2540 Reader section, and for the 2540 Punch section, requires three subchannels (UCWs).
- A single control unit that services the requirements of several devices at once. For example, the 3830 Storage Control (used with 3330 Disk Storage devices) requires a subchannel for each device.
- A single control unit that services the requirements of several I/O units (one at a time). For example, each 3272 Control Unit requires an exclusive UCW, but all 3277 Displays serviced by that control unit share the UCW assigned to the 3272.

UCW Pool

Up to 1,024 subchannels (UCWs) can be configured. A minimum storage allocation of 8,192 bytes is reserved for the first 128 UCWs. (Each UCW occupies 64 bytes.) If more than 128 UCWs are required, additional groups of 32 UCWs (2,048 bytes of main storage) can be allocated, up to the maximum 1,024.

Refer to the *IBM 4341 Processor Operator's Guide* for information pertaining to the generation of UCWs.

Each group of 32 additional UCWs reduces the usable processor storage by 2,048 bytes. For example, 128 UCWs require 8K of storage, and 1,024 UCWs require 64K of storage. For each increase in the number of UCWs, the additional UCW storage requirement limits the high address boundary.

Sixteen UCWs (0F0 through 0FF) are reserved for natively attached I/O and for internal functions.

The UCW contains all control information necessary for a channel to perform I/O operations to an attached device.

Each channel has a UCW through which all unassigned I/O devices can present asynchronous interruptions.

Each device attached to a channel must have a UCW associated with its address.

Channel UCW Directory

Each channel has a channel directory. Each directory has 256 entries, one for each possible device address (00-FF) on the channel.

Each directory entry contains the reference number of its associated UCW.

UCW Assignment

UCWs for natively attached I/O devices (such as the 3278-2A) are preassigned. The logical addresses for these devices are assigned by the user. All other I/O devices must be described to the processor by the service representative at the user's request. This information is kept in the UCW directory tables for automatic assignment of UCWs by the processor.

All UCW assignments are written onto the system diskette and become effective after subsequent IMLs. Note, however, that UCW reassignments for natively attached equipment become effective immediately (without a re-IML).

Device Considerations

Devices that share a control unit and operate in selector mode can use one common shared UCW on the channel. The SEL mode bit must be on in the directory entry.

Devices capable of running in block-multiplexer mode may use an *unshared* UCW for each actual device attached, or one *shared* UCW for all devices attached to the control unit. Normally, UCWs for the block-multiplexer channel are unshared, with SEL mode off.

Devices, such as the IBM 3272 Control Unit, require one exclusive UCW for each control unit on the channel. Each 3277 attached to that control unit then shares that control unit's UCW. The shared bit must be on in the directory entry for that UCW. Magnetic tape devices use a shared UCW and operate in selector mode.

Channel-to-channel adapters are treated as control units, and require one UCW for each interface attached to the processor.

Channel Operation

Initial selection, interruptions, and channel status are controlled by microcode and hardware circuits. When the channel operation is set up, hardware controls the data transfer. Each channel has 256 bytes of channel data buffer. Depending on the data length and address boundary, hardware controls the transfer of data to processor storage:

- in 64-byte blocks
- in a partial block to line up to a 64-byte address boundary
- in a partial block to complete a data transfer
- in a partial block for very short records
- in a partial block for byte-multiplexer operation.

Display Console

The IBM 3278 Model 2A Display Console or the IBM 3279-2C Color Display Console is the principal device provided for the operator to communicate with the system. The operator can use the keyboard and the display console to control system operation, as well as to display system status. The display consoles and printers are attached to the 4341 via the support processor.

The minimum 4341 configuration requires one 3278-2A Display Console or 3729-2C Color Display Console.

Depending on the mode of operations, three additional devices may be attached to the support processor ports:

Port Devices

- 0 3278-2A or 3279-2C
- 1 3278-2A or 3279-2C or 3268-2 or 3287
- 2 3278-2A or 3279-2C or 3268-2 or 3287
- 3 3278-2A or 3279-2C or 3268-2 or 3287

Two operating modes are described in this section: Display mode and Printer/Keyboard Emulation mode.

Display Consoles

The display console can be a 3278 Model 2A Display Console or a 3279 Model 2C Color Display Console with their keyboard. Both uppercase and lowercase characters can be entered and displayed.

The 3278-2A or 3279-2C has a total screen size of 2000 characters (25 lines of 80 characters each). The character positions on line 25 are used exclusively for indicating display console status. The remaining display screen is partitioned into two regions. The upper 20-line region is used and managed by the operating system. Lines 21 through 24 are reserved for

displaying system status. Lines 21 through 25 are not available to the user.

Printers

The 3268 Model 2 Printer is a matrix printer which can be attached as a console printer. This printer can operate at up to 340 characters per second, with 10 or 16.7 horizontal spacing, and 3, 4, 6 or 8 lines per inch vertical spacing, with all selections manually operated.

The 3287 Printer has a nominal print speed of 80 characters per second (Models 1 and 1C) or 120 characters per second (Models 2 and 2C). A print operation can be initiated by application programs or by the operator using the Copy key. The 'C' version of the 3287 Printer can print in color.

Operator Control Panel

The operator control panel (OCP), used with the display console, controls and monitors 4341 operation. With the OCP, the operator can:

- Power the system on and off
- Load the support processor microcode
- Observe system status

3278-2A or 3279-2C Display Console

The display console allows visual communication between the operator and the system. The operating system uses the display to pass messages to the operator, to present operating modes for selection and further definition, and to display information accessed or entered at the keyboard. The 3279-2C Color Display Console displays the same information as the

DISPLAY MODE
PTR/KYBD MODE
MANUAL CONTROL
DISCONNECTED

INSERT MODE (or blank)

INHIBITED PTR-BUSY PTR-INTV REQD PTR-CHECK USAGE CONFLICT

Figure 4. Display Console Status Indicator Layout (Line 25)

3278-2A. However, the information may be displayed in up to four colors (red, green, blue, and white) for improved readability.

4341 Display Console Indicators

Messages on line 25 of the 3278-2A or 3279-2C Display Console screen indicate the status of the device (Figure 4). Except when in Test mode, the indicator line is identified by the ▶ symbol as soon as the display console 'power on response' is accepted. Only one indicator in each field is active at any given time.

Display Mode indicates that the display console is available to the host operating system, using 3272 control unit interface or equivalent support.

Ptr/Kybd Mode (printer/keyboard emulation mode) indicates that the display console is available to the host operating system, using 1052, 3210, or 3215 Console Printer/Keyboard interface support.

Manual Control indicates that the display console is under the manual functions control control and is not available to the operating system. In this mode, a Start I/O (SIO) to the display console is accepted but is held pending until the device is available to the operating system. If an SIO is enqueued, the audible alarm sounds and SYSTEM MESSAGE WAITING is displayed on line 23.

Disconnected indicates that the device is not logically connected to the 4341, is not available to the operating system, and is not in use for manual functions. This condition exists when no unit address has been assigned to the display console.

Insert Mode (keyboard insert mode) is displayed after the Insert key has been pressed and is reset by pressing the Reset key.

Inhibited indicates that the keyboard input is inhibited because:

- An Attention key (a PF key, ENTER, CNCL (PA2), REQ (PA1), MODE SEL, CHG DPLY, or DIAG) was pressed.
- The operator attempted to alter a protected field.
- The operator attempted to insert a character into a field that had no nulls.

- Keyboard overrun (caused by multiple simultaneous key entries) occurred.
- The Start, Stop, INTR, MODE SEL, CHG DPLY, or DIAG keys were pressed when *another* display console was already in Manual Control mode, or when a previous request had not been completed.

When the keyboard is locked, the Inhibited indication is displayed. Certain functions, such as the Reset key, are accepted and processed when the keyboard is locked.

This indicator is reset by:

- Pressing the Reset key, or
- When the application program issues a Write command which specifies 'keyboard restore.' (This also resets the PTR-Busy, PTR-INTV REQD, and PTR-Check indicators.)

Ptr-Busy (hard-copy printer busy) is displayed when a Copy request is issued to a hard-copy printer that is busy with a previous Copy request or with an SIO from the operating system. The Copy request is ignored.

This indicator is reset by pressing the Reset key.

Ptr-Intv Reqd (hard-copy printer intervention required) is displayed if a Copy request is rejected because of an error condition from which the operator can recover, such as:

- End of forms
- Power off or in Test mode
- No hard-copy device assigned.

This indicator is reset by pressing the Reset key.

Ptr-Check (hard-copy printer equipment check) appears when an equipment check condition is detected while attempting to perform a Copy request.

This indicator is reset by pressing the Reset key.

Usage Conflict This indicator displays when a function is not allowed at the present time. This occurs if the START, STOP, INTR, MODE SEL, CHG DPLY, or DIAG key, etc., is pressed when another display console is already in manual functions control mode or when a previous request has not been completed. This indicator is reset by pressing the RESET key.

3278-2A Switches and Controls

The switches located on the 3278-2A Display Console control the operation of the unit.

Power On/Off and Normal/Test Switches can be used to make the device ready (On and Normal positions) or not ready (Off or Test positions). In the not ready case, intervention required in the sense byte is set and Start I/O instructions to the device are rejected. When a not-ready-to-ready transition occurs, a device-end status is presented. Whenever the device is made not ready this way, the current display console image is lost and is not recoverable. The top 20 lines of the screen are blank when the device is made ready.

Mono/Dual Switch determines whether lowercase alphabetic characters are displayed in lowercase or uppercase. When the switch is set to Mono, all alphabetic characters are entered and displayed in uppercase. When set to Dual, both uppercase and lowercase characters are displayed.

Contrast/Brightness and Alarm-Volume Controls are used to set up display console and alarm conditions appropriate to the operating environment or operator preference. These controls have no effect on the operation of the device and cause no error conditions.

3279-2C Switches and Controls

- Power on/off
- Test/Normal
- Mono/Dual
- Brightness
- Volume
- Two Color/Four Color

The 3279-2C switches and controls function the same as those on the 3278-2A. The Two Color/Four Color switch (00/0000) enables the display to change from a base four-color presentation to a two-color (green/white) presentation which duplicates the 3278-2A mono-color presentation with intensification display.

Display Console Keyboard

The display console keyboard is the operator's primary input device to the system. The keyboard controls the display console and provides a means of signaling the program. The keyboard allows operator communication with the processor to:

- Enter data
- Answer program generated requests
- Perform manual functions
- Enter system configuration.

The functional key groups are: Shift keys, alphanumeric and graphic keys, cursor control keys, input control keys, Program Attention keys, system function keys, and the Copy key.

- Shift keys generate unique codes that are interpreted and acted upon accordingly.
- Alphanumeric and graphic keys are interpreted as data and displayed on the screen. Alphameric data characters appear on the display console screen at the cursor location, unless the cursor is in a protected field or in an attribute character location. The Enter or Program Function keys are used to indicate that data entry is complete. The alphabetic characters are displayed on the screen in either uppercase or lowercase.
- Cursor control keys reposition the cursor on the screen.
- Input control keys cause the character(s) in the input field(s) to be inserted, deleted, or erased.
- Program Attention (PA) keys generate an attention interruption to notify the application program. An attention identification (AID) character is generated at the time of interruption to identify which Program Attention key caused the interruption.

		Alternate	
Normal/Shift Mode		Mode	
Function	AID	Function	AID
CANCEL (PA2)	6E	PF1	F1
REQUEST (PA1)	6C	PF2	F2
ENTER	7D	PF3	F3
		PF4	F4
		PF5	F5
		PF6	F6
		PF7	F7
		PF8	F8
		PF9	F9
		PF10	7A
		PF11	7B
		PF12	7C

- System function keys (such as Start, Stop, and INTR) are used to control the 4341 processor, and are transparent to the user program.
- Copy key performs a print operation, and is active only when the manual functions screen is in use. The data printed on the printer appears in the same characters and format that appears on the display console screen.

Display Console Testing

The 3278-2A or 3279-2C Display Console has facilities to test the refresh buffer, keyboard, and execution of device Write commands.

A description of the available tests and their functions and controls are contained in the Problem Determination Guide located in the console keyboard.

Display Mode

Display mode supports the 3278-2A Display Console or 3279-2C Color Display Console and an optional printer.

The user screen size is limited to 20 lines; only 1600 bytes of the 3278-2A or 3279-2C Display Console device buffer is available to the operating system. For the printer, 24 lines (1920 bytes) are available.

The display consoles and printers have unique unit addresses and are treated as independent devices. The device addresses (X'000'-X'0EF') on channel 0 can be selected by the operator at any time.

When the operating system or application program requires the display console for service, it issues an SIO to the channel. The channel starts the display console operation by issuing one of the following 3270-mode CCW commands. For more detailed information about 3270 commands, see the *IBM 3270 Information Display System Component Description*, GA27-2749.

3270-Mode CCW Commands

<u> </u>	CON C	Initial	Ending	Async.
Code	CCW Command	Status	Status	Status
01	Write	00	CE	DE
05	Erase/Write	00	CE	DE
02	Read Buffer	00	CE,DE	
06	Read Modified	00	CE,DE	
0B	Select	CE		DE
0F	Erase All Unprotected	CE		DE
03	No Operation	CE,DE		
04	Sense	00	CE,DE	
E4	Sense I/O	00	CE,DE	

Device Status

The display console status byte presented to the operating system can be generated synchronously or asynchronously. Synchronous status is passed to the host channel as ending status to a command.

- Bit Status
- 0 Attention
- 1 Not used
- 2 Not used
- 3 Busy
- 4 Channel end
- 5 Device end
- 6 Unit check
- 7 Unit exception

Initial Status reflects the condition of the selected device upon receipt of a command, and indicates to the channel if the command can be executed.

Ending Status reflects the condition of the selected device after all channel interface operations of a non-immediate command are completed.

Asynchronous Status reflects ending status for an immediate command (other than no operation), a second ending status for a Write or Erase/Write command, or an equipment condition or operator action (attention) not associated with command execution.

Sense Byte

When an error is detected, the appropriate bit(s) are set as follows to describe the condition that caused the error. This sense byte is sent to the host channel when a Sense command is executed. Sense data is reset by every command except NOP.

- Bit Sense
- 0 Command reject
- 1 Intervention required
- 2 Bus-out check
- 3 Equipment check
- 4 Data check
- 5 Unit specify
- 6 Control check
- 7 Operation check

Write Data Record

Programming for a display console differs from most other I/O devices in that the CCW commands for a display console are fairly elementary, and the detailed data positioning and control attributes are imbedded within the data record that is transferred by the CCW.

To control the information displayed on a display console, the application program must provide a write data record that includes a write control character (WCC), buffer control orders, buffer address, attribute, and data. Only the attribute and data are stored into the device buffer. Invalid or undefined data (EBCDIC) in the data record is displayed as a hyphen.

Bit	Write	Control	Character	(WCC
Du	<i>w</i> rne	Control	Churacier	$(m \cup \cup$

- 0 Not used
- 1 1
- 2,3* Printout format
 00=the NL order in the data record; determines print line length
 01=specifies 40-character print line
 10=specifies 64-character print line
 11=specifies 80-character print line
- 4* Start print
- 5 Sound alarm
- 6 Restore keyboard
- 7 Reset MDT
- * Used only when the Write is directed to the 3287 Printer.

Code Buffer Control Orders

- 1D Start field (SF)
- 11 Set buffer address (SBA)
- 13 Insert cursor (IC)
- 05 Program tab (PT)
- 3C Repeat to address (RA)
- 12 Erase unprotected to address (EUA)

- Bit Attribute Character
- 0 Determined by the contents of bits 2-7
- 1 1
- 2 0=Unprotected 1= Protected field
- 3 Bits 2,3=11 cause an automatic skip
- 4,5 00=Normal display
 01=Normal display
 10=Intensified display
 11=Nondisplay, nonprint
- 6
- 7 Modified data tag (MDT)
 - 0=Field has not been modified
 - 1 = Field has not been modified by the
 - operator, or set by program in the
 - data record

If a buffer address is specified past the end of the user buffer area (20 lines for the 3278-2A and 3279-2C, 24 lines for the 3268-2 or 3287), operation check sense is set. The command is ended with device end and unit check status. When the operation check occurs because of an invalid buffer address, channel end may have been sent before the operation check was detected. The residual count filed is unpredictable in this case, and cannot be used to precisely determine the location of the error in the write data record. The CCW address stored in the CSW may also be past the CCW pointing to the invalid data if data chaining is specified for the write command. The write data record must be inspected to locate the invalid address specification.

Color Console and Printer

The colors presented on the color console and printer are determined by the existing protection and intensification attribute bits of the displayed field.

The attribute bits, their field characteristics, and colors presented are:

Bits	Attribute Characteristic	Color
00	Unprotected, Unintensified	Green*
01	Unprotected, Intensified	Red
10	Protected, Unintensified	Blue
11	Protected, Intensified	White
	Black* (3287-1C, 2C)	

*Green and black may be interchanged on the 3287-1C or 2C when the feature is ordered.

Read Data Record

Depending on the command, these types of Read data records are generated:

• Read Buffer data record: In response to the Read Buffer command, the read data record is generated, beginning with a three-character read heading that consists of the Attention Identification (AID) character, followed by a two-character cursor address.

The contents of all device buffer locations (lines 1-20 for 3278-2A or 3279-2C, lines 1-24 for 3268-2 or 3287) are transferred, including nulls. Start field (SF) order codes are inserted before each attribute character to identify the beginning of each field.

• Read Modified data record: In response to the attention interruption as a result of pressing the Enter key or a Program Function key, the application program issues a Read Modified command to the display console. The first three-byte read heading of the read data record is always the AID code and the two-byte cursor address.

Following the read heading is the data of each modified field (lines 1-20 for 3278-2A or 3279-2C, lines 1-24 for 3268-2 or 3287). The data for each field is preceded in the data record by a Set Buffer Address (SBA) order code, followed by the two-byte buffer address of the first character position in that field (the attribute address +1).

Programming Information

The concepts of protected or unprotected data, and the modified data tag function, are basic to the modes of operation.

Protected Data Autolock

A program-controlled facility of the display console allows fields to be defined as *protected* or *unprotected*. The operator may not alter the protected field in any way. If an attempt is made to enter an alphameric character when the cursor is located in a protected field, the keyboard becomes disabled by the autolock function.

In an unprotected field, the operator can enter, modify, and erase alphameric data.

Modified Data Tag

To identify data fields that have been modified, the modified data tag bit (bit 7 of the attribute character) is set to a 1. This process is called *tagging* the field.

Data in protected or unprotected fields can be tagged as having been modified. In protected fields, the tags are set under program control. In unprotected fields, modified data tags are also set by keyboard operation. These tagged fields are the only fields transferred upon execution of a Read Modified command.

Display Mode Error Handling

When attempting any recovery after a unit check is reported to the operating system, use error recovery procedures in the *IBM 3270 Information Display System Component Description*, Order No. GA27-2749.

3268 Model 2 and 3287 Printer

This optional printer provides a hard copy of the information that is displayed on a 3278-2A or 3279-2C Display Console, or of information written from the application program.

All 3270-mode CCW commands and data record to the display console can be issued to the printer. Printouts can be formatted the same way as a display image. Cursor information is ignored by the printer.

When a print operation is specified by a Write command addressed to the printer, the print line format in which the data is to be printed can be specified as part of the command in one of three printer formats. These formats simply define the print line length: 40, 64, or 80 character positions per line. If a format is not specified, the print line length is set to the default for the Printer.

Printer control orders (NL, EM, and FF) are transferred as part of the data record from the application program. They are stored in the printer buffer as data.

Note: If a write buffer has been issued but the buffer not printed, the printer appears busy to a Copy key request.

Error Handling

The recovery procedures for errors detected by the 3268-2 and 3287 are categorized in three ways:

- Automatic recovery The Alphanumeric Readout (ANR) may or may not be used to indicate the reason code.
- Manual intervention recovery Printer check light is turned on and ANR indicates the reason code.
- Machine stop Printer check light is turned on and ANR indicates the reason code.

Printer/Keyboard Emulation Mode

The printer/keyboard emulation mode allows the processor to run operating systems and programs designed for devices such as the 1052, 3210, or 3215. Although physically different, these devices accept the same commands and respond in a similar fashion. They are emulated on a Display Console with a recommended optional Printer coupled as a hard-copy device.

Printer/keyboard emulation mode requires one 3278-2A or 3279-2C Display Console. Options allow three additional 3278-2A, 3279-2C, or 3268-2 or 3287 devices to be attached to the support processor and configured as needed.

In *coupled* mode, all data read from or written to the device is printed on the Printer. In *uncoupled* mode, only up to the last 18 lines of data are displayed on the 3278-2A or 3279-2C. No hard-copy record is made. To couple, assign the same device address to a display console and printer by using the console functions program load (L) screen.

Both the display console and the printer appear to the operating system as one device. The device addresses used are selected by the operator and can be modified at any time. A maximum of two printer/keyboard devices can be configured at a time.

In printer/keyboard emulation mode, the display console accepts printer/keyboard commands and responds with status and sense information. The data received or sent with the commands is formatted to appear on the display console and optional printer in a manner similar to the actual devices being emulated. Indicators and keys to emulate printer/keyboard functions are displayed on the 3278-2A or 3279-2C Display Console to allow the functions to be performed by the operator.

Screen Management

When the display console is in printer/keyboard emulation mode, the entire screen is controlled by the emulation facility. The screen is initialized the first time printer/keyboard emulation mode is entered, either via operator action (Attention keys or making device ready) or by channel action (SIO or reset). After the screen is initialized, the operator can only enter data on the screen when a Read command is issued. At all other times, the entire display is protected from operator alteration.

The screen is divided into three areas. Lines 1-18 contain the message area for displaying the operator input and host program output messages. Lines 19-20 contain the operator input area. Up to 126 characters are allowed for each read operation. Part of line 20 displays 1052 keyboard indicators (REQUEST, PROCEED, ALARM, and INTV REQD). Lines 21-24 display the system status. The character positions on line 25 are exclusively for indicating display console status.

A message in the input area transferred to the host program is also displayed in the message area. Input and output messages are displayed sequentially starting from the top (line 1) as if they were printed on the printer/keyboard device.

When the display message reaches the bottom (line 18), scrolling takes place. The top six lines (lines 1-6) are rolled off the screen and the bottom six lines (lines 13-18) are made available for subsequent input.

Display Area

The display area starts in line 1, column 1, and ends on line 18, column 79. The display area is protected by an attribute at line 20, column 80. This area shows the last 18 lines of data that has been read or written to the device. The data is organized into lines corresponding to lines of printer/keyboard output, with each new line beginning in column 1 of a display area line. A line longer than 80 characters (the display area line length) wraps to the next line in the display area, to a maximum of 126 characters. To allow the maximum amount of data to be displayed, multiple new line characters (NL = X'15') in a write data record result in only one blank line on the display console. This prevents data from being scrolled off the screen when new line characters are added to space the printer output. The new line suppression only applies within a command; each new command is treated separately. Therefore, multiple blank lines caused by two or more commands are not suppressed.

Input Area

The input area is defined by an attribute at line 18, column 80, and contains 126 characters on lines 19 and 20. This area is unprotected only when a Read command is in progress (Proceed indicator displayed) to permit data entry or modification. At all other times, it is protected from operator alteration.

Pressing the Enter or Cancel key transfers the data in the input area across the channel. Data accepted by the channel for a Read command is then moved to the display area and printed. If the channel terminates data transfer before all data entered by the operator was sent, only that data accepted appears in the display area.

The field initially contains all nulls (X'00') that are compressed out of the Read data and not transferred to the channel. If a space (X'40') is desired in the response, the Space key must be used to enter it. Cursor movement keys leave nulls in the input area. These nulls are compressed from the data.

Indicator Area

The indicator area is a brightened field defined by an attribute on line 20, column 47, and is 32 characters in length. This area contains the following visual indicators:

Proceed appears whenever a Read command is in progress for entry of data in the input area. Proceed is cleared when the Read command is ended by any of the following conditions:

- Normal ending caused by Enter or Cancel key
- A system or selective reset
- A Halt Device or Halt I/O instruction issued to the device.

Request indicates that a REQ (Request) key attention status has been stacked because the device is busy

executing a command. When the current operation completes, the attention status is presented to the channel, and the Request indicator is reset. A system or selective reset also resets the Request indicator.

Alarm appears, and the audible alarm sounds, whenever an Alarm command is received. The indicator is reset by:

- Pressing a Program Function (PF) key at any time
- Pressing the Enter or Cancel key when Proceed is not displayed
- System or selective reset occurs

Intervention Required (INTV-REQD) appears if the 3268-2 or 3287 Printer coupled as a hard-copy device becomes Not Ready. This condition occurs when end-of-forms, power off, or other check conditions are present. At this time, a Start I/O is not accepted. When the check condition is cleared, the indicator is erased, and device-end status is presented to the channel.

3278-2A and 3279-2C Console Keyboard Operation

The 3278-2A or 3279-2C Display Console keyboard is used to communicate with the processor. In printer/keyboard emulation mode, the display console always contains a formatted screen defining protected and unprotected fields. The response for the different keys is:

Alphanumeric and Special Character Keys can only be used when an unprotected field is present, which is only during a Read command (Proceed indicator displayed). Such use causes the character to appear and the cursor to advance. At all other times, or if the cursor is not in the input area, using these keys causes the Inhibited message to appear on line 25 of the display console. The inhibited condition can be cleared with the Reset key.

Request (REQ) Key initiates communication with the operating system by sending an attention status (X'80'). If the request status cannot be sent immediately because of a busy condition, the Request indicator appears, and the attention is stacked. When the device becomes not busy, the attention is sent, and the Request indicator is reset.

Note: Do not confuse the REQ and INTR keys: the red INTR key causes an External Interrupt, and is not used with printer/keyboard emulation mode operation.

Enter Key is used during a Read operation when data entry is complete. The data is read, and the input area is cleared and protected. This key is equivalent to the End-of-Block (EOB) or End keys on the emulated device. If no Read is in progress, the Alarm indicator is reset, and the keyboard is unlocked.

Cancel (CNCL) Key serves the same function as the printer/keyboard Cancel key. During a Read operation, this key sends a cancel response to the channel. The cancel response is a channel end with unit exception (X'09'). The input area is cleared and protected, and an asterisk (*) is written in the display area. No data is transferred. If no Read is in progress, the Alarm indicator is reset and the keyboard unlocked.

Program Function Keys (PF1-PF12) reset the Alarm indicator and unlock the keyboard.

Cursor Movement Keys move the cursor without causing any modifications to the screen. They may be used at any time and cause no errors. The cursor is positioned at the start of the input area when a Read command is initiated. The cursor must be in the input area to enter data. The Tab and Backtab keys always position the cursor at the beginning of the input area if a Read is in progress. Pressing the New-Line key places the cursor in the first column of either line 19 or 20.

Erase Input and Erase EOF Keys clear the input area when a Read is in progress. The Erase Input key clears the entire field and repositions the cursor. The Erase EOF (End-Of-Field) key clears the field from the current cursor position to the end of the input area.

Insert and Delete Keys selectively add and delete characters when entering and altering data in the input area.

3268 Model 2 or 3287 Printer

A 3268-2 or 3287 (optionally coupled as a hard-copy device in printer/keyboard emulation mode) can print all data that has been transferred. The data is arranged in lines of 126 or fewer characters, as they would appear on the emulated printer. Each line is printed as soon as it is completed; that is, when any of the following conditions occurs:

- 126 bytes of data have been received since the beginning of the line.
- A New-Line character (X'15') is found in a Write or Write-ACR command data record.
- All data has been received on a Write-ACR command.
- After a Read command is executed.
- A system or selective reset occurs.

Note that, after a Write command (X'01') is executed but a complete line has not been accumulated, the last partial line is not printed. The data is accumulated in the printer buffer and is printed when the line is completed by succeeding commands. The printer is busy to a Copy key request at this time.

A system or selective reset causes a line feed to be performed if no data is present in the buffer.

Printer/Keyboard Commands

Printer/keyboard emulation mode accepts and executes all commands that are valid for the emulated device. The valid commands, and the normal status responses received, are:

		Initial	Ending	Async.
Code	Command	Status	Status	Status
01	Write	00	CE	DE
09	Write ACR	00	CE	DE
0A	Read Inquiry	00	CE	DE
04	Sense	00	CE,DE	
03	NOP	CE,DE		
0B	Alarm	CE,DE		
E4	Sense I/O	00	CE,DE	

It is assumed that an alternate console printer has been coupled as a hard-copy device. If not, ending status for a command is presented to the channel after the data has been placed in the Display Console area.

Write

The Write command transfers data from the channel to an internal buffer, and then processes a line at a time to the display area on the display console and prints it on the printer. When all of the data has been received, channel-end status is returned and the last line is processed. If the last line is not complete (ended with New-Line character or exactly 126 bytes), device-end status is sent after the data has been placed in the display area and saved in the printer buffer without printing. If the command ends with a complete line, device-end status is sent when printing is finished. The Write command continues requesting data from the channel until the channel stops data transfer (when the CCW count reaches zero). Therefore, Incorrect Length channel status is always indicated with a zero residual count unless the Suppress Incorrect Length Indicator (SILI) CCW flag is on.

Write-ACR (Automatic Carriage Return)

This command is similar to the Write command, except that at the end of processing, the last line begins printing, and an automatic carriage return (ACR) is performed.

Read-Inquiry

When the Read-Inquiry command is received, the input area is unprotected and the Proceed indicator is turned on. The command then waits until the operator signals that data entry is complete (by pressing the Enter or Cancel key).

If the Enter key is pressed, the data is read from the input area and transferred to the channel. (Channel end is presented after data transfer.) Then any data the channel accepted is written to the display area and printed on the printer.

If the Cancel key is pressed, a channel end with unit exception is presented to the channel, and an asterisk (*) is written to the display area and printed.

After printing is completed, a device end is sent to terminate the command.

Sense

The Sense command transfers one byte of sense information. After the sense byte is accepted, channel end and device end are presented as ending status. The sense byte is reset at the initiation of any command except Sense. The sense byte is defined as:

Bit	Code	Sense Information
0	CR	Command Reject
1	IR	Intervention Required
2	BOC	Bus-Out Check
3	EC	Equipment Check
4		Unused; always 0
5		Unused; always 0
6		Unused; always 0
7		Unused; always 0

Sense I/O

This command is used to identify device type. If the device is not busy or not ready, seven bytes of fixed data are presented:

Printer/Keyboard Mode PTR/KBD	X'FF434100105200'
Display Mode	
3278-2A	X'FF43410032782A'
3279-2C	X'FF43410032792C'
3287	X'FF434100328700'
3268-2	X'FF434100328700'

No Operation (NOP)

This control command results in an immediate channel end and device end. No action is performed in the device. This command can be used to clear any outstanding status, or to cause a command chain to end with a channel end and device end together.

Alarm

This control command sounds the audible alarm on the display console and displays the Alarm indicator. If the display is Not Ready, the command functions as a NOP. Immediate channel end and device end are presented as initial status. The command is accepted even if an intervention required condition exists.

Device Status

The following status bits are set:

Bit	Status

- 0 Attention
- 1 Unused
- 2 Unused
- 3 Busy
- 4 Channel end
- 5 Device end
- 6 Unit check
- 7 Unit exception

Unlike the display console, the following interfaces are specific for the printer/keyboard Read operation:

- The attention bit is set to one when the Request key is pressed. This attention status is presented to the channel when no other operation is in progress. The host program should react to the attention interruption by issuing a Read command.
- No AID is generated for an attention interruption.

Error Conditions

The following error conditions can occur while operating in printer/keyboard emulation mode:

3278-2A and 3279-2C Display Consoles Not Ready

When the 3278-2A or 3279-2C is not ready because of one of the following conditions, a Write, Write-ACR, Read, or NOP command is rejected. The Sense and Alarm commands are always accepted.

Power Off: When the Display Console is not powered on, an intervention required condition exists and is reported to a Sense command.

Test Mode: When the 3278-2A or 3279-2C is in test mode (Test/Normal switch in Test position), an intervention required condition exists and is reported to a Sense command.

Device Not Functional: If the 3278-2A or 3279-2C hardware is failing, or the device is incorrectly configured or not connected, an intervention required or equipment check condition exists and is reported to a Sense command.

Console Printer Not Ready

The 3268-2 or 3287 Printer is in a not ready state because of any of the following conditions:

Power Off: When power is off to the 3268-2 or 3287, an intervention required condition exists and is reported to a Sense command.

Test Mode: When the printer is in test mode (by pressing the Test switch), an intervention required condition exists and is reported to a Sense command.

End-of-Forms: When the End-of-Forms switch indicates that no paper is in the printer and a one-minute timeout occurs, an intervention required condition exists and is reported to a Sense command. The INTV-REQD indicator is displayed only after the timeout condition occurs. This delay allows the operator to correct the end-of-forms condition and have the continue printing with no errors reported or software retry required. The end-of-forms condition causes the audible alarm to sound until the HOLD PRINT switch is pressed. The end-of-forms condition is cleared by pressing the HOLD PRINT switch, replacing the forms, then pressing the ENABLE PRINT switch. *Hold Print Timeout:* When the HOLD PRINT condition lasts longer than ten minutes, an intervention required condition exists and is reported to a Sense command. The INTV-REQD indicator is displayed only after the ten-minute timeout occurs.

Device Not Functional: When the console printer is not operating because of error conditions, incorrect connection, or invalid configuration, an intervention required or equipment check condition exists.

Restrictions and Functional Differences

When operating in printer/keyboard emulation mode, consider the following:

Timing

The length of time taken to execute a command may differ from the emulated devices. Typically, the time from Start I/O until the associated channel end is much shorter, and the time from channel end to device end is longer, due to buffering, because printing is not started until a complete line is received.

The printing speed of the 3268-2 and 3287 is much faster than the speed of the 1052 and the other emulated devices. Using printer/keyboard emulation mode results in a considerable gain in throughput.

When running in uncoupled mode (3278-2A or 3279-2C without a Printer), a timing delay of 0.5 seconds per line is added to improve display console readability.

Keyboard Differences

The 3278-2A or 3279-2C and emulated keyboards differ in both the number of keys and in the keyboard layout. Because all of the keys on the 3278-2A or 3279-2C can be used even if the emulated printer/keyboard has no corresponding key, it is possible to read and write characters that are not implemented on the emulated device. Note that the Carriage Return key on the emulated keyboard is not implemented. Programs that require the use of this key are not supported in printer/keyboard emulation mode.

Note that the Carriage Return key on the emulated keyboard is not implemented.

Model-Dependent Information

This section addresses the 4341 implementation of certain 4300 Processors facilities and functions. (The terms used here are defined elsewhere in this manual, or in the *IBM 4300 Processors Principles of Operation for ECPS:VSE Mode* or in the *IBM System/370 Principles of Operation.*)

Addressing of Natively Attached I/O Devices: I/O addresses 0F0 through 0FF are reserved for internal use for natively attached devices.

Condition Code Setting: The nullification of the NC, CLC, OC, XC, TRT and ZAP instruction execution does not cause the instruction to be executed as if it were specified as a no-operation. The Condition Code may be altered by any of these six instructions even though the instruction has been nullified.

Detection of a PSW Loop: A continuous string of interruptions (PSW loop) may be indicated if pressing the Stop key does not stop processing. Under this condition, a system reset may be necessary to stop.

Segment Table Entry: The 4341 does not check bits 4-7, 29, and 30 of the segment table entry for zeros.

Timing Facility Damage: The 4341 does not distinguish among the failure of the three timing facilities: TOD clock, CPU timer, and clock comparator. Any failure of hardware timing facilities causes all three facilities to enter the error state.

Storage Size and Page Capacity Count (PCC): In ECPS:VSE mode, the virtual storage size is 16,777,216 bytes and cannot be altered (as described in the IBM 4300 Processors Principles of Operation for ECPS:VSE Mode). The display console has no storage size control. The value of the page capacity count (PCC) is always 8192.

Timer and Clock Resolution: The interval timer is updated every 3.328 milliseconds. The processor skips one update every 625 updates to derive the average of 3.333 milliseconds. The conditions of losing a decrement update are:

- When the Interval Timer control on the display console is set to Off
- When the processor is not in the operating state

• When the Rate Control on the display console is set to Instruction Step

The time-of-day clock resolution is one microsecond. The 1-MHz oscillator has a tolerance of 0.0027 percent.

The CPU timer and clock comparator have the same resolution as the time-of-day clock. Time-of-day clock updates that are interrupted during an instruction retry are readjusted after the retry.

Reference and Change Recording: The recording of reference and change bits is accurate, with the following exceptions:

- The reference bit may be set because of storage operand fetching of a nullified or suppressed instruction.
- The reference bit may be set because of prefetching of an instruction (instruction buffering). Prefetching can be from a minimum of one instruction to a maximum of eight instructions ahead. This can be a minimum of ten bytes and, up to a maximum of sixteen bytes in advance of the beginning of the current instruction.
- The reference bit may be set because of channel prefetching of CCW, IDAW, or data during an output operation.
- Change bits may be set for the operands of a unit of operation that is nullified because of a page translation exception. For example, the destination operand of a Move instruction may cross a page boundary and encounter a page translation exception at the page boundary. In this case, the instruction is nullified so that the portion of the operand up to the page boundary is restored to its original value, but change bits may remain set after the nullification of the instruction.
- On the Model Group 12 there are two additional situations for which the reference bit will not be set on:

- The reference bit will not be set, if following a RRB or SSK instruction that sets the reference bit to zero, all subsequent accesses to the 2K storage block find the data in the high speed buffer.
- If reference to a 2K storage block is because of a branch instruction destination with alignment 2, 4, or 6 relative to a doubleword boundary, the reference bit will not be set on for this access.

Nontransparent Suppression and Nullification: The channel may observe the effect of temporary storage change of a partially executed, but nullified or suppressed, instruction. This can occur because of an operand access exception, access retry (due to lack of pretest), and instruction retry.

Machine Check Handling: Machine check handling is implemented in the 4341 as follows:

- CR14 Machine Check Control Bits:
 - Bit 4: Recovery Report Mask
 - Bit 6: External Damage Report Mask
- Machine Check Interruption Code (MCIC):
 - Bit 0: System Damage
 - Bit 1: Instruction Processing Damage
 - Bit 2: System Recovery
 - Bit 3: Interval Timer Damage
 - Bit 4: Timing Facility Damage
 - Bit 15: Delayed
 - Bit 16: Storage Error
 - Bit 18: Storage Key Error Uncorrected
 - Bit 20: PSW EMWP Validity
 - Bit 21: PSW Mask and Key Validity
 - Bit 22: PSW Program Mask and Condition Code Validity
 - Bit 23: PSW Instruction Address Validity
 - Bit 24: Failing Storage Address Validity
 - Bit 27: Floating-Point Register Validity
 - Bit 28: General Register Validity
 - Bit 29: Control Register Validity
 - Bit 31: Storage Logical Validity
 - Bit 46: CPU Timer Validity

Bit 47: Clock Comparator Validity

• Other CR14 and MCIC bits are not set by the 4341.

Machine Check Interruptions: Machine check interruptions do not report failures detected during channel operations. A machine check interruption can occur when the malfunction also causes an IPU operation to be retried or terminated.

CCW Prefetch: The channels do not prefetch a CCW for data chaining on input operation.

Power-On State: Before the operator control panel (OCP) indicates power-on complete, the following components must have completed power-on in the sequence:

- 1. Support Processor
 - a. System Diskette Drive
 - b. Adapters
 - c. Display Console
- 2. Processor
- 3. Channel-to-Channel Adapter
- 4. Channel Attached I/O Devices

Power to these components is controlled by the OCP power control.

The optional 3278-2A or 3279-2C Display Consoles (alternate console) and 3268-2 or 3287 Printers enter the power-on state by operator activation of their power-on switches.

Version Code in Store CPU ID Instruction: The version code returned in bits 0-7 of the operand of the Store CPU ID instruction is:

- Hex '05' for 4341 Model Group 9
- Hex '01' for 4341 Model Group 10
- Hex '03' for 4341 Model Group 11
- Hex '04' for 4341 Model Group 12.
Facility Descriptions

This section describes some of the significant 4341 facilities.

Instruction Set

The universal instruction set is implemented in the 4341. For details on instruction word formats and definitions, refer to the *IBM 4300 Processors Principles* of Operation for ECPS:VSE Mode or the *IBM System/370 Principles of Operation*.

Model Groups 9, 10, 11, and 12 support the Engineering Scientific Assist described in the *Multiply Add Facility* manual, GA22-7082, listed in the preface to this document.

Channel-to-Channel Adapter

The channel-to-channel adapter provides a path for operations to take place between two channels and synchronizes those operations. The adapter uses one control unit position on each of the two channels, but only one of the two connected channels requires the adapter to be installed. The adapter is attached to a block-multiplexer channel on the 4341. The other system can be any System/370 or 4300 processor.

The adapter operates in burst mode and transmits data at the rate of the lower-speed channel. The adapter is selected and responds the same as any control unit. The adapter accepts and decodes commands from the channel; however, it differs from a control unit in that it does not use these commands to operate and control input/output devices. The adapter instead uses the commands to open a path between the two channels it connects and then synchronizes the operations performed between the two channels.

| Appendix A. Model Groups 9 and 10

Storage Characteristics

The 4341 Model Groups 9 and 10 are available in the following processor storage sizes:

Model J9: 1,048,576 bytes (1 megabyte) Models K9 and K10: 2,097,152 bytes (2 megabytes) Models L9 and L10: 4,194,304 bytes (4 megabytes)

Virtual storage increases the effective use of processor storage.

Note: The contents of this storage is not saved when power is removed.

Storage Access Width, Data Transfer Rate, and Internal Cycle Time

The processor has an eight-byte parallel dataflow. There is also an eight-byte path among processor, storage, and channels. The processor cycle time is 150- to 300-nanoseconds.

The storage access width is eight bytes. The high-speed buffer storage cycle time for reading or writing of data already in the buffer is 0.225 microseconds.

Detailed timings of the Model Groups 9 and 10 instruction set are given in "Instruction Timing Information."

System Storage Requirements

A portion of processor storage is required for dynamic tables. This reduces the amount of processor storage available for user programming. Depending upon the processor configuration, the reduction of available processor storage may be from 18,432 bytes to 116,736 bytes. The reduction is the sum of the requirements of user selectable options:

- Installed storage size (processor model), plus
- Number of unit control words (UCWs) selected, plus
- Mode of operation, as shown below:

	Models J9 (1 Megabyte) K9. K10 (2 Megabytes)	Models L9, L10 (4 Megabytes)
	Processor	Processor
Mode of	Storage	Storage
Operation	Required	Required
ECPS:VSE	49,152 Bytes	51,200 Bytes
System/370	10,240 Bytes	10,240 Bytes

Number	Processor
of UCWs	Storage Required
128	8,192 Bytes
next 32	+2,048 Bytes
next 32	+2,048 Bytes
etc.,	etc.,
up to:	up to:
1024	65,536 Bytes

Channel Configurations

The input/output channel configurations available for the Model Groups 9 and 10 are:

• Channel Group 1 (standard): One bytemultiplexer channel (1-megabyte/second data rate) and two block-multiplexer channels. In *byte mode*, simultaneous operation of several low-speed devices is permitted. Data transfer can be interleaved.

Note: I/O devices that are subject to data overrun (that is, the possibility of data loss), such as magnetic tape units, are not supported in burst mode on the byte-multiplexer channel.

The two block-multiplexer channels permit simultaneous operation of high-speed devices. Block-multiplexer channels are for relatively high-speed burst operations. They can multiplex complete blocks of data, and thereby permit a device to disconnect only after channel end, or after a halt instruction has been executed. This facility allows the interleaved execution of several channel programs by one channel.

The actual data rates depend upon the types of devices being serviced and the effect of concurrent processor and channel activity. For channel data rates, refer to the Data Rate chart.

 Channel Group 2 (optional): This feature provides three additional block-multiplexer channels. Channel 4 can be selected as a byte-multiplexer channel having a 2-megabyte/second data rate. With this option, the configuration includes two byte-multiplexer channels, one 1-Mb blockmultiplexer channel, one 2-Mb block-multiplexer channel and two 3-Mb block-multiplexer channels. The standard channel configuration is:

Channel 0 - Byte Multiplexer Channel 1 - Block Multiplexer Channel 2 - Block Multiplexer

When the additional channels are installed, the channel configuration is:

- Channel 0 Byte Multiplexer
- Channel 1 Block Multiplexer
- Channel 2 Block Multiplexer
- Channel 3 Block Multiplexer
- Channel 4 Block (or Byte) Multiplexer
- Channel 5 Block Multiplexer

Basic Channel Set (3 Channels)

Channel

Byte-Multiplexer Channel 0* Block-Multiplexer Channel 1 Block-Multiplexer Channel 2

Full Channel Set (6 Channels)

Channel

Byte-Multiplexer Channel 0* Block-Multiplexer Channel 1 Block-Multiplexer Channel 2 Block-Multiplexer Channel 3 Block-Multiplexer Channel 4 (or Byte-Multiplexer Channel 4) Block-Multiplexer Channel 5 Data Rate (See Table Below) 3 Megabytes/Second 3 Megabytes/Second 2 Megabytes/Second (See Table Below) 1 Megabyte/Second

Data Rate

(See Table Below)

3 Megabytes/Second

3 Megabytes/Second

* These devices attached natively to the Support Processor are addressed from the IPU through addresses on the byte-multiplexer channel 0.

Note: The aggregate data rate for two block- multiplexer channels operating concurrently is up to 6 megabytes per second. When five block-multiplexer channels are operating concurrently, the aggregate data rate is up to 11 megabytes per second.

Byte-Mode Operation (Channels 0 and 4)

These data rates are with no other channel activity. For data rates with other channel activity, see 4341 Processor Channel Characteristics.

Byte-Mode Type	Byte Channel 0 Data Rate	Byte Channel 4 Data Rate
Single-Byte Transfer	16 kb/Second	22 kb/Second
Transfer	64 kb/Second	88 kb/Second
Burst-Mode Operation	(Channel 0)	(Channel 4)
Device Type	Channel Data Rat	e
Buffered Input:	1000 kb/Second	2000 kb/Second
Output:	(1000 x device ra divided by (1000 + device r	ate in kb) ate in kb)
For device rates eq	ual to or less than	1000 kb.
Note: These data ra cable and control-u	ites assume small i nit generated delay	nterface- ys.

Figure A-1. Channel Configurations and Data Rates

Instruction Timing Information

This section describes the basic Model Group 9 and 10 instruction timings and storage cycle times.

Processor Instruction Timings

Listed below are the formulas for determining the instruction execution times in nanoseconds. These formulas do not include any allowance for high-speed buffer storage misses, TLB misses, or processor interference because of I/O operations, processor storage refresh, or interval timer updates. These factors must be considered separately.

Accuracy Codes: Formulas with no identified accuracy code yield instruction execution times accurately.

Formulas with Accuracy Code I may not yield exact instruction times, but represent *average* execution times that can be expected in representative instruction sequences.

Formulas with Accuracy Code II yield instruction timing values given below, within 5 percent.

4341 Model Groups 9 and 10 Instruction Timing List

	For-	Mne-	Op-	Accure	acy
Instruction Name	mat	monic	Code	Code	Formula and Comments
Add	RR	AR	1 A		3/75
Add	RX	Α	5A		600+DW2•225+X•150
Add Decimal	SS	AP	FA		1275+RC•525+ZR•1050+(2•DW1+DW2)•225
					(for $L1 \leq 8$ and $L2 \leq 8$)
					1950+RC•675+ZR•1050+(2•DW1+DW2+EDW2)•225
					(for $L1 \leq 8$ and $L2 > 8$)
					2625+RC•675+ZR•1050+(2•DW1+2•EDW1+DW2)•225
					(for $L1>8$ and $L2\leq 8$)
					2700+RC•675+ZR•1050
					+(2•DW1+2•EDW1+DW2+EDW2)•225
					(for L1>8 and L2>8)
Add Halfword	RX	AH	4A		600+DW2•225+X•150
Add Logical	RR	ALR	1E		375
Add Logical	RX	AL	5E		600+DW2•225+X•150
Add Normalized (Extended)	RR	AXR	36	I	3834
Add Normalized (Long)	RR	ADR	2A	I	1050
Add Normalized (Long)	RX	AD	6A	I	1425+DW2•225+X•150
Add Normalized (Short)	RR	AER	3A	I	1134
Add Normalized (Short)	RX	AE	7A	I	1434+DW2•225+X•150
Add Unnormalized (Long)	RR	AWR	2E	I	993
Add Unnormalized (Long)	RX	AW	6E	I	1518+DW2•225+X•150
Add Unnormalized (Short)	RR	AUR	3E	I	1166
Add Unnormalized (Short)	RX	AU	7E	I	1691+DW2•225+X•150
And	RR	NR	14		375
And	RX	N	54		600+DW2•225+X•150
And (character)	SS	NC	D4		1275+DW1•375+DW2•225 (for L≤8)
					900+L8•825+LZ8•75+MA•(floor of AL/8)•225 (for L>8)
					900+L•825+LZ•75 (for operands overlap)
And (Immediate)	SI	NI	94		825
Branch and Link	RR	BALR	05	I	750+S•(750-AB•300)
Branch and Link	RX	BAL	45	I	1275
Branch on Condition	RR	BCR	07	I	150+S•(750-AB•300)
					1125 (for serialization case)
Branch on Condition	RX	BC	47	I	150+S•(600-AB•300)

		For-	Mne-	Op-	Accura	асу
	Instruction Name	mat	monic	Code	Code	Formula and Comments
	Branch on Count	RR	BCTR	06	I	$525+S \cdot (600-AB \cdot 300)$ (for $R2 \neq 0$) 300 (for $R2=0$)
	Branch on Count	RX	BCT	46	I	$375 + S \cdot (600 - AB \cdot 300)$
	Branch on Index High	RS	вхн	86	I	750+S•(600-AB•300)
	Branch on Index Low or Equal	RS	BXLE	87	I	750+S•(600-AB•300)
I	Clear Page	c		D215	п	100000
	Compare	ם סס	CDR	10	11	100000
	Compare	RX	C	59		575 $600 \pm DW2 + 225 \pm X + 150$
	Compare (long)	RR	CDR	29		450 (for E1 - E2)
	compare (rong)		ODR	27		825 (for E1=22)
	Compare (long)	RX	CD	69		$825 + DW2 \cdot 225 + X \cdot 150$ (for E1=E2)
						$1200 + DW2 \cdot 225 + X \cdot 150$ (for E1 \neq E2)
	Compare (short)	RR	CER	39		450 (for E1=E2)
						1125 (for E1≠E2)
	Compare (short)	RX	CE	79		$825 + DW2 \cdot 225 + X \cdot 150$ (for E1=E2)
						1500+DW2•225+X•150 (for E1 <e2)< td=""></e2)<>
						1350+DW2•225+X•150 (for E1>E2)
	Compare and Swap	RS	CS	BA		975 (for OP1≠OP2)
						1050 (for OP1=OP2)
	Compare Decimal	SS	СР	F9		$1275 + (DW1 + DW2) \cdot 225$ (for $L1 \le 8, L2 \le 8$)
						$1800+(DW1+EDW2+DW2) \cdot 225 \text{ (for } L1 \le 8, L2 > 8)$
						$1950+(DW1+EDW1+DW2) \cdot 225 \text{ (for } L1>8, L2\leq8)$
		DC	CDC			$2025+(DW1+EDW1+DW2+EDW2) \cdot 225$ (for L1>8, L2>8)
	Compare Double and Swap	RS	CDS	вв		$1650 \text{ (for OP1 \neq OP2)}$
	Compare Halfword	DV	СЦ	40		18/5 (10f OP1=OP2) 600 + DW2 + 225 + V + 150
	Compare Logical	RA DD		49		375
	Compare Logical	RX		55		$600 \pm DW_{2} \cdot 225 \pm X \cdot 150$
	Compare Logical (character)	SS		D5		975 (for I < 8)
		55	020	20		$600 + L8 \cdot 525 + MA \cdot (floor of L/8) \cdot 225 (for L>8)$
	Compare Logical (immediate)	SI	CLI	95		600
	Compare Logical Characters					
	under Mask	RS	CLM	BD	I	825+CM•150+DW2•225+M•75
	Compare Logical Long	RR	CLCL	0F	II	675+C256•6750+L8C•675+MA•L8C•225
•						(for C256>0 and F256=0)
						$675 + F256 \cdot (7650 + (L1 > 0) \cdot 300) + L8F \cdot 600$
						(for C256=0 and F256>0)
						$-2225 + F256(7650 + (L1 > L2) \cdot 300) + C256 \cdot 6750$
						+L8F•600+L8C•675+MA•L8C•225
	Comment De se	DC	CTD	DO		(for C256>0 and F256>0) $= 5225 (-2075)$
	Connect Page	KS DV	CVP	BU		5325-C+3075
	Convert to Decimal	RA DY		4F 4F		1275 + NDD2 + 750 + DW2 + 223 + X + 150
	Convert to Decimal	KA	CVD	40		1273•(1+1411D)+K•130+Dw2•225+X•130
	Deconfigure Page	S	DĘP	B21B		13650
	Diagnose			83		
	Disconnect Page	S	DCTP	B21C		3375+C•14425
	Divide	RR	DR	1D		6525+N2•150+N1•375
	Divide	RX	D	5D		7200+N1•375+N2•150+DW2•225+X•150
	Divide (long)	RR	DDR	2D		10050+PN•225
	Divide (long)	RX	DD	6D		10575+PN•225+DW2•225+X•150
	Divide (short)	RR	DER	3D		5550+PN+225
	Divide (short)	КX	DE	7 D		6U/S+(FN+DW2)+225+X+150

.

	For-	Mne-	Op-	Accura	cy
Instruction Name	mat	monic	Code	Code	Formula and Comments
Divide Decimal	SS	DP	FD	Π	$7350+(L1>8) \cdot 1500-(L1=16) \cdot 300$ -(NDD1=0) \cdot 150+DWR \cdot 225 (for NDD1 <ndd2) 11800+(NDD1-NDD2) \cdot 1350 (for NDD1 \ge NDD2 and L1 \le 8) 13300+(1350+LC \cdot 150) \cdot (NDD1-NDD2) +(NDD1>15) \cdot 825-(L1=16) \cdot 300+(DWQ+DWR) \cdot 225 (for NDD1 \ge NDD2 and L1>8)</ndd2)
					Note: For cases with NDD1 \geq NDD2 times assume equal probability of digits 0 through 9.
Edit	SS	ED	DE	II	975+(Cond.Code≠0)•225+Sum(EB+SA) (Where the sum is taken over all pattern characters. Refer to Figure A-3.)
Edit and Mark	SS	EDMK	DF	II	1425+(Cond.Code≠0)•225+Sum(EB+SA+EMK) (Where the sum is taken over all pattern characters. Refer to Figure A-3.)
Exclusive OR	RR	XR	17		375
Exclusive OR	RX	х	57		600+DW2•225+X•150
Exclusive OR (character)	SS	xc	D7		1275+DW1•375+DW2•225 (for L≤8) 1200+L8•825+LZ8•75+MA•(floor of L/8)•225 (for L>8) 900+L•825+LZ•75 (for operands overlap) 1350+L8•225 [for L>8 and ADDR(OP1)=ADDR(OP2)]
Exclusive OR (immediate)	SI	XI	97		825
Execute	RX	EX	44		1125+SI+R•3300+X•150 (for R1=0) 1425+SI+R•3300+X•150 (for R1≠0)
Halve (long)	RR	HDR	24		1125+PN•225
Halve (short)	RR	HER	34		975+PN•225
Insert Character	RX	IC	43		375+X•150
Insert Characters under Mask	RS	ICM	BF	I	900+DW2•225+IM•150
Insert Page Bits	RS	IPB	B4		2250+C•525
Insert PSW Key	S	IPK	B20B		1350
Invalidate Page Table Entry	S	IPTE	B221		6075 (2K Page Size) 6375 (4K Page Size)
Insert Storage Key	RR	ISK	09		1425 (for BC, S/370 mode)
					1575 (for EC, S/370 mode)
					1725 (for BC, ECPS:VSE mode)
					2475 (for EC, ECPS:VSE mode)
Load	RR	LR	18		300
Load	RX	L	58		375+DW2•225+X•150
Load (long)	RR	LDR	28		300
Load (long)	RX	LD	68		375+DW2•225+X•150
Load (short)	RR	LER	38		300
Load (short)	RX	LE	78		375+DW2•225+X•150
Load Address	RX	LA	41		300+X•150
Load and Test	RR	LTR	12		375
Load and Test (long)	RR	LTDR	22		525
Load and Test (short)	RR	LTER	32		525
Load Complement	RR	LCR	13		375
Load Complement (long)	RR	LCDR	23		675
Load Complement (short)	RR	LCER	33		675

		For-	Mne-	Op-	Accura	сy
	Instruction Name	mat	monic	Code	Code	Formula and Comments
	Load Control	RS	LCTL	B7		12750+PU•5400 (for S/370 mode, load all Regs)
						12000 (for ECPS:VSE, load all Regs)
						5025+PU•5400 (for S/370 mode, load Regs 0,1)
						3825 (for ECPS:VSE, load Regs 0,1)
	Load Frame Index	RS	LFI	B 8		2175
	Load Halfword	RX	LH	48		375+DW2•225+X•150
	Load Multiple	RS	LM	98		150+N•225
	Load Negative	RR	LNR	11		525-N2•150
	Load Negative (long)	RR	LNDR	21		525
	Load Negative (short)	RR	LNER	31		675
	Load Positive	RR	LPR	10		375+N2•225
	Load Positive (long)	RR	LPDR	20		675
	Load Positive (short)	RR	LPER	30		675
	Load PSW	S	LPSW	82		5250 (for BC/BC)
						5550 (for EC/BC)
						6150 (for BC-EC/EC)
	Load Real Address	RX	LRA	B 1		4800 (for 1M segment)
						5100 (for 64K segment)
	Load Rounded (extended to					
	long)	RR	LRDR	25		750
	Load Rounded (long to short)	RR	LRER	35		900
	Make Addressable	S	MAD	B21D		3075-A•150
	Make Unaddressable	S	MUN	B21E		2925+A•450
	Monitor Call	SI	MC	AF		13875
						1050 (for NO-OP case)
	Move (character)	SS	MVC	D2		$825 + DW1 \cdot 225 + DW2 \cdot 150$ (for $L \le 8$)
						675+L8•450+MA•(floor of L/8)•225 (for L>8)
_						900+L•450 (for L ≤ 8 and operands overlap)
						975+L•450 (for L>8 and operands overlap, $A1 \neq A2+1$)
						1725+L8•225(for L>8 and A1=A2+1)
	Move (immediate)	SI	MVI	92		375
	Move Inverse	SS	MVCIN	E8		750+L•750
	Move Long	RR	MVCL	0E	11	375+M256•6863+P256•7688-MEOB•5738+L8M•450
						+L8P•225
1						(for M256>0 and P256>0)
I						375+M256•6863+MEOB•150+L8M•450
						(for M256>0 and P256=0)
						-600+P256•7688+L8P•225
						(for M256=0 and P256>0)
	Move Numerics	SS	MVN	D1		$1800 + DW2 \cdot 225 + DW1 \cdot 375$ (for $L \le 8$)
						(L8+1)•1125+MA•(floor of L/8)•225 (for L>8)
						$825+L \cdot 1125+DW1 \cdot 150$ (for L ≤ 8 and operands overlap)
						825+L•1125 (for L>8 and operands overlap)

•

Instruction Name nunt Code Code Formula and Comment Move with Offset SS MVO F1 2400+(A2 <a1+l1):375+(2-dw1+dw2):225 ((or L15&L25)) Move with Offset SS MVO F1 2400+(A2<a1+l1):375+(2-dw1+dw2):225 ((or L15&L25)) SS0+(A2<a1+l1):375+(2-dw1+dw2):225 ((or L15&L25)) 3000+(A2<a1+l1):375+(2-dw1+dw2):225 ((or L15&L25)) 3000+(A2<a1+l1):375+(2-dw1+dw2):225 ((or L15&L25)) Move Zones SS MVZ D3 Stabub/225 ((or or unspin)) Matipity R MR IC Stabu-123+DW2-225 ((or L28)) ((1+1):100+24+DW2+2DW2):225 + (1-28) Multipity RX M SC 3000+N2-450+DW2-225 ((or L28)) Multipity RX M SC 3000+N2-450+DW2-225 ((or L28)) Multipity RX M SC 3000+N2-450+DW2-225 + X-150 Multipity (cong to extended) RR MXR S2 Stabu-PN-75 Multipity (cong) RX MD C 3825+PN-75 Multipity (cong) RX MD C 3825+PN-25 Multipity (cong) RX</a1+l1):375+(2-dw1+dw2):225 </a1+l1):375+(2-dw1+dw2):225 </a1+l1):375+(2-dw1+dw2):225 </a1+l1):375+(2-dw1+dw2):225 </a1+l1):375+(2-dw1+dw2):225 		For-	Mne-	Op-	Accuracy
Move with Offset SS MVO F1 2400+(A2 <a1+l1)-1375+(2-dw1+dw2)+225 (or L15K1.25K) Move with Offset SS MVO F1 2550+(A2<a1+l1)-1375+(2-dw1+dw2+edw2)+225 (or L15K1.25K) Move Zones SS MVZ D3 3000+(A2<a1+l1)-1375+(2-edw1+2-dw1+dw2)+225 (or L15K1.25K) Move Zones SS MVZ D3 1800+DW1-375+DW2-225 (or L5K) Multiply RK MR IC 1500+N1-105+U2-25 (or urbanping case) Multiply RK MR IC 1500+N2-25 (or U-5K) Multiply (carended) RK MXR 2C 15525+PN-75 Multiply (cong cestended) RK MXR 2C 4550+PN-75 Multiply (long) RX MD 2C 4550+PN-75 Multiply (long) RX MD 2C 4550+PN-225 Multiply</a1+l1)-1375+(2-edw1+2-dw1+dw2)+225 </a1+l1)-1375+(2-dw1+dw2+edw2)+225 </a1+l1)-1375+(2-dw1+dw2)+225 	Instruction Name	mat	monic	Code	Code Formula and Comments
Move with Offset SS MVD FI 2400+(A2 <a1+l1):175+(2-dw1+dw2):225 (for L1≤8.1.2≤5) Move with Offset SS MVD FI (for L1≤8.1.2≤5) Move Zones SS MVZ D3 (for L1≤8.1.2≤5) Move Zones SS MVZ D3 (for L1>8.1.2≤5) Move Zones SS MVZ D3 (for L1>8.1.2≤8) Multiply RR MR IC (for L1>8.1.2≤8) Multiply RR MR IC (for L1>8.1.2≤8) Multiply (fong to extended) RR MR 26 (for L3%1.255 And operands overlap) Multiply (fong to extended) RR MR 26 (for S152+PN-15) Multiply (fong to extended) RR MRE 3C (for S152+PN-15) Multiply (fong to extended) RR MRE 3C (for S152+PN-15) Multiply (fong to extended) RR MRE 3C (for ND1+02) Multiply (fong to extended) RR MRE 3C (for ND1+02) Multiply (fong to extend</a1+l1):175+(2-dw1+dw2):225 					
(or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 8,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § N, S,L2 § N) (or L1 § 1, S,L2 § N) (or L1 § N, S,L2 § N) (or L1 § N) (or L1 § N, S,L2 § N) (or L1 § N,L2 § N) (or ND) ND (or ND) (or ND) R M2 \$ C (lipt() (or ND) ND (li	Move with Offset	SS	MVO	F1	2400+(A2 <a1+l1)•375+(2•dw1+dw2)•225< td=""></a1+l1)•375+(2•dw1+dw2)•225<>
Image: Second					$(for L1 \le 8, L2 \le 8)$
(or 11 § k.12.8) 3600+(A2 < A1+L1). 375+(2-DW1+2.DW1+DW2).225					$2850+(A2$
3600+(A2 <a1+l1)-375+(2+dw1+2+dw1+dw2)+225 </a1+l1)-375+(2+dw1+2+dw1+dw2)+225 (for L1>8,L2<8) (gout-12 <a1+l1)-375 </a1+l1)-375 +(2+DW1+2-DW1+DW2+EDW2)-225 (for L1>8,L2>8) 3900+min(L1,L2)+1050+(L1>L2)+251+ 					(for L1≤8,L2>8)
Instrument (or L1>8,12.8)(or L1>8,12.8)900+(A2 <a1+l1)-375 </a1+l1)-375 +(2-EDW1+2.DW1+DW2+EDW2)-225 (or L1>8,12.2)Move ZonesSSMVZD3MultiplyRRMRMultiplyRRMRMultiplyRRMRMultiply (extended)RRMRMultiply (extended)RRMRMultiply (on to extended)RRMXRMultiply (on to extended)RRMXRMultiply (on to for long)RRMXDMultiply (ong to extended)RRMXDMultiply (ong to extended)RRMXDMultiply (ong to extended)RRMZDMultiply (ong to extended)RRMZDMultiply (ong to extended)RRMZDMultiply (ong)RRMDDMultiply (ong)RRMDDMultiply (ong)RXMEMultiply (ong)RRMDDMultiply (ong)RRMDDMultiply (ong)RRMDDMultiply (ong)RRMDDMultiply (ong)RRMDRMultiply (ong) <td></td> <td></td> <td></td> <td></td> <td>$3600 + (A2 < A1 + L1) \cdot 375 + (2 \cdot EDW1 + 2 \cdot DW1 + DW2) \cdot 225$</td>					$3600 + (A2 < A1 + L1) \cdot 375 + (2 \cdot EDW1 + 2 \cdot DW1 + DW2) \cdot 225$
Move Zones SS MVZ D3 1800+(A2 <a1+l1)-375 +(2-EDW1+2+DW1+DW2+EDW2)-225 ((Gr L1>&L2>S) Move Zones SS MVZ D3 1800+DW1-375+DW2-222 (Gr L2>S) (L1>L12>B)-225 (Gr L2>B) Mutiply RR MR IC 125+DW1-150 (L1>L2>S) Mutiply RR MR IC 3000+X2450 Mutiply (Grad Carendo) RR MR 26 3000+X2450 Mutiply (Grad Carendo) RR MXD 27 5850+PN-75 Mutiply (Grad Carendo) RR MXD 27 5850+PN-75 Mutiply (Grad Carendo) RR MXD 27 5850+PN-75 Mutiply (Grad Carendo) RR MZD 27 5850+PN-75 Mutiply (Grad Carendo) RX MZD 27 5850+PN-75 Mutiply (Grad Carendo) RX MZD 27 5850+PN-75 Mutiply (Grad Carendo) RX MZD 27 5850+PN-225 Mutiply (Grad Carendo) RX MZD 27 5850+PN-75 Mutiply (Grad Carendo) RX<td></td><td></td><td></td><td></td><td>(for L1>8,L2≤8)</td></a1+l1)-375 					(for L1>8,L2≤8)
Move Zones SS MVZ D3 1800+PW1+2108)+(12-28)-225 (for vertapping case) Move Zones SS MVZ D3 1800+PW1+2108)+(12-28)-225 (for vertapping case) Multiply RR MR 1C 3000+N2-450 Multiply RX M 5C 3000+N2-450 Multiply (catended) RR MX 26 15525+PN-75 Multiply (catended) RR MXD 67 6375+PN-75 Multiply (catended) RX MZD 67 6375+					$3900+(A2$
Move Zones SS MVZ D3 (for L1>2.8.1-2.8) 3000+min(L1,L2)-1050+(L1>L2)-525+ (L1>2.8)[-225 (for overlapping case) Multiply RR MR D3 (B00+DW1-375+DW2-225 (for L28) (L3+L1)-1125+MA-(floor JL/8)-225) 825+L-1125+DW1-150 (L58 and operands overlap) 825+L-1125 (L58 and operands overlap) 825+L-1125 (L58 and operands overlap) 825+L-1125 (L58 and operands overlap) Multiply RR MR 1C 3000+X-450 1050 (for VP=0) Multiply (castended) RR MXD 27 5500+FN-75 Multiply (cont o toxended) RR MXD 67 6375+FN-75+DW2-225+X-150 Multiply (cont o long) RR MDR 3C 3825+FN-225 Multiply (long) RR MDR 2C 4875+FN-225 Multiply (long) RR MDR 2C 4875+FN-225 Multiply (long) RR MDR 2C 4875+FN-225 Multiply (long) RX MD 6C 5400+FN-225+MV2-25+X-150 Multiply (long) RX MD 6C 5400+FN-225+MV2-25+X-150 Multiply (long) RX MD 6C 5400+FN-225+MV2-25+X-150					+(2•EDW1+2•DW1+DW2+EDW2)•225
Move ZonesSSMVZD33000+min(1.1.2)-1050+(1.2)-815+(1.2)-					(for L1>8,L2>8)
Move Zones SS MZ D3 1800+DW1:375+UM2-225 (for L>28) (L3+H):1125+UA+(IDO of L/S):225 (for L>8) 225+L-1125 (L>8 and operands overlap) 825+L-1125 (L>8 and operands overlap) 825+L-1125 (L>8 and operands overlap) Multiply RR MR IC 3000+N2:450 1050 (for VP=0) Multiply (extended) RR MXD 26 3502+PN-75 Multiply (ong to extended) RR MXD 7 6352+PN-75 Multiply (short to long) RX MZ 300+PN-225+X-150 Multiply (ong) RR MER 3C 3825+PN-25 Multiply (ong) RR MER 3C 3825+PN-25 Multiply (ong) RR MDR 2C 4352+PN-225 Multiply (ong) RX MD 6C 5400+PN-225+DW2-225+X-150 Multiply Decimal SS MP FC 4352+DW2-225+(10+DW1+DW2-225+(10+DW1+D					$3900 + \min(L1, L2) \cdot 1050 + (L1 > L2) \cdot 525 +$
Move Zones SS MVZ D3 1800+DW1.075+DW2-225 (for L_S) (L8+1)-1125+DW1-150 (L_S) and operands overlap) 825+L-1125 (L_S) and operands overlap) Multiply RR MR IC 3000+N2.450 1050 (for VP=0) Multiply (cstended) RR MX 26 1552+PN.75 Multiply (cstended) RR MXR 27 5850+PN.75 Multiply (store to long) RR MXR 26 3252+PN.75 Multiply (store to long) RR MRR 3C 3325+PN.252 Multiply (short to long) RR MRR 3C 3325+PN.225 Multiply (long) RX MD 6C 5400+PN.225+DW2.225+X.150 Multiply Decimal SS MP FC 4755+MD2.1050+CL2.525+X.150 Multiply Lange RX MD 6C 5400+PN.225+DW1.202					[(L1>L2+8)+(L1>8)+(L2>8)]•225 (for overlapping case)
Image: Second	Move Zones	SS	MVZ	D3	$1800 + DW1 \cdot 375 + DW2 \cdot 225$ (for $L \le 8$)
Multiply RR MR IC 3604+N2-450 1050 (for VP=0) Multiply (extended) RR MX SC 3000+N2-450 1050 (for VP=0) Multiply (extended) RR MXR SC 3000+N2-450-DW2-225+X-150 Multiply (extended) RR MXR SC 3000+N2-450-DW2-225+X-150 Multiply (extended) RR MXD S7 S850+FN-75 Multiply (forg to extended) RR MXD G7 G375+FN-75+DW2-225+X-150 Multiply (forg) RR MZD C 4350+FN-225 Multiply (forg) RR MD C 4875+FN-225 Multiply (forg) RX MD C 4875+FN-225 Multiply Decimal SS MD FC 4875+FN-225 Multiply Decimal SS MD FC 4875+FND2-15DH1-DU1-D1+2+EDW1)-225 (for NDD1>5 for NDD1=ND1 ST FD0+1-DW1+2+EDW1)-225 (for NDD1+5 Multiply Decimal SS MD FC 4875+KD02-150+IE2+FD02+225+K1-210 Multiply Decimal					(L8+1)•1125+MA•(floor of L/8)•225 (for L>8)
Multiply RR RR IC 3600+N2-450 Multiply RX M SC 3000+N2-450+DW2-225+X+150 Multiply (cented) RR MX 26 15525+PN-75 Multiply (long to extended) RR MXD 27 850+PN-75 Multiply (long to extended) RR MXD 67 6375+PN-757+DW2-225+X-150 Multiply (short to long) RX ME 7C 4350+PN-225 Multiply (long) RR MDR 2C 4875+PN-225 Multiply (long) RR MDR 2C 4875+PN-225 Multiply (long) RR MDR 2C 4875+PN-225 Multiply Decimal RS MP FC 4875-(NDD2-0) STOO+(FDW1+PW2)-225+U12-25+U12-25+U12-25+U12-25+U12-25 (for NDD1>51 and NDD2-0) 5100+CCL2-525+(3-DW1+DW1)-225 STOO+(FDW1+DW2)-225+(13.03+FDW1-DW1+2-EDW1)-225 (for NDD1=15 and NDD2-0) 11625+CL2+52+U12+5+DW1-24+EDW1)-225 STOO+(FDW1+DW2)-225+(13.04+DW2)-225+(13.04+DW2)-225+(12.05+W1+DW2)-225 (for NDD1=15 and NDD2-0) 11625+CL2+525+(3-DW1+DW1)-225 STOO					$825+L \cdot 1125+DW1 \cdot 150$ (L ≤ 8 and operands overlap)
Multiply RR RR RR C 3600+N2450 1050 (for VP=0) 3000+N2450+DW2-225+X:150 Multiply (extended) RR MXR 26 15525+FN-75 Multiply (not o extended) RR MXD 67 6375+FN-75+DW2-225+X:150 Multiply (not o extended) RR MXD 67 6375+FN-75+DW2-225+X:150 Multiply (short to long) RR MER 3C 3825+FN-225 Multiply (long) RR MDD 2C 4375+FN-225 Multiply (long) RX MD 6C 5400+FN-225+DW2-225+X:150 Multiply Decimal SS MP FC 4375+FN-225 Multiply Decimal SS MP FC 4375+FND20+25 (for ND1=0) 5100+CLZ+525+(3-DW1+DW1+2+DW1)+225 (for ND1=0) ST00+(FDW1+DW2)+22D (for ND1=15 and NDD2-0)					$825+L \cdot 1125$ (L>8 and operands overlap)
Instruct Instruct Instruct Multiply (extended) RR MXR 26 3900+N2.450+DW2.225+X.150 Multiply (long to extended) RR MXDR 27 5850+PN.75 Multiply (long to extended) RR MXDR 27 5850+PN.75 Multiply (short to long) RX MER 3 6375+PN.75+DW2.225+X.150 Multiply (short to long) RX MER 3 6352+PN.225 Multiply (long) RR MDR 2 4875+PN.225 Multiply (long) RR MDR 2 4875+PN.225 Multiply Decimal SS MD 6C 5400+PN.225+DW2.225+X.150 Multiply Decimal SS MD FC 4875-(NDD2=0).150+ (3.DW1+DW2).225+(1.33+FDW1-DW1+2.EDW1).225 (for NDD1=)51 and NDD2=0) 5100+(FDW1+DW2).225+(1.33+FDW1-DW1+2.EDW1).225 (for NDD1 SIGO+(FDW1+DW2).225+(1.50+(DW1+DW2).225+(1.158)+(3.4FDW1-DW1+2.EDW1).225 (for NDD1=15 and NDD2=0) I1625-CL2.525+(3.0W1+DW2).225 (for NDD1=ND2=15) 14275+300-NDD2+(FDW1+DW2).225 (for NDD1=ND2=15) 14275+300-N	Multiply	RR	MR	1C	3600+N2•450
Multiply Multiply (extended) RX MX SC 3900+N2+250+DW2+225+X+150 Multiply (long to extended) RR MXDR 27 5850+PN-75 Multiply (long to extended) RX MXD 67 6375+PN-75 Multiply (long to extended) RX MXD 67 6375+PN-75+DW2-225+X-150 Multiply (long) RX ME 7C 4350+PN-225 Multiply (long) RX ME 7C 4350+PN-225+DW2-225+X-150 Multiply (long) RX MD 6C 5400+PN-225+DW2-225+X-150 Multiply Decimal SS MP FC 4875-(NDD2-0) Multiply Decimal SS MP FC 4875-(NDD2-0) S100-CLZ-s25+G100+DW1+DW1+2-EDW1+2-EDW1)+225 (for NDD1+DW2+2-EDW1+2-DW1)+225 (for NDD1+DW2+2-EDW1+2-DW1)-225 (for NDD1-S1 and NDD2=0) 5100-CLZ-s25+(3-DW1-DW1+2-EDW1)+225 (for NDD1-S1 and NDD2-0) S115+NDD2-150-CLZ-s25+(3-DW1+DW2)-225 +(L1>8)-(3+FDW1-DW1+2-EDW1)+225 (for NDD1-S1 and NDD2-0) S1162-CLZ-s25+(3-DW1-DW2+2-EDW1+2-DW1)+225 (for NDD1-ND1=15 and NDD2-0) 11625-CLZ-s25+(3-DW1+DW2)-225					1050 (for VP=0)
Multiply (extended) RR MXR 26 15525+FN-75 Multiply (long to extended) RR MXD 77 5850+FN-75 Multiply (long to extended) RX MXD 67 6375+FN-75+DW2-225+X-150 Multiply (short to long) RR MER 3C 3825+FN-225 Multiply (long) RX ME 7C 4350+FN-225+DW2-225+X-150 Multiply (long) RX MD 6C 5400+FN-225+DW2-225+X-150 Multiply (long) RX MD 6C 5400+FN-225+DW2-225+X-150 Multiply Decimal SS MP FC 4875+(ND2=0)-150+ (long) RX MD 6C 500+(FDW1+DW2+2:EDW1)-01+2:EDW1)-225 (for NDD1=0) 5100-CLZ-525+(1.03+FDW1-DW1+2:EDW1)-225 (for NDD1=15 and ND2=0) S100-CLZ-525+(3-DW1+DW2)-225 +(L1>8)+(3+FDW1-DW1+2:EDW1)-225 (for NDD1=15 and ND2>20) I650+WP-600+(FDW1+DW2+2:EDW1+2:DW1)+225 (for NDD1=15 and ND2>20) 1162:S2+(3:0+ND1+DW2)-225 (for NDD1=15 and ND2>0 6150+WP-600+(FDW1+DW2+2:EDW1+2:DW1)+225 (for NDD1=15 and ND2>0) I619+WP-600+(FD	Multiply	RX	М	5C	$3900 + N2 \cdot 450 + DW2 \cdot 225 + X \cdot 150$
Multiply (long to extended) RR MXDR 27 \$\$850+PN-75 Multiply (long to extended) RX MXD 67 6375+PN-75+DW2-225+X-150 Multiply (short to long) RX ME 3C 4350+PN-225 Multiply (short to long) RX ME 7C 4350+PN-225+DW2-225+X-150 Multiply (long) RX MD 6C \$400+PN-225+DW2-225+X-150 Multiply Decimal SS MP FC 4875-(ND2=0)-150+ Multiply Decimal SS MP FC 4875-(ND2=0)-150+ (3-DW1+DW2)+225+(1.33+FDW1-DW1+2-EDW1)+225 (for ND01=0) 5700+(FDW1+DW2+2-EDW1+2-DW1)+225 (for NDD1>15 and NDD2=0) \$100-CLZ+525+(3-DW1+DW1)+225 (for NDD1=S) (for NDD1 <s15 and="" ndd2="0)</td"> 8175+NDD2-150-CLZ+525+(3-DW1+DW1)+225 (for NDD1<s15 and="" ndd2="0)</td"> 11625-CLZ+525+(3-DW1+DW1)+225 (for NDD1<s15 and="" ndd2="0)</td"> 11625-CLZ+525+(3-DW1+DW1)+225 (for NDD1<s15 and="" ndd2<0)<="" td=""> 11625-CLZ+525+(3-DW1+DW1)+225 (for NDD1=NDW1+2+EDW1)+225 (for NDD1=NDM2+51) 14275+300-ND2+(FDW1+DW2+2+EDW1)+225 (for NDD1=NDM2+51) Multiply HaffwordR</s15></s15></s15></s15>	Multiply (extended)	RR	MXR	26	15525+PN•75
Multiply (long to extended) RX MXD 67 6375+PN-75+DW2+225+X+150 Multiply (short to long) RX ME 3C 3825+PN-225 Multiply (long) RX ME 7C 4350+PN-225+DW2+225+X+150 Multiply (long) RX MD 6C 5400+PN-225+DW2+225+X+150 Multiply (long) RX MD 6C 5400+PN-225+DW2+225+X+150 Multiply Decimal SS MP FC 4875-(NDD2=0)+150+ Multiply Decimal SS MP FC (3-DW1+DW2)+225+(1.33+FDW1-DW1+2+EDW1)+225	Multiply (long to extended)	RR	MXDR	27	5850+PN•75
Multiply (short to long) RR MER 3C 3825+PN-225 Multiply (long) RX MD 7C 4350+PN-225+DW2-225+X-150 Multiply (long) RX MD 6C 5400+PN-225+DW2-225+X-150 Multiply (long) RX MD 6C 5400+PN-225+L03+2DW1+20W1+2.EDW1+2.EDW1)-225 Multiply Decimal SS MP FC 4875-(NDZ=0)-150+ Multiply Decimal SS MP FC 4875-(NDZ=0)-150+ Som DD1=0 5700+(FDW1+DW2+2.EDW1+2.EDW1)-225 (for NDD1>15 and NDD2=0) 5100-CLZ-525+(3.0W1+DW2)-225 Som DD1 SIG NDD SIG NDD SIG Multiply Halfword RX MH 4C 3075+(NZ+X)-150 11625-(LZ+32)+(X-10W1+2)-225 Multiply Halfword RX MH 4C 3075+(NZ+X)-150 060+(DW1+DW2+2.EDW1+2.EDW1)-225 Multiply Halfword SS OC D6 1275+DW1-375+DW1-225 (for L≤8) Multiply Halfword RX MH 4C 3075+(N2+X)-150+(N2+2)-EDW1+2.EDW1)-225 Multiply Halfword	Multiply (long to extended)	RX	MXD	67	6375+PN•75+DW2•225+X•150
Multiply (short to long) RX ME 7C 4350+PN-225+DW2+225+X+150 Multiply (long) RX MD 6C 5400+PN-225+DW2+225+X+150 Multiply Decimal SS MP FC 4875+(NDD2=0)+150+ Multiply Decimal SS MP FC 4875-(NDD2=0)+150+ (3-DW1+DW2)+225+(1.33+FDW1-DW1+2+EDW1)+225 (for NDD1=0) 5700+(FDW1+DW2)+225+(1.33+FDW1-DW1+2+EDW1)+225 (for NDD1>15 and NDD2=0) 5100-CLZ+525+(3-DW1+DW2)+225 +(L1>8)+(3+FDW1-DW1+2+EDW1)+225 (for NDD1<5 15 and NDD2=0)	Multiply (short to long)	RR	MER	3C	3825+PN•225
Multiply (long) RR MDR 2C 4875+PN-225 Multiply (long) RX MD 6C 5400+PN-225+X-150 Multiply Decimal SS MP FC 4875-(NDD2=0)-150+ (3-DW1+DW2)-225+(1.33+FDW1-DW1+2-EDW1)-225 (for NDD1=0) 5700+(FDW1+DW2)-225+(1.33+FDW1-DW1+2-EDW1)-225 (for NDD1=0) 5700+(FDW1+DW2)-225+(3-DW1+DW2)-225 (for NDD1=15 and NDD2=0) 8175+ND2150-CL2-525+(3-DW1+DW2)-225 (for NDD1<51 and NDD2=0)	Multiply (short to long)	RX	ME	7C	4350+PN•225+DW2•225+X•150
Multiply (long) RX MD 6C 5400+PN-225+DW2-225+X×150 Multiply Decimal SS MP FC 4875-(NDD2=0)+150+ (3-DW1+DW2)+225+(1.33+FDW1-DW1+2*EDW1)*225 (for NDD1=0) 5700+(FDW1+DW2)*225+(1.33+FDW1-DW1)*225 (for NDD1=0) 5700+(FDW1+DW2)*25(3+FDW1-DW1)*225 (for NDD1=15 and NDD2=0) 8175+NDD2*150-CLZ*525+(3-DW1+DW2)*225 +(L1>8)*(3,67+FDW1-DW1)*225 (for NDD1<15 and NDD2=0)	Multiply (long)	RR	MDR	2C	4875+PN•225
Multiply Decimal SS MP FC 4875-(NDD2=0).150+ (3-DW1+DW2)-225+(1.33+FDW1-DW1+2•EDW1)+225 (for NDD1=0) 5706+(FDW1+DW2)+22EDW1+2•EDW1+2•EDW1)+225 (for NDD1>15 and NDD2=0) 5100-CL2*525+(3-DW1+DW2)+225 +(L1>8)•(3.67+FDW1-DW1+2•EDW1)+225 (for NDD15] is and NDD2=0) 8175+NDD2•150-CL2*525+(3-DW1+DW2)+225 +(L1>8)•(3+FDW1-DW1+2•EDW1)+225 (for NDD1<15 and NDD2=0)	Multiply (long)	RX	MD	6C	5400+PN•225+DW2•225+X•150
(3+DW1+DW2)+225+(1.33+FDW1-DW1+2*EDW1)*225 (for NDD1=0) 5700+(FDW1+DW2+2*EDW1+2*DW1)*225 (for NDD1=5) 5700-CLZ+525+(3+DW1+DW2)*225 +(L1>8)*(3-67+FDW1-DW1+2*EDW1)*225 (for NDD1≤15 and NDD2=0) 8175+NDD2*150-CLZ*525+(3+DW1+DW2)*225 +(L1>8)*(3+FDW1-DW1+2*EDW1)*225 (for NDD1<15 and NDD2>0) 11625-CLZ*525+(3+DW1+DW2)*225 +(L1>8)*(3+FDW1-DW1+2*EDW1)*225 (for NDD1<15 and NDD2>0) 11625-CLZ*525+(3+DW1+DW2)*225 +(L1>8)*(3+FDW1-DW1+2*EDW1)*225 (for NDD1=15 and NDD2>1) 18600-CLZ*525+(3+DW1+DW2)*225 (for NDD1=15 and NDD2>25) 18600-CLZ*525+(3+DW1+DW2)*225 (for NDD1=15 and NDD2>25) 18600-CLZ*525+(3+DW1+DW2)*225 (for NDD1=15 and NDD2>1) 18600-CLZ*525+(3+DW1+DW2)*225 (for NDD1=15 and NDD2>25) 18600-CLZ*525+(3+DW1+DW2)*225 (for NDD1=15 and NDD2>25) 18600-CLZ*525+(3+DW1+DW2)*225 (for NDD1=29 and NDD2=15) 14275+300*ND2+(FDW1+DW2+2*EDW1+2*DW1)*225 (for NDD1=29 and NDD2=1) Multiply Halfword RX MH 4C 3075+(N2+X)*150+DW2*225 (for ND1=29 (for ND1=29) (for ND1=29) OR RX 0 56 (for 0+DW2*225+X*150 OR RX 0 56 (for DW1*375+DW1*25 (for L≤8) 900+L *825+LZ8*75+MA*(floor of L/8)*225 (for L>8) 900+L *825+LZ8*75+MA*(floor of L/8)*225 (for L>8)	Multiply Decimal	SS	MP	FC	$4875 - (NDD2 = 0) \cdot 150 +$
(for NDD1=0) 5700+(FDW1+DW2+2*DW1)+2*DW1)•225 (for NDD1>15 and NDD2=0) 5100-CLZ*525+(3*DW1+DW2)•225 +(L1>8)•(3.67+FDW1-DW1+2*EDW1)•225 (for NDD1≤15 and NDD2=0) 8175+NDD2•150-CLZ*525+(3*DW1+DW2)•225 +(L1>8)•(3.67+FDW1-DW1+2*EDW1)•225 (for NDD1≤15 and NDD2=0) 8175+NDD2•150-CLZ*525+(3*DW1+DW2)•225 +(L1>8)•(3+FDW1-DW1+2*EDW1)•225 (for NDD1=15 and NDD2<15)					$(3 \cdot DW1 + DW2) \cdot 225 + (1.33 + FDW1 - DW1 + 2 \cdot EDW1) \cdot 225$
5700+(FDW1+DW2+2·EDW1+2.DW1)·225 (for NDD1>15 and NDD2=0) 5100-CLZ·525+(3.DW1+DW2)·225 +(L1>8)·(3.67+FDW1-DW1+2.EDW1)·225 (for NDD1≤15 and NDD2=0) 8175+NDD2·150-CLZ·525+(3.DW1+DW2)·225 +(L1>8)·(3+FDW1-DW1+2.EDW1)·225 (for NDD1≤15 and NDD2>0) 11625-CLZ·525+(3.DW1+DW2)·225 +(L1>8)·(3+FDW1-DW1+2.EDW1)·225 (for NDD1=15 and NDD2>0) 11625-CLZ·525+(3.DW1+DW2)·225 +(L1>8)·(3+FDW1-DW1+2.EDW1)·225 (for NDD1=15 and NDD2<15)					(for NDD1=0)
(for NDD1>15 and NDD2=0) 5100-CLZ.525+(3.0W1+DW2).225 +(L1>8):(3.67+FDW1-DW1+2*EDW1).225 (for NDD1≤15 and NDD2=0) 8175+NDD2:150-CLZ.525+(3.0W1+DW2).225 +(L1>8):(3+FDW1-DW1+2*EDW1).225 (for NDD1<15 and NDD2>0) 11625-CLZ.525+(3.0W1+DW2).225 +(L1>8):(3+FDW1-DW1+2*EDW1).225 (for NDD1=15 and NDD2>0) 11625-CLZ.525+(3.0W1+DW2).225 +(L1>8):(3+FDW1-DW1-2*EDW1).225 (for NDD1=15 and NDD2=15) 18600-CLZ.525+(3.0W1+DW2).225 +(L1>8):(3+FDW1-DW1-2*EDW1).225 (for NDD1=NDD2=15) 14275+300*NDD2+(FDW1+DW2+2*EDW1+2*DW1).225 (for NDD1=15 and NDD2>0) 6150+VP*600+(FDW1+DW2+2*EDW1+2*DW1).225 (for NDD1=29 and NDD2=1) Multiply Halfword RX MR 4C 3075+(N2+X)*150+DW2*225 OR RR OR RX OR RX OR RX OR RX OK 600+L27525+(50FL28) 900+L8*25+L28*75+MA*(floor of L/8)*225 (for L>8) 900+L8*825+L28*75+MA*(floor of L/8)*225 (for L>8)					5700+(FDW1+DW2+2•EDW1+2•DW1)•225
S100-CL2.525+(3.DW1+DW2).225 + (L1>8).(3.67+FDW1-DW1+2.EDW1).225 (for NDD1≤15 and NDD2=0) 8175+NDD2.150-CL2.525+(3.DW1+DW2).225 + (L1>8).(3.4FDW1-DW1+2.EDW1).225 (for NDD1<15 and NDD2>0) 11625-CL2.525+(3.DW1+DW2).225 + (L1>8).(3+FDW1-DW1+2.EDW1).225 (for NDD1<15 and NDD2>0) 11625-CL2.525+(3.DW1+DW2).225 + (L1>8).(3+FDW1-DW1+2.EDW1).225 (for NDD1=15 and NDD2<15)					(for NDD1>15 and NDD2=0)
+(L1>8)•(3.67+FDW1-DW1+2*EDW1)•225 (for NDD1≤15 and NDD2=0) 8175+NDD2•150-CL2•525+(3•DW1+DW2)•225 +(L1>8)•(3+FDW1-DW1+2•EDW1)•225 (for NDD1<15 and NDD2>0) 11625-CL2•525+(3•DW1+DW2)•225 +(L1>8)•(3+FDW1-DW1+2•EDW1)•225 (for NDD1=15 and NDD2<15)					$5100-CLZ \cdot 525 + (3 \cdot DW1 + DW2) \cdot 225$
(for NDD1≤15 and NDD2=0) 8175+NDD2.150-CLZ.525+(3.DW1+DW2).225 +(L1>8).(3+FDW1-DW1+2.EDW1).225 (for NDD1<15 and NDD2>0) 11625-CLZ.525+(3.DW1+DW2).225 +(L1>8).(3+FDW1-DW1+2.EDW1).225 (for NDD1=15 and NDD2<15)					$+(L1>8) \cdot (3.67+FDW1-DW1+2 \cdot EDW1) \cdot 225$
81/3+NDD2+130-CL2+323+(3+DW1+DW2)+223 +(L1>8)+(3+FDW1-DW1+2+EDW1)+225 (for NDD1<15 and NDD2>0) 11625-CL2+325+(3+DW1+DW2)+225 +(L1>8)+(3+FDW1-DW1+2+EDW1)+225 (for NDD1=15 and NDD2<15)					$(101 \text{ NDD1} \le 15 \text{ and } \text{NDD2} = 0)$
if (L1>8) (3+FDW1-DW1+2+2EDW1)*223 (for NDD1<15 and NDD2>0) 11625-CL2.\$25+(3.0W1+DW2)*225 +(L1>8) (3+FDW1-DW1+2+EDW1)*225 (for NDD1=15 and NDD2<15)					$81/3 + NDD2 \cdot 130 - CL2 \cdot 323 + (3 \cdot DW1 + DW2) \cdot 223$
(i) NDD1<15 and NDD2>0 11625-CLZ.525+(3.DW1+DW2).225 +(L1>8).(3+FDW1-DW1+2.EDW1).225 (for NDD1=15 and NDD2<15)					$+(L1>0) \cdot (3+FDW1-DW1+2 \cdot EDW1) \cdot 223$ (for NDD1>15 and NDD2>0)
Hilds Oblest (0) Hild Hilds +(L1>8)•(3+FDW1-DW1+2•EDW1)•225 (for NDD1=15 and NDD2<15)					$11625-C1.7 \cdot 525 + (3 \cdot DW1 + DW2) \cdot 225$
(for NDD1=15 and NDD2<15)					$+(L_1>8) \cdot (3+FDW1-DW1+2 \cdot FDW1) \cdot 225$
$\begin{tabular}{l lllllllllllllllllllllllllllllllllll$					(for NDD1 = 15 and NDD2 < 15)
$\begin{array}{c c c c c c c c c c c c c c c c c c c $					$18600-CLZ \cdot 525 + (3 \cdot DW1 + DW2) \cdot 225$
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $					+(L1>8)•(3+FDW1-DW1-2•EDW1)•225
$\begin{array}{c c c c c c c c c c c c c c c c c c c $					(for NDD1=NDD2=15)
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $					14275+300•NDD2+(FDW1+DW2+2•EDW1+2•DW1)•225
$\begin{array}{c c c c c c c c c c c c c c c c c c c $					(for NDD1>15 and NDD2>0)
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $					6150+VP•600+(FDW1+DW2+2•EDW1+2•DW1)•225
Multiply HalfwordRXMH4C $3075+(N2+X)\cdot150+DW2\cdot225$ ORRROR16 375 ORRXO56 $600+DW2\cdot225+X\cdot150$ OR (character)SSOCD6 $1275+DW1\cdot375+DW1\cdot225$ (for L ≤ 8) $900+L8\cdot825+LZ8\cdot75+MA\cdot(floor of L/8)\cdot225$ (for L>8) $900+L \cdot825+LZ8\cdot75$ (for operands overlap)					(for NDD1=29 and NDD2=1)
OR RR OR 16 375 OR RX O 56 $600+DW2\cdot225+X\cdot150$ OR (character) SS OC D6 $1275+DW1\cdot375+DW1\cdot225$ (for L≤8) 900+L8*825+LZ8*75+MA*(floor of L/8)*225 (for L>8) 900+L *825+LZ8*75+MA*(floor of L/8)*225 (for L>8)	Multiply Halfword	RX	ΜН	4C	3075+(N2+X)•150+DW2•225
OR RX OK 10 $3/3$ OR RX O 56 $600+DW2 \cdot 225+X \cdot 150$ OR (character) SS OC D6 $1275+DW1 \cdot 375+DW1 \cdot 225$ (for L ≤ 8) 900+L8 \cdot 825+LZ8 \cdot 75+MA \cdot (floor of L/8) \cdot 225 (for L > 8) 900+L \cdot 825+LZ8 \cdot 75+MA \cdot (floor of L/8) \cdot 225 (for L > 8)	OP	סס	OP	16	275
NA O 50 $000+Dw2\cdot223+A\cdot130$ OR (character) SS OC D6 $1275+DW1\cdot375+DW1\cdot225$ (for L ≤ 8) $900+L8\cdot825+LZ8\cdot75+MA\cdot(floor of L/8)\cdot225$ (for L>8) $900+L \cdot825+LZ8\cdot75+MA\cdot(floor of L/8)\cdot225$ (for L>8)		RK DV		56	575 $600 \pm DW2 - 225 \pm Y - 150$
$\frac{1273 + D \times 1 \cdot 373 + D \times 1 \cdot 223}{101 L \le 8}$ $900 + L8 \cdot 825 + LZ8 \cdot 75 + MA \cdot (floor of L/8) \cdot 225 (for L>8)$ $900 + L \cdot 825 + L7 \cdot 75 (for operands overlap)$	OR (character)	лл 99		-00 -00	$1000 \pm D W 2 \cdot 223 \pm A \cdot 130$ $1275 \pm D W 1 \cdot 375 + D W 1 \cdot 225 (for 1 < 0)$
$900\pm1.825\pm1.7.75$ (for operands overlap)	ON (Character)	33	00	00	$1213 \pm 10113 \pm 1013 \pm 10112223$ (101 $1 \ge 0$) 900118.8251178.75 $\pm MA$. (floor of 1/9). 225 (for 1 ≤ 0)
					$900+L \cdot 825+LZ \cdot 75$ (for operands overlap)

	For-	Mne-	Op-	Accuracy
Instruction Name	mat	monic	Code	Code Formula and Comments
OR (immediate)	SI	OI	96	825
Pack	55	PACK	F2	$1725 \pm (DW1 \pm DW2) + 225 (for 1.1 < 8.1.2 < 8)$
Tack	33	IACK	1.7	$(DW_1 + DW_2) \cdot 223 (101 E1 \le 6, E2 \le 6)$ 2100+(DW_1 + EDW_1 + DW_2) \cdot 225 (for 1.1 > 8.1.2 < 8)
				$2100 + (DW1 + DW2 + FDW2) \cdot 225 (101 E1 > 8, E2 \le 8)$ $2400 + (DW1 + DW2 + FDW2) \cdot 225 (101 E1 < 8, E2 \le 8)$
				$3000 \pm (DW1 \pm DW1 \pm DW2 \pm DW2) + 225 (for L1 > 8 L2 > 8)$
Purge TLB	S	PTLB	B20D	6750
Reset Reference Bit	s	PPR	R213	$3225 \pm C_{\star} 150$ (for ECDS: VSE mode)
Reset Reference Br	5	KKD	D 215	1875 (for S/370 mode)
Retrieve Status and Page	22	DCD	20	13775 (101 5/ 5/0 mode)
Retrieve Status and Fage	33	KSI	108	133423
Set Clock	S	SCK	B204	7650 (for Enable ON)
				1950 (for Enable OFF)
Set Clock Comparator	S	SCKC	B206	5325
Set CPU Timer	S	SPT	B208	11175
Set Page Bits	RS	SPB	B5	3675
Set Program Mask	RR	SPM	04	750
Set PSW Key from Address	S	SPKA	B20A	2325
Set Storage Key	RR	SSK	08	7650 (for S/370 mode)
				4725 (for ECPS:VSE mode)
Set System Mask	S	SSM	80	1725 (for BC; S/370 mode or ECPS:VSE mode)
				5475 (for EC; ECPS:VSE mode)
				4875 (for EC; S/370 mode)
Shift and Round Decimal	SS	SRP	F0	3600-ZR•225 (for L \leq 8, shift left and amount < 16)
				3300-ZR \cdot 150 (for L \leq 8, shift right and amount $<$ 16)
				2400 (for $L \le 8$, shift left and amount ≥ 16)
				2175 (for $L \leq 8$, shift right and amount>16)
				2400 (for $L \le 8$, shift right and amount=16)
				5850 -ZR \cdot 150 (for L>8, shift left)
				5175 -ZR \cdot 150 (for L>8, shift right and amount <16)
				4875-ZR • 150 (for L>8, shift right and amount=16)
				5025 -ZR \cdot 150 (for L>8, shift right and 16 <amount<32)< td=""></amount<32)<>
	B G		0.77	4725-ZR • 150 (for L>8, shift right and amount=32)
Shift Left Double	RS	SLDA	8F	1500+N1•750
Shift Left Double Logical	RS	SLDL	8D	1200
Shift Left Single	RS	SLA	88	1050+N1•150
Shift Diskt Dauble	K5 DC	SLL	89	600
Shift Right Double	KS DC	SKDA	8E	1275
Shift Right Single	K) K)	SRUL SD A	0C 8 A	1200
Shift Dight Single Logical	KS DC	SKA SDI	0A 99	62J 600
Store	KS DV	SKL St	00 50	$375 \pm DW_{2},225 \pm V \pm 150$
Store (long)	RA RY	STD	50 60	$375 \pm DW2 \cdot 223 \pm X \cdot 150$
Store (long)	ΓΛ DV	STE	70	$375 \pm DW2 \cdot 223 \pm X \cdot 150$
Store Canacity Counts	S S	STCAP	70 B21E	2025
Store Character	ט סע	STCAP	021F 12	2023 $375 \pm X \cdot 150$
Store Characters Under Meet	DC DC	STCM	42 RF	575 ± 0.150
Store Clock	сл 2	STON	BOUL	073+0120 2175
Store Clock Comparator	3 5	STOR	B203 B207	21/3
Store Control	5 10 C	STORU	D207 R4	13/3
Store Control	ЦЭ	JICIL	DU	$\pm [floor of((N_ODS)/2)] + DSA = 225$
Store CPU ID	s	פתודא	R 202	T[1001 01((11-0K3/ 2)]•K3A•223
	3	STIDE	D404	1300

	For-	Mne-	Op-	Accura	cy
Instruction Name	mat	monic	Code	Code	Formula and Comments
Store CPU Timer	S	STPT	B209		2550
Store Halfword	RX	STH	40		375+DW2•225+X•150
Store Multiple	RS	STM	90		150+N•225
Store then AND System Mask	SI	STNSM	AC		1650 (for BC; S/370 mode or ECPS:VSE mode)
					5400 (for EC; ECPS:VSE mode)
					4800 (for EC; S/370 mode)
Store then OR System Mask	SI	STOSM	AD		1650 (for BC; S/370 mode or ECPS:VSE mode)
					5400 (for EC; ECPS:VSE mode)
					4575 (for EC; S/370 mode)
Subtract	RR	SR	1B		375
Subtract	RX	S	5B		600+DW2•225+X•150
Subtract Decimal	SS	SP	FB		1275+RC•525+ZR•1050+(2.5•DW1+DW2)•225
					(for $L1 \leq 8$ and $L2 \leq 8$)
					1950+RC•675+ZR•1050+(2•DW1+DW2+EDW2)•225
					(for $L1 \leq 8$ and $L2 > 8$)
					2625+RC•675+ZR•1050+(2•DW1+2•EDW1+DW2)•225
					(for L1>8 and L2 \leq 8)
					2700+RC•675+ZR•1050
					+(2•DW1+2•EDW1+DW2+EDW2)•225
					(for L1>8 and L2>8)
Subtract Halfword	RX	SH	4B		600+DW2•225+X•150
Subtract Logical	RR	SLR	1F		375
Subtract Logical	RX	SL	5F		$600 + DW2 \cdot 225 + X \cdot 150$
Subtract Normalized					
(extended)	RR	SXR	37	I	3834
Subtract Normalized (long)	RR	SDR	2B	I	1050
Subtract Normalized (long)	RX	SD	6B	I	1425+DW2•225+X•150
Subtract Normalized (short)	RR	SER	3B	Ι	1134
Subtract Normalized (short)	RX	SE	7B	Ι	$1434 + DW2 \cdot 225 + X \cdot 150$
Subtract Unnormalized (long)	RR	SWR	2F	I	986
Subtract Unnormalized (long)	RX	SW	6F	Ι	1518+DW2•225+X•150
Subtract Unnormalized (short)	RR	SUR	3F	Ι	1166
Subtract Unnormalized (short)	RX	SU	7F	Ι	1694+DW2•225+X•150
Supervisor Call	RR	SVC	0A	I	6075 (for BC/BC)
					6675 (for EC/BC)
					6975 (for BC-EC/EC)
Test and Set	S	TS	93		1125
Test under Mask	SI	ТМ	91		450
Test Protect	SS	TPRT	E501		8400 Condition Code 0, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
					9000 Condition Code 1, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
					9000 Condition COde 2, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
					8400 Condition Code 3, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
T 1.4.	66	тр	DC		600+1 •975
I ranslate	33 66	і К Трт	טע חח		600+NP+975+CCV+150
Translate and Test	33	1 1 1	UU		000 + 11 - 2/3 + 00 + 130

Instruct	tion N	ame	For- mat	Mne- monic	Op- Code	Accurae Code	cy Formula	and (Comments
Unpacl	k		SS	UNPK	F3		2175+(DW1 2700+(DW1 2550+(DW1	+D' +EI +D'	W2)•225 (for L1≤8, L2≤8) DW1+DW2)•225 (for L1>8, L2≤8) W2+EDW2)•225 (for L1≤8, L2>8)
Zero and Add SS		SS	ZAP		$3225 + (DW1 + EDW1 + DW2 + EDW2) \cdot 225 \text{ (for } L1 > 8, L2 > 8)$ $2475 + RC \cdot 525 + ZR \cdot 1050 + (A2 < A1 + L1) \cdot 375$ (for L1 ≤ 8 and L2 ≤ 8) $3150 + RC \cdot 675 + ZR \cdot 1050 + (A2 < A1 + L1) \cdot 375$ (for L1 ≤ 8 and L2 > 8) $3150 + RC \cdot 675 + ZR \cdot 1050 + (A2 < A1 + L1) \cdot 375$ (for L1 > 8 and L2 ≤ 8) $3225 + RC \cdot 675 + ZR \cdot 1050 + (A2 < A1 + L1) \cdot 375$				
 L aga	nd f	on Timing For	mulas				FMK	<u> </u>	EDMK adjustment Refer to Figure A 3
Lege	nu j	or Timing For	nuius				F1	_	exponent corresponding to operand 1
A	=	1 if block is address	able; else	0.			E2	_	exponent corresponding to operand 7
AB	=	1 if branch target is	on doubl	eword bou	ndary,el	se 0.	FDW1	=	1 if the first 8 bytes of operand 1 cross a doubleword
41	=	address of operand	1.				5256		boundary, else 0.
A2	=	address of operand	2.				F236	=	spanned by that portion of the longer operand that is
CCV	=	0 if condition code	(CC) is 0).					compared to the pad character, extending through the first doubleword, if any, which compare unequal.
		5 if CC is 1	,				IM	=	See Figure A-2.
		6 if CC is 2.					к	=	Number of hex digits with value greater than 9.
CLZ	=	1 if multiplicand has else 0	s 8 bytes o	or more of	leading	zeros;	L	=	length of the operand in bytes.
256	=	number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand which				aries, /hich	LC	=	0 if NDD1 \leq 15 1 if NDD1 > 15 and NDD1-NDD2 \leq 14 2 if NDD1 > 15 and NDD1-NDD2 > 14
		extending through t compares unequal. refers to operand 1	he first do If operand alignment	bubleword, ds have eq	, if any, v ual lengt	which h, this	LZ	=	number of contiguous zero result bytes starting from the beginning of the result.
СМ	=	See Figure A-2.	-				LLO	-	starting from the beginning of the result.
DW1	=	1 for L1≤8 if opera	nd 1 cros	ses a doub	leword		LI	=	length of operand 1 in bytes.
		boundary.					L2	=	length of operand 2 in bytes.
D WA	=	doubleword bounda	ry, else 0	ytes of ope	rand I c	ross a	L8	=	number of operand 1 doublewords spanned by processing.
JW2	=	 for L2≤8 if opera boundary. for L2>8 if the las doubleword bounda 	nd 2 cross st eight by rv. else 0.	ses a doub ytes of ope	leword rand 2 c	ross a	L8C	=	number of doublewords in the longer operand which are actually compared to the shorter operand. If operands have equal length, this refers to operand 1.
DWB	=	1 if first 8 bytes cros	ss a doubl	eword bou	ındary, e	lse 0.	L8F	=	number of doublewords in the longer operand which
DWQ	=	1 for (length of quo double word bounda	tient) ≤8 ary.	if the quo	tient cro	sses a	L8M	=	number of doublewords spanned by move in operand
	=	1 for (length of quo the quotient cross a	tient) >8 doublewa	if the last	eight by ry else	tes of 0	L8P	=	number of doublewords spanned by pad in operand 1.
DWR	=	1 if the remainder cr	osses a d	oubleword	bounda	rv:	М	_	0 if all test mask bits selected are 0; else 1.
- D		else 0.		T 11 0		- , ,	MA	=	0 if operands are mutually aligned on doubleword
ED EDWQ	=	1 for (length of quotient encoded)	tient) > 8	if the first	(L1-8)	bytes	MEOB	=	1 if move ends on other than a 256 byte boundary in operand 1: else 0.
EDW1	=	1 for L1>8 and L1	≤ 16 if the	first (L1- boundary	8) bytes	of	M256	=	number of 256 byte blocks, on 256 byte boundaries, spanned by move in operand 1.
FDW2	_	1 for L2N8 and L2	16 if the	first (1.2	(, ease 0.	of	N	=	number of registers in LM, STM.
<u>م ۱</u> ۱ میں		operand 2 cross a do	oubleword	l boundary	v, else 0.	51	N1	-	1 if operand 1 is negative; else 0.

N2	=	1 if operand 2 is negative; else 0.
NDD1	=	number of significant decimal digits in operand 1.
NDD2	=	number of significant decimal digits in operand 2.
NHD	=	number of significant hex digits.
NP	=	number of bytes processed.
NWB	=	number of source or destination words that cross doubleword boundaries.
ORS	-	1 if starting register number is odd, else 0.
Р	=	0 if processing halts before the last doubleword is reached; 1 if the last doubleword is processed.
PN	=	1 if post normalization is required; else 0.
PU	-	1 if Translation Lookaside Buffer purge is required; else 0.
P256	-	number of 256 byte blocks, on 256 byte boundaries, spanned by the pad in operand 1.
R	=	1 if returned to the instruction following the execute instruction after completing the subject instruction; else 0.
RC	=	1 if recomplementation is required; else 0.
RMN	=	1 if remainder is negative; else 0.
RSA	=	1 if starting register number is even and op2 is off doubleword boundary or if starting register number is odd and op2 is on a doubleword boundary; else 0.
S	=	1 if successful branch; else 0.
SA	=	sign adjustment. Refer to Figure A-3.
SI	=	time to execute the subject instruction.
VP	=	value of the operand 2.
Х	=	1 if index register number is not zero; else 0.
ZR	=	1 if result is zero; else 0.
•	=	multiply.
()	=	1 if the logical condition within () is satisfied; else 0.

Figures A-2 and A-3 are used to determine the processing time for each pattern character occurring in ED and EDMK instructions. For each such pattern character, determine the EB value and add it to the instruction time. If the pattern character is a digit selector or a significance starter, and if the source digit is the lowest-order digit in its field, add in also the sign adjustment (SA). If the instruction is EDMK and a nonzero source digit is encountered with the significance indicator off, add the EMK adjustment to the instruction time.

Mask	СМ	IM	Mask	СМ	IM
0000	0	0	1000	1	1
0001	1	0	1001	2	2
0010	1	1	1010	2	2
0011	1	0	1011	3	4
0100	1	1	1100	1	1
0101	2	2	1101	3	4
0110	1	1	1110	2	3
0111	2	2	1111	1	0

Figure A-2. Mask Character

High-Speed Buffer Storage and TLB Miss Service Times

These timings (in microseconds) are added to the instruction execution time for each miss in the high-speed buffer storage.

2.36 µs	High-speed buffer storage miss where page to
	be replaced has not been altered.
3.75 to 6.45 µs	High-speed buffer storage miss where page to
	be replaced has been altered (castout case).

For typical workloads, the average high-speed buffer storage miss time is approximately 2.85 microseconds. The frequency depends upon the addressing pattern of the program being executed.

For each miss in the Translation Lookaside Buffer (TLB), these times are added:

1.43 μs	TLB miss for S/370 mode, BC mode or EC
	DAT off.
1.43 μs	TLB miss for ECPS:VSE mode.
5.80 µs	TLB miss for S/370 mode, EC mode DAT on.
0.15 μs	Add for TLB miss if S/370 mode with DAT on
	if 64K segment.

These times assume that the necessary translation table entries are found in the high-speed buffer storage. The TLB miss frequency depends on the addressing pattern of the program being executed. Values outside the ranges given are possible.

Performance Degradation of PER Mask Setting

With bit 1 of the PSW on, EC mode on, and the PER mask bits all zero, there is no performance degradation. With PER mask bits enabled for the various categories of events shown below, and with typical instruction mixes, PER degradation (time in addition to normal time) is approximately as follows:

Pattern Character Type	Significance Indicator	Source Digit	EB Value	SA - Sign Adju Add if digit is i last before:	EMK Adjustment		
				Sign B	Sign D	+Sign	add if EDMK
	Off	0	1838	487	487	487	0
Digit	Off	1-9	2138	712	937	937	300
Selector	On	0	1838	712	937	1087	0
	On	1-9	1988	712	937	1087	0
Signifi-	Off	0	2513	712	937	782	0
cance	Off	1-9	2663	712	937	937	300
Starter	On	0	2288	637	862	1012	0
	On	1-9	2438	637	862	1012	0
		First					L
		Byte					
		of					
		Pattern					
Field		Yes	900				
Separator		No	1275				
Message	Off	Yes	900				
Char.	Off	No	1275				
	On		1050				

Figure A-3. ED and EDMK Pattern Character Timings

PER Enabled For:	Degradation
Register alteration	350 percent
Instruction fetch	415 percent
Storage alteration	320 percent
Successful branch	300 percent

These values are for cases in which no PER events are raised. The accumulative degradation with more than one event selected is not an addition of the above figures.

Effect of Hardware-Assist Features on Performance

The ECPS:VM Assist and ECPS:VS1 Assist facilities

simulate certain frequently used functions in hardware. The effect that these have on performance depends on the workload being executed and the frequency with which it requests services that are assisted.

ECPS:VM Assist and ECPS:VS1 Assist are mutually-exclusive facilities. Also, ECPS:VM Assist and ECPS:MVS are mutually-exclusive facilities.

The Engineering Scientific Assist will improve the native assembler code of equal function up to 30 percent.

4341 Processor Model Groups 9 and 10 Complex Configurator

Byte-Multiplexer STANDARD FEATURES Channel 0 Model Processor Storage (Note 1) System Diskette Drive System/370 Mode .19 1,048,576 bytes ECPS:VS1 Assist (*) K9, K10 L9, L10 ECPS:VM/370 Assist (*) 2,097,152 bytes Support Processor ECPS:MVS (*) (Note 2) 4,194,304 bytes System/370 Universal Instruction Set Note 1. The available processor Time-of-Day Clock storage depends upon the number of Translation 3278-2A Display UCWs, system microcode and other Virtual Storage Console or 3279-20 variables. Color Display (*) S/370 Mode only Console (with Basic Control (BC) Mode Byte-Oriented Operands Operator Control Panel) Channel Group 1: Note 2. Extended facility without low-address protect or common One Byte-Multiplexer Channel Two Block-Multiplexer Channels segment bit. When the machine has Channel Command Retry (Block only) been set via IML for ECPS:VSE mode, Optional 3278-2A Channel Indirect Data Addressing (*) all of the facilities described in Clear I/O Clock Comparator and CPU Timer Devices or 3279-20 IBM Processor Principles of Operation for ECPS:VSE mode are available Color Devices Commercial Instruction Set with the following restrictions: Conditional Swapping SIOF instruction is executed as a SID instruction **Control Registers** Optional 3268-2 or - There is only one virtual storage Data Streaming 3287 Printers or Decimal Instructions size (16M) therefore there is no 3287-1C or 2C Dynamic Address Translation storage size control. Color Printers (in System/370 mode) Eight Byte Parallel Data Flow 150 to 300 Nanosecond Processor Cycle OPTIONAL FEATURES Engineering Scientific Assist Error Checking and Correction Channel Group 2: System Diskette (ECC) in Processor Storage Three additional Block-Drive (not avail-Extended Control (EC) Mode Multiplexer channels or one Byteable to user as an Extended Control-Program Support 1/O device) and two Block-Multiplexer channels. (ECPS:VSE) Mode Channel-to-Channel Adapter Extended Precision Floating Point Additional Channel Control Unit External Signal Positions Floating-Point Instructions Block-Multiplexer Remote Support Facility Remote Operator Console Facility Channels 1 and 2 Halt Device 3279-2C Color Display Console High-Speed Buffer Storage I/O Error Alert and 3287-1C or 2C Color Printer Instruction Retry Optional Block-Interval Timer Note: Any combination of three 3278-2A, 3279-2C, 3268-2, and/or Limited Channel Logout Multiplexer 3287 devices is optional on the Channels 3, 4, 5 Machine Check Handling Move Inverse Instruction 4341 (in addition to the 3278-2A Program Event Recording (PER) or 3279-2C Display Console). PSW Key Handling Reloadable Control Storage Storage Protection (Store and Fetch) PREREQUISITES Store Status (System/370 Mode) or Save (ECPS:VSE Mode) 3278 Model 2A Display Console or Subchannels (128 to 1024) 3279 Model 2C Color Display Console Support Processor (both with Operator Control Panel)

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Figure A-4. 4341 Processor Model Groups 9 and 10 Complex Configurator

Appendix B. Model Group 11

Storage Characteristics

The 4341 Model Group 11 is available in three processor storage sizes:

Model K11: 2,097,152 bytes (2 megabytes) Model L11: 4,194,304 bytes (4 megabytes) Model M11: 8,388,608 bytes (8 megabytes)

Virtual storage increases the effective use of processor storage.

Note: The contents of this storage is not saved when power is removed.

Storage Access Width, Data Transfer Rate, and Internal Cycle Time

The processor has an eight-byte parallel dataflow. There is also an eight-byte path among processor, storage and channels. The processor cycle time is 120to 240-nanoseconds.

The high-speed buffer storage cycle time is 180 nanoseconds for reading or writing of data already in the buffer.

Detailed timings of the Model Group 11 instruction set are given in Instruction Timing Information in this appendix.

System Storage Requirements

A portion of processor storage is required for dynamic tables. This reduces the amount of processor storage available for user programming. Depending upon the processor configuration, the reduction of available processor storage may be from 18,432 bytes to 120,832 bytes. The reduction is the sum of the requirements of user selectable options:

- Installed storage size (processor model), plus
- Number of unit control words (UCWs) selected, plus
- Mode of operation, as shown below:

Mode of Operation	Model K11 (2 Megabytes) Processor Storage Required	Model L11 (4 Megabytes) Processor Storage Required	Model M11 (8 Megabytes) Processor Storage Required		
ECPS:VSE	49,152 Bytes	51,200 Bytes	55,296 Bytes		
System/370	10,240 Bytes	10,240 Bytes	10,240 Bytes		

Number	Processor
of UCWs	Storage Required
128	8,192 Bytes
next 32	+2,048 Bytes
next 32	+2,048 Bytes
etc.,	etc.,
up to:	up to:
1024	65,536 Bytes

Channel Configurations

The input/output channel configurations available for the Model Group 11 are:

• One byte-multiplexer channel and five block-multiplexer channels. In *byte mode*, simultaneous operation of several low-speed devices is permitted. Data transfer can be interleaved.

Note: I/O devices that are subject to data overrun (that is, the possibility of data loss), such as magnetic tape units, are not supported in burst mode on the byte-multiplexer channel.

The five block-multiplexer channels permit simultaneous operation of high-speed devices. Block-multiplexer channels are for relatively high-speed burst operations. They can multiplex complete blocks of data, and thereby permit a device to disconnect only after channel end, or after a halt instruction has been executed. This facility allows the interleaved execution of several channel programs by one channel.

The actual data rates depend upon the types of devices being serviced and the effect of concurrent processor and channel activity. For channel data rates, refer to the Data Rate chart.

• Channel 4 can be selected as a second byte-multiplexer channel. With this option, the configuration includes two byte-multiplexer channels, two 3-Mb block-multiplexer channels and two 2-Mb block-multiplexer channels.

Basic Channel Set (6 Channels)

Channel

Data Rate

Byte-Multiplexer Channel 0* Block-Multiplexer Channel 1 Block-Multiplexer Channel 2 Block-Multiplexer Channel 3 Block-Multiplexer Channel 4 (or Byte-Multiplexer Channel 5)

(See Table Below) `3 Megabytes/Second 3 Megabytes/Second 2 Megabytes/Second (See Table Below) 2 Megabytes/Second

* These devices attached natively to the Support Processor are addressed from the IPU through addresses on the byte-multiplexer channel 0.

Note: When five block-multiplexer channels are operating concurrently, the aggregate data rate is up to 12 megabytes per second.

Byte-Mode Operation (Channels 0 and 4)

These data rates are with no other channel activity. For data rates with other channel activity, see 4341 Processor Channel Characteristics.

	Byte Channel 0 Data Rate	Byte Channel 4 Data Rate				
Byte-Mode Type						
Single-Byte						
Transfer	24 kb/Second	24 kb/Second				
Four-Byte Transfer	96 kb/Second	96 kb/Second				
Burst-Mode Operation	(Channel 0)	(Channel 4)				
Device Type	Channel Data Rate					
Buffered						
Input:	2.0 mb/Second	2.0 mb/Second				
Output:	Average kb/Secon	d =				
	(1000 x device rate in kb/sec) divided by					
	(1000 + 0.8 device rate in kb/sec)					
Note: These data ro generated delays.	ates assume small interj	face-cable and control-unit				

Figure B-1. Channel Configurations and Data Rates

Instruction Timing Information

This section describes the basic Model Group 11 instruction timings and storage cycle times.

Processor Instruction Timings

Listed below are the formulas for determining the instruction execution times in nanoseconds. These formulas do not include any allowance for high-speed buffer storage misses, TLB misses, or processor interference because of I/O operations, processor storage refresh, or interval timer updates. These factors must be considered separately.

Accuracy Codes: Formulas with no identified accuracy code yield instruction execution times accurately.

Formulas with Accuracy Code I may not yield exact instruction times, but represent *average* execution times that can be expected in representative instruction sequences.

Formulas with Accuracy Code II yield instruction timing values given below, within 5 percent.

	For-	Mne-	Op-	Accura	асу
Instruction Name	mat	monic	Code	Code	Formula and Comments
Add	RR	AR	1 A		300
Add	RX	Α	5A		480+DW2•180+X•120
Add Decimal	SS	AP	FA		1020+RC•420+ZR•840+(2•DW1+DW2)•180
					(for $L1 \leq 8$ and $L2 \leq 8$)
					1560+RC•540+ZR•840+(2•DW1+DW2+EDW2)•180
					(for $L1 \leq 8$ and $L2 > 8$)
					2100+RC•540+ZR•840+(2•DW1+2•EDW1+DW2)•180
					(for L1>8 and L2 \leq 8)
					2160+RC•540+ZR•840
					+(2•DW1+2•EDW1+DW2+EDW2)•180
					(for L1>8 and L2>8)
Add Halfword	RX	AH	4A		480+DW2•180+X•120
Add Logical	RR	ALR	1E		300
Add Logical	RX	AL	5E		480+DW2•180+X•120
Add Normalized (Extended)	RR	AXR	36	I	3068
Add Normalized (Long)	RR	ADR	2A	I	840
Add Normalized (Long)	RX	AD	6A	I	1140+DW2•180+X•120
Add Normalized (Short)	RR	AER	3A	I	908
Add Normalized (Short)	RX	AE	7A	I	1147+DW2•180+X•120
Add Unnormalized (Long)	RŔ	AWR	2E	I	795
Add Unnormalized (Long)	RX	AW	6E	I	1214+DW2•180+X•120
Add Unnormalized (Short)	RR	AUR	3E	I	933
Add Unnormalized (Short)	RX	AU	7E	I	1353+DW2•180+X•120
And	RR	NR	14		300
And	RX	Ν	54		480+DW2•180+X•120
And (character)	SS	NC	D4		$1020 + DW1 \cdot 300 + DW2 \cdot 180$ (for L ≤ 8)
					720+L8•660+LZ8•60+MA•(floor of AL/8)•180 (for L>8)
					720+L•660+LZ•60 (for operands overlap)
And (Immediate)	SI	NI	94		660
Branch and Link	RR	BALR	05	I	600+S•(600-AB•240)
Branch and Link	RX	BAL	45	I	1020
Branch on Condition	RR	BCR	07	I	$120 + S \cdot (600 - AB \cdot 240)$
					900 (for serialization case)
Branch on Condition	RX	BC	47	I	120+S•(480-AB•240)
Branch on Count	RR	BCTR	06	I	420+S•(480-AB•240) (for R2≠0)
					240 (for R2=0)
Branch on Count	RX	BCT	46	I	300+S•(480-AB•240)

4341 Model Group 11 Instruction Timing List

		For-	Mne-	Op-	Accura	асу
	Instruction Name	mat	monic	Code	Code	Formula and Comments
	Propoh on Index High	DC	DVII	97	Ţ	(00 · 5 (480 AB 240)
	Branch on Index Low or Equal	KS DC	BAH	80	1	$600 + S \cdot (480 - AB \cdot 240)$
	branen on maex Low of Equal	КS	DALL	07	1	000+3•(480-AB•240)
I	Clear Page	S	CLRP	B215	II	100000
	Compare	RR	CR	19		300
	Compare	RX	С	59		480+DW2•180+X•120
	Compare (long)	RR	CDR	29		360 (for E1=E2)
						660 (for E1≠E2)
	Compare (long)	RX	CD	69		$660+DW2 \cdot 180+X \cdot 120$ (for E1=E2)
						960+DW2•180+X•120 (for E1≠E2)
	Compare (short)	RR	CER	39		360 (for E1=E2)
						900 (for E1≠E2)
	Compare (short)	RX	CE	79		660+DW2•180+X•120 (for E1=E2)
						1200+DW2•180+X•120 (for E1 <e2)< td=""></e2)<>
						1080+DW2•180+X•120 (for E1>E2)
	Compare and Swap	RS	CS	BA		780 (for OP1≠OP2)
						840 (for OP1=OP2)
	Compare Decimal	SS	СР	F9		$1020 + (DW1 + DW2) \cdot 180$ (for $L1 \le 8, L2 \le 8$)
						$1440+(DW1+EDW2+DW2) \cdot 180$ (for $L1 \le 8, L2 > 8$)
						$1560+(DW1+EDW1+DW2) \cdot 180$ (for L1>8, L2 ≤ 8)
	_					1620+(DW1+EDW1+DW2+EDW2)•180 (for L1>8, L2>8)
	Compare Double and Swap	RS	CDS	BB		1320 (for OP1≠OP2)
						1500 (for OP1=OP2)
	Compare Halfword	RX	СН	49		480+DW2•180+X•120
	Compare Logical	RR	CLR	15		300
	Compare Logical	RX	CL	55		480+DW2•180+X•120
	Compare Logical (character)	SS	CLC	D5		780 (for $L \le 8$)
						480+L8•420+MA•(floor of L/8)•180 (for L>8)
	Compare Logical (immediate)	SI	CLI	95		480
	Compare Logical Characters					
•	under Mask	RS	CLM	BD	I 	$660 + CM \cdot 120 + DW2 \cdot 180 + M \cdot 60$
I	Compare Logical Long	ĸĸ	CLCL	OF	11	540+C256•5400+L8C•540+MA•L8C•180
						(for C256>0 and F256=0)
						540+F256+(6120+(L1>0)+240)+L8F+480
						$(100 \times C256=0 \text{ and } F256>0)$
I						-1780 + F250(6120 + (L1 > L2) + 240) + C256 + 5400
						$+L_{0}F \cdot 4_{0}U + L_{0}C \cdot 54U + MA \cdot L_{0}C \cdot 18U$
	Connect Page	RC	СТР	BO		(101 C 250 > 0 and F 250 > 0)
	Convert to Binary	RY	CVB	4E		4200-C+2400
	Convert to Decimal	RX	CVD	41 4E		$1020 + NDD2 \cdot 000 + DW2 \cdot 180 + X \cdot 120$ 1020 - (1 + NHD) + K - 120 + DW2 - 180 + X - 120
		КХ		412		$1020 \cdot (1 + NHD) + K \cdot 120 + DW2 \cdot 180 + X \cdot 120$
	Deconfigure Page	S	DEP	B21B		10920
	Diagnose			83		
	Disconnect Page	S	DCTP	B21C		2700+C•11540
	Divide	RR	DR	1D		5220+N2•120+N1•300
	Divide	RX	D	5D		5760+N1•300+N2•120+DW2•180+X•120
	Divide (long)	RR	DDR	2D		8040+PN•180
	Divide (long)	RX	DD	6D		8460+PN•180+DW2•180+X•120
	Divide (short)	RR	DER	3D		4440+PN•180
	Divide (short)	RX	DE	7D		4860+(PN+DW2)•180+X•120

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	For-	Mne-	Op-	Accura	<i>cy</i>
Instruction Name	mat	monic	Code	Code	Formula and Comments
Divide Decimal	SS	DP	FD	II	$5880 + (L1 > 8) \cdot 1200 - (L1 = 16) \cdot 240$
					$-(NDD1=0) \cdot 120 + DWR \cdot 180$
					(for NDD1 <ndd2)< td=""></ndd2)<>
					9440+(NDD1-NDD2)+1080
					(for NDD1 \geq NDD2 and L1 \leq 8)
					$10640 + (1080 + LC \cdot 120) \cdot (NDD1 - NDD2)$
					$+(NDD1>15)\cdot 660-(L1=16)\cdot 240+(DwQ+DwR)\cdot 180$ (for NDD1>NDD2 and L1>8)
					Note: For cases with NDD1≥NDD2 times assume equal probability of digits 0 through 9.
Edit	SS	ED	DE	П	780+(Cond.Code≠0)•180+Sum(EB+SA)
					(where the sum is taken over all pattern
	99	EDM	DE		characters. Refer to Figure $D-3.7$
Edit and Mark	55	EDMK	DF	11	$1140 + (Cond.Code \neq 0) \cdot 180 + Sum(EB + SA + EMR)$
					(where the sum is taken over an pattern
	DD	VD	17		
Exclusive OR	KK DV	XK V	57		$480 \pm DW_{2} = 180 \pm X = 120$
Exclusive OR	KA SS		57		$480 \pm DW2 \cdot 180 \pm X \cdot 120$ 1020 ± DW1 • 300 ± DW2 • 180 (for L < 8)
Exclusive OR (character)	33	лс	DT		$960 \pm 1.8 \cdot 660 \pm 1.78 \cdot 60 \pm MA \cdot (floor of L/8) \cdot 180 (for L)$
					$720\pm1.660\pm1.7.60$ (for operands overlap)
					1080 ± 180 [for L>8 and ADDR(OP1)=ADDR(OP2)
Evolusius OR (immediate)	SI	XI	97		660
Exclusive OR (infinediate)	DY	EX	44		$900 + SI + R \cdot 2640 + X \cdot 120$ (for R1=0)
Execute	KX	LA			$1140 + SI + R \cdot 2640 + X \cdot 120 \text{ (for } R1 \neq 0)$
Halve (long)	RR	HDR	24		900+PN•180
Halve (short)	RR	HER	34		$780 + PN \cdot 180$
Insert Character	RX	IC	43		300+X•120
Insert Characters under Mask	RS	ICM	BF	I	720+DW2•180+IM•120
Insert Page Bits	RS	IPB	B4		$1800 + C \cdot 420$
Insert PSW Key	S	IPK	B20B		1080
Invalidate Page Table Entry	S	IPTE	B221		4860 (2K Page Size)
					5100 (4K Page Size)
Insert Storage Key	RR	ISK	09		1140 (for BC, S/370 mode)
					1260 (for EC, S/370 mode)
					1380 (for BC, ECPS: VSE mode) 1980 (for EC, ECPS: VSE mode)
Load	RR	LR	18		240
Load	RX	L	58		$300 + DW2 \cdot 180 + X \cdot 120$
Load (long)	RR	- LDR	28		240
Load (long)	RX	LD	68		300+DW2•180+X•120
Load (short)	RR	LER	38		240
Load (short)	RX	LE	78		300+DW2•180+X•120
Load Address	RX	LA	41		240+X•120
Load and Test	RR	LTR	12		300
Load and Test (long)	RR	LTDR	22		420
	DD	LTER	32		420
Load and Test (short)	ĸĸ				
Load and Test (short) Load Complement	RR	LCR	13		300
Load and Test (short) Load Complement Load Complement (long)	RR RR	LCR LCDR	13 23		300 540

		For-	Mne-	Op-	Accura	cy
	Instruction Name	mat	monic	Code	Code	Formula and Comments
r	Load Control	RS	LCTL	B7		10200+PU•4320 (for S/370 mode, load all Regs)
						9600 (for ECPS:VSE, load all Regs)
	·					4020+PU•4320 (for S/370 mode, load Regs 0,1)
						3060 (for ECPS:VSE, load Regs 0,1)
	Load Frame Index	RS	LFI	B 8		1740
	Load Halfword	RX	LH	48		300+DW2•180+X•120
	Load Multiple	RS	LM	98		120+N•180
	Load Negative	RR	LNR	11		420-N2•120
	Load Negative (long)	RR	LNDR	21		420
	Load Negative (short)	RR	LNER	31		540
	Load Positive	RR	LPR	10		300+N2•180
	Load Positive (long)	RR	LPDR	20		540
	Load Positive (short)	RR	LPER	30		540
	Load PSW	S	LPSW	82		4200 (for BC/BC)
						4440 (for EC/BC)
						4920 (for BC-EC/EC)
	Load Real Address	RX	LRA	B 1		3840 (for 1M segment)
						4080 (for 64K segment)
	Load Rounded (extended to					
	long)	RR	LRDR	25		600
	Load Rounded (long to short)	RR	LRER	35		720
	Make Addressable	S	MAD	B21D		2460-A•120
	Make Unaddressable	S	MUN	B21E		2340+A•360
	Monitor Call	SI	MC	AF		11100
						840 (for NO-OP case)
	Move (character)	SS	MVC	D2		$660 + DW1 \cdot 180 + DW2 \cdot 120$ (for $L \le 8$)
						540+L8•360+MA•(floor of L/8)•180 (for L>8)
						720+L•360 (for L ≤ 8 and operands overlap)
						780+L•360 (for L>8 and operands overlap, $A1 \neq A2+1$)
1						$1380 + L8 \cdot 180$ (for L>8 and A1=A2+1)
	Move (immediate)	SI	MVI	92		300
	Move Inverse	SS	MVCIN	E8		600+L•600
	Move Long	RR	MVCL	0E	II	300+M256•5491+P256•6151-MEOB•4591+L8M•360 +L8P•180
						(for M256>0 and P256>0)
						300+M256•5491+MEOB•120+L8M•360
						(for M256>0 and P256=0)
						$-480 + P256 \cdot 6151 + L8P \cdot 180$
						(for M256=0 and P256>0)
	Move Numerics	SS	MVN	D1		$1440 + DW^{2} \cdot 180 + DW^{1} \cdot 300 (for 1 < 8)$
			• • •			$(18+1) \cdot 900 + MA \cdot (floor of 1/8) + 180 (for 1 > 8)$
						(50 + 1) = 200 + 100 +
						$660 \pm L = 200 \pm D$ w 1 ± 120 (for L > 8 and operands overlap)
						oou+L. 900 (for L>8 and operands overlap)

 $\tilde{r}_{\tilde{t}}$

	For-	Mne-	Op-	Accuracy
Instruction Name	mat	monic	Code	Code Formula and Comments
Move with Offset	SS	MVO	F1	$1920 + (A2 < A1 + L1) \cdot 300 + (2 \cdot DW1 + DW2) \cdot 180$
				(for $L1 \le 8, L2 \le 8$)
				$2280 + (A2 < A1 + L1) \cdot 300 + (2 \cdot DW1 + DW2 + EDW2) \cdot 180$
				(for $L1 \le 8, L2 > 8$)
				2880+(A2 <a1+l1)•300+(2•edw1+2•dw1+dw2)•180< td=""></a1+l1)•300+(2•edw1+2•dw1+dw2)•180<>
				(for $L1 > 8, L2 \le 8$)
				$3120 + (A2 < A1 + L1) \cdot 300$
				+(2•EDW1+2•DW1+DW2+EDW2)•180 (for L1>8,L2>8)
				$3120 + \min(L1, L2) \cdot 840 + (L1 > L2) \cdot 420 +$
				$[(L1>L2+8)+(L1>8)+(L2>8)] \cdot 180$ (for overlapping case)
Move Zones	SS	MVZ	D3	$1440 + DW1 \cdot 300 + DW2 \cdot 180$ (for L ≤ 8)
				$(L8+1) \cdot 900 + MA \cdot (floor of L/8) \cdot 180 (for L>8)$
				$660+L \cdot 900+DW1 \cdot 120$ (L ≤ 8 and operands overlap)
				$660+L \cdot 900$ (L>8 and operands overlap)
Multiply	RR	MR	1 C	2880+N2•360
				840 (for VP=0)
Multiply	RX	М	5C	3120+N2•360+DW2•180+X•120
Multiply (extended)	RR	MXR	26	12420+PN•60
Multiply (long to extended)	RR	MXDR	27	4680+PN•60
Multiply (long to extended)	RX	MXD	67	5100+PN•60+DW2•180+X•120
Aultiply (short to long)	RR	MER	3C	3060+PN•180
Aultiply (short to long)	RX	ME	7C	3480+PN•180+DW2•180+X•120
Multiply (long)	RR	MDR	2C	3900+PN•180
Aultiply (long)	RX	MD	6C	4320+PN•180+DW2•180+X•120
Aultiply Decimal	SS	MP	FC	$3900-(NDD2=0) \cdot 120 +$
				(3•DW1+DW2)•180+(1.33+FDW1-DW1+2•EDW1)•180 (for NDD1=0)
				4560+(FDW1+DW2+2•EDW1+2•DW1)•180
				(for NDD1>15 and NDD2=0) $($
				4080-CLZ•420+(3•DW1+DW2)•180
				$+(L1>8) \cdot (3.67+FDW1-DW1+2 \cdot EDW1) \cdot 180$
				(for NDD1 \leq 15 and NDD2=0)
				$6540 + NDD2 \cdot 120 - CLZ \cdot 420 + (3 \cdot DW1 + DW2) \cdot 180$
				$+(L1>8) \cdot (3+FDW1-DW1+2 \cdot EDW1) \cdot 180$
				(for NDD1<15 and NDD2>0)
				9300-CLZ•420+(3•DW1+DW2)•180
				+(L1>8)•(3+FDW1-DW1+2•EDW1)•180
				(for NDD1=15 and NDD2<15)
				14880-CLZ•420+(3•DW1+DW2)•180
				$+(L1>8) \cdot (3+FDW1-DW1-2 \cdot EDW1) \cdot 180$
				(for NDD1=NDD2=15)
				11420+240•NDD2+(FDW1+DW2+2•EDW1+2•DW1)•186
				(for NDD1>15 and NDD2>0)
				4920+VP•480+(FDW1+DW2+2•EDW1+2•DW1)•180
				(for NDD1=29 and NDD2=1) $($
fultiply Halfword	RX	МН	4C	2460+(N2+X)•120+DW2•180
R	RR	OR	16	300
)R	RX	0	56	480+DW2•180+X•120
)R (character)	SS	OC	D6	$1020 + DW1 \cdot 300 + DW1 \cdot 180$ (for L ≤ 8)
				720+L8•660+LZ8•60+MA•(floor of L/8)•180 (for L>8)
				$720+L \cdot 660+LZ \cdot 60$ (for operands overlap)

	For-	Mne-	Op-	Accuracy
Instruction Name	mat	monic	Code	Code Formula and Comments
OR (immediate)	SI	OI	96	660
Pack	SS	PACK	F2	$1380 + (DW1 + DW2) \cdot 180$ (for $L1 \le 8$, $L2 \le 8$)
				$1680 + (DW1 + EDW1 + DW2) \cdot 180 \text{ (for } L1 > 8, L2 \le 8)$
				$1920 + (DW1 + DW2 + EDW2) \cdot 180$ (for $L1 \le 8, L2 > 8$)
				$2400+(DW1+EDW1+DW2+EDW2) \cdot 180$ (for L1>8, L2>8)
Purge TLB	S	PTLB	B20D	5400
Reset Reference Bit	S	RRB	B213	$2580 + C \cdot 120$ (for ECPS:VSE mode)
				1500 (for S/370 mode)
Retrieve Status and Page	SS	RSP	D8	106740
Set Clock	S	SCK	B204	6120 (for Enable ON)
				1560 (for Enable OFF)
Set Clock Comparator	S	SCKC	B206	4260
Set CPU Timer	S	SPT	B208	8940
Set Page Bits	RS	SPB	B5	2940
Set Program Mask	RR	SPM	04	600
Set PSW Key from Address	S	SPKA	B20A	1860
Set Storage Key	RR	SSK	08	6120 (for S/370 mode)
				3780 (for ECPS:VSE mode)
Set System Mask	S	SSM	80	1380 (for BC; S/370 mode or ECPS:VSE mode)
				4380 (for EC; ECPS:VSE mode)
				3900 (for EC; S/370 mode)
Shift and Round Decimal	SS	SRP	F0	2880-ZR•180 (for L \leq 8, shift left and amount $<$ 16)
				2640-ZR•120 (for L ≤ 8 , shift right and amount < 16)
				1920 (for $L \leq 8$, shift left and amount ≥ 16)
				1740 (for $L \le 8$, shift right and amount>16)
				1920 (for $L \le 8$, shift right and amount=16)
				4680-ZR • 120 (for L>8, shift left)
				4140-ZR•120 (for L>8, shift right and amount<16)
				$3900-ZR \cdot 120$ (for L>8, shift right and amount=16)
				4020-ZR•120 (for L>8, shift right and 16 <amount<32)< td=""></amount<32)<>
				3780-ZR • 120 (for L>8, shift right and amount=32)
Shift Left Double	RS	SLDA	8F	$1200 + N1 \cdot 600$
Shift Left Double Logical	RS	SLDL	8D	960
Shift Left Single	RS	SLA	8B	840+N1•120
Shift Left Single Logical	RS	SLL	89	480
Shift Right Double	RS	SRDA	8E	1020
Shift Right Double Logical	RS	SRDL	8C	960
Shift Right Single	RS	SRA	8A	660
Shift Right Single Logical	RS	SRL	88	480
Store	RX	ST	50	300+DW2•180+X•120
Store (long)	RX	STD	60	300+DW2•180+X•120
Store (short)	RX	STE	70	300+DW2•180+X•120
Store Capacity Counts	S	STCAP	B21F	1620
Store Character	RX	STC	42	300+X•120
Store Characters Under Mask	RS	STCM	BE	540+CM•120
Store Clock	S	STCK	B205	1740
Store Clock Comparator	S	STCKC	B207	1260
Store Control	RS	STCTL	B6	720+[ORS+ceiling of(N-ORS)/2]•360
				+[floor of($(N-ORS)/2$)]•RSA•180
Store CPU ID	S	STIDP	B202	1200

	For-	Mne-	Op-	Accura	псу
Instruction Name	mat	monic	Code	Code	Formula and Comments
Store CPU Timer	S	STPT	B209		2040
Store Halfword	RX	STH	40		300+DW2•180+X•120
Store Multiple	RS	STM	90		$120 + N \cdot 180$
Store then AND System Mask	SI	STNSM	AC		1320 (for BC; S/370 mode or ECPS:VSE mode)
					4320 (for EC; ECPS:VSE mode)
					3840 (for EC; S/370 mode)
Store then OR System Mask	SI	STOSM	AD		1320 (for BC; S/370 mode or ECPS:VSE mode)
					4320 (for EC; ECPS:VSE mode)
					3660 (for EC; S/370 mode)
Subtract	RR	SR	1B		300
Subtract	RX	S	5B		480+DW2•180+X•120
Subtract Decimal	SS	SP	FB		$1020 + RC \cdot 420 + ZR \cdot 840 + (2.5 \cdot DW1 + DW2) \cdot 180$ (for L1 < 8 and L2 < 8)
					$1560+RC \cdot 540+ZR \cdot 840+(2 \cdot DW1+DW2+FDW2) \cdot 180$
					(for L1<8 and L2>8)
					$2100 + RC \cdot 540 + 7R \cdot 840 + (2 \cdot DW1 + 2 \cdot EDW1 + DW2) = 18$
					(for L1>8 and L2<8)
					$2160+BC \cdot 540+7B \cdot 840$
					$+(2 \cdot DW1 + 2 \cdot EDW1 + DW2 + EDW2) \cdot 180$
					(for L1>8 and L2>8)
Subtract Halfword	RX	SH	4B		$480 + DW2 \cdot 180 + X \cdot 120$
Subtract Logical	RR	SLR	1F		300
Subtract Logical	RX	SL	5F		480+DW2•180+X•120
Subtract Normalized					
(extended)	RR	SXR	37	I	3068
Subtract Normalized (long)	RR	SDR	2B	I	840
Subtract Normalized (long)	RX	SD	6B	I	1140+DW2•180+X•120
Subtract Normalized (short)	RR	SER	3B	I	908
Subtract Normalized (short)	RX	SE	7B	I	1148+DW2•180+X•120
Subtract Unnormalized (long)	RR	SWR	2F	I	789
Subtract Unnormalized (long)	RX	SW	6F	Ι	1215+DW2•180+X•120
Subtract Unnormalized (short)	RR	SUR	3F	I	933
Subtract Unnormalized (short)	RX	SU	7F	I	1356+DW2•180+X•120
Supervisor Call	RR	SVC	0A	I	4860 (for BC/BC)
					5340 (for EC/BC)
					5580 (for BC-EC/EC)
Test and Set	S	TS	93		900
Test under Mask	SI	ТМ	91		360
Test Protect	SS	TPRT	E501		6720 Condition Code 0, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
					7200 Condition Code 1, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
					7200 Condition COde 2, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
					6720 Condition Code 3, 1Meg Seg.)
					(add 300 ns for 64K Seg.)
Translate	55	TR	DC		480+1.780
Franslate and Test	SS	TRT	DD		$480 \pm NP \cdot 780 \pm CCV \cdot 120$
	55		~ ~		100 1 11 - 100 1 - 00 + - 120

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Instruction	n Na	me	For- Mne- Op- Accuracy e mat monic Code Code Formula and Comments			Somments			
Unpack			SS	UNPK	F3		1740+(DW1- 2160+(DW1- 2040+(DW1-	+DV +ED +DV	$V2) \cdot 180 \text{ (for } L1 \le 8, L2 \le 8)$ $V1+DW2) \cdot 180 \text{ (for } L1 > 8, L2 \le 8)$ $V2+EDW2) \cdot 180 \text{ (for } L1 \le 8, L2 > 8)$
Zero and Add			SS	ZAP	F8		2580+(DW1- 1980+RC-42 (for L1< 2520+RC-54 (for L1< 2520+RC-54 (for L1>	+ED 20+2 8 and 40+2 8 and 40+2 8 and 8 and	$W1+DW2+EDW2) \cdot 180 \text{ (for } L1>8, L2>8)$ $ZR \cdot 840+(A2 d L2 \leq 8) ZR \cdot 840+(A2 d L2>8) ZR \cdot 840+(A2 d L2 \leq 8)$
							2580+RC•54 (for L1>	10+2 8 and	$2R \cdot 840 + (A2 < A1 + L1) \cdot 300$ d L2>8)
		Timina Fam					F1		exponent corresponding to operand 1
Legen	u je _	1 if block is addressa	iuius ble: else	• 0			E2	_	exponent corresponding to operand 2.
	-	1 if bronch torget is c	n doubl	ourd hou	undary a	lee ()	FDW1	_	1 if the first 8 bytes of operand 1 cross a doubleword
AB	=	I il branch target is c	n uouo	leword bot	inuary,e	ise 0.	10.01	-	boundary, else 0.
A1 A2	-	address of operand 1. address of operand 2.					F256	=	number of 256-byte blocks, on 256-byte boundaries,
A2	=								spanned by that portion of the longer operand that is
CCV	=	1 II block is connecte	From $C(C)$ is 0				first doubleword, if any, which compare unequal.		
	Ξ	5 if CC is 1	(0,0) is 0				IM	=	See Figure B-2.
		6 if CC is 2.					ĸ	=	Number of hex digits with value greater than 9.
CLZ	=	1 if multiplicand has	8 bytes	or more of	f leading	zeros;	T.	_	length of the operand in bytes
	else 0.			LC	=	0 if NDD1 < 15			
C256 = number of 256-byte blocks, on 256-byte spanned by that portion of the longer oper is actually compared to the shorter operan extending through the first doubleword, i compares unequal. If operands have equa		te bound operand rand,	laries, which			1 if NDD1 > 15 and NDD1-NDD2 \leq 14 2 if NDD1 > 15 and NDD1-NDD2 > 14			
		e first d f operar	oubleword ids have ec	l, if any, qual leng	which ;th, this	LZ	=	number of contiguous zero result bytes starting from the beginning of the result.	
CM		See Figure B-2	ngnmen	IL.			LZ8		starting from the beginning of the result.
	_	1 for $I \le 8$ if operan	d 1 cros	sses a dout	bleword		L1	=	length of operand 1 in bytes.
DWI	-	boundary.		55C5 a dout	ne word		L2	-	length of operand 2 in bytes.
	=	1 for L1>8 if the las doubleword boundar	t eight t y, else (oytes of op).	erand 1	cross a	L8	=	number of operand 1 doublewords spanned by processing.
DW2	=	1 for $L2 \le 8$ if operandoper	id 2 cros t eight b	sses a dout	erand 2	cross a	L8C	-	number of doublewords in the longer operand which are actually compared to the shorter operand. If operands have equal length, this refers to operand 1.
DWB	=	1 if first 8 bytes cros	y, eise (s a doub). Deword bo	oundarv	else 0	L8F	-	number of doublewords in the longer operand which
DWO	=	1 for (length of quot	ient) < 9	B if the and	otient cr	osses a	1011		are actually compared to the pad character.
	=	double word bounda 1 for (length of quot	ry. ient) >8	B if the last	t eight b	ytes of	L8M 1.8P	-	number of doublewords spanned by move in operand 1. number of doublewords spanned by pad in operand 1
		the quotient cross a d	loublew	ord bound	iary, else	÷ 0.	M	_	Ω if all test mask hits calcoted are Ω , also 1
DWR	=	1 if the remainder croelse 0.	osses a (doublewor	d bound	ary;	MA	=	0 if operands are mutually aligned on doubleword
EB	=	ED, EDMK base val	ue. Refe	er to Figur	e B-3.				boundaries, else 1.
EDWQ	=	1 for (length of quot of the quotient cross	ient) >8 a doub	8 if the firs leword bou	st (L1-8) andary, e) bytes else 0.	MEOB	-	1 if move ends on other than a 256 byte boundary in operand 1; else 0.
EDW1	=	1 for L1>8 and L1 \leq operand 1 cross a do	16 if th ublewo	e first (L1 rd boundai	-8) byte ry, else 0	s of).	M256	-	number of 256 byte blocks, on 256 byte boundaries, spanned by move in operand 1.
EDW2	=	1 for L2>8 and L2 \leq	16 if th	e first (L2	-8) byte	s of	Ν	=	number of registers in LM, STM.
						۱	N11		1 if an around 1 is represented along O
		operand 2 cross a do	ublewo	rd boundar	ry, eise u).	NI	=	The operand This negative; else 0.

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NDD1	=	number of significant decimal digits in operand 1.
NDD2	=	number of significant decimal digits in operand 2.
NHD	=	number of significant hex digits.
NP	=	number of bytes processed.
NWB	=	number of source or destination words that cross doubleword boundaries.
ORS	=	1 if starting register number is odd, else 0.
Р	=	0 if processing halts before the last doubleword is reached; 1 if the last doubleword is processed.
PN	=	1 if post normalization is required; else 0.
PU	=	1 if Translation Lookaside Buffer purge is required; else 0.
P256		number of 256 byte blocks, on 256 byte boundaries, spanned by the pad in operand 1.
R	=	1 if returned to the instruction following the execute instruction after completing the subject instruction; else 0.
RC	=	1 if recomplementation is required; else 0.
RMN	=	1 if remainder is negative; else 0.
RSA	=	1 if starting register number is even and op2 is off doubleword boundary or if starting register number is odd and op2 is on a doubleword boundary; else 0.
S	=	1 if successful branch; else 0.
SA	=	sign adjustment. Refer to Figure B-3.
SI	=	time to execute the subject instruction.
VP	=	value of the operand 2.
х	=	1 if index register number is not zero; else 0.
ZR	=	1 if result is zero; else 0.
•	=	multiply.
()	=	1 if the logical condition within () is satisfied; else 0.

Figures B-2 and B-3 are used to determine the processing time for each pattern character occurring in ED and EDMK instructions. For each such pattern character, determine the EB value and add it to the instruction time. If the pattern character is a digit selector or a significance starter, and if the source digit is the lowest-order digit in its field, add in also the sign adjustment (SA). If the instruction is EDMK and a nonzero source digit is encountered with the significance indicator off, add the EMK adjustment to the instruction time.

Mask	СМ	IM	Mask	СМ	IM
0000	0	0	1000	1	1
0001	1	0	1001	2	2
0010	1	1	1010	2	2
0011	1	0	1011	3	4
0100	1	1	1100	1	1
0101	2	2	1101	3	4
0110	1	1	1110	2	3
0111	2	2	1111	1	0

Figure B-2. Mask Character

High-Speed Buffer Storage and TLB Miss Service Times

These timings (in microseconds) are added to the instruction execution time for each miss in the high-speed buffer storage.

l	1.89 μ s Read or Write	High-speed buffer storage miss where page to be replaced has not been altered.
	3.00 to 5.16 μ s	High-speed buffer storage miss where page to be replaced has been altered (castout case)

For typical workloads, the average high-speed buffer storage miss time is approximately 2.28 microseconds. The frequency depends upon the addressing pattern of the program being executed.

For each miss in the Translation Lookaside Buffer (TLB), these times are added:

	1.144 μs	TLB miss for S/370 mode, BC mode or EC DAT off.
I	1.144 µs	TLB miss for ECPS:VSE mode.
I	4.64 µs	TLB miss for S/370 mode, EC mode DAT on

These times assume that the necessary translation table entries are found in the high-speed buffer storage. The TLB miss frequency depends on the addressing pattern of the program being executed. Values outside the ranges given are possible.

Performance Degradation of PER Mask Setting

With bit 1 of the PSW on, EC mode turned on, and the PER mask bits all zero, there is no performance degradation. With PER masks enabled for the various categories of events shown below, and with typical instruction mixes, PER degradation (time in addition to normal time) is approximately as follows:

PER Enabled For:	Degradation
Register alteration	430 percent
Instruction fetch	510 percent
Storage alteration	400 percent
Successful branch	360 percent

These are for cases in which no PER events are raised. The accumulative degradation with more than one event selected is not an addition of the above figures.

Pattern Character Type	Significance Indicator	Source Digit	EB Value	SA - Sign Adjı Add if digit is last before:	ustment the	EMK Adjustment add if	
				Sign B	Sign D	+Sign	EDMK
Digit	Off Off On	0 1-9 0	1470 1710 1470	390 570 570	390 750 750	390 750 870	0 240 0
30100101	On	1-9	1590	570	750	870	0
Signifi- cance Starter	Off Off On On	0 1-9 0 1-9	2010 2130 1830 1950	570 570 510 510	750 750 690 690	626 750 810 810	0 240 0 0
		First Byte of Pattern					
Field Separator		Yes No	720 1020				
Message Char.	Off Off On	Yes No 	720 1020 840				

Figure B-3. ED and EDMK Pattern Character Timings

Effect of Hardware-Assist Features on Performance

The ECPS:VM Assist and ECPS:VS1 Assist facilities simulate certain frequently used functions in hardware. The effect that these have on performance depends on the workload being executed and the frequency with which it requests services that are assisted. ECPS:VM Assist and ECPS:VS1 Assist are mutually-exclusive facilities. Also, ECPS:VM Assist and ECPS:MVS are mutually-exclusive facilities.

The Engineering Scientific Assist will improve the native assembler code of equal function up to 30 percent.





Appendix C. Model Group 12

Storage Characteristics

The 4341 Model Group 12 is available in five processor storage sizes:

Model K12:	2,097,152 bytes (2 megabytes)
Model L12:	4,194,304 bytes (4 megabytes)
Model M12:	8,388,608 bytes (8 megabytes)
Model N12:	12,582,912 bytes (12 megabytes)
Model P12:	16,777,216 bytes (16 megabytes)

Virtual storage increases the effective use of processor storage.

Note: The contents of this storage is not saved when power is removed.

Storage Access Width, Data Transfer Rate, and Internal Cycle Time

The processor has an eight-byte parallel dataflow. There is also an eight-byte path among processor, storage and channels. The processor cycle time is 115to 230-nanoseconds. The high-speed buffer storage cycle time is 115 nanoseconds for reading and 173 nanoseconds for writing of data already in the buffer.

Detailed timings of the Model Group 12 instruction set are given in Instruction Timing Information in this appendix.

System Storage Requirements

A portion of processor storage is required for dynamic tables. This reduces the amount of processor storage available for user programming. Depending upon the processor configuration, the reduction of available processor storage may be from 18,432 bytes to 129,024 bytes. The reduction is the sum of the requirements of user selectable options:

- Installed storage size (processor model), plus
- Number of unit control words (UCWs) selected, plus
- Mode of operation, as shown in the following:

Mode of Operation	Model K12 (2 Megabytes) Processor Storage Required	Model L12 (4 Megabytes) Processor Storage Required	Model M12 (8 Megabytes) Processor Storage Required	Model N12 (12 Megabytes) Processor Storage Required	Model P12 (16 Megabytes) Processor Storage Required
ECPS:VSE System/370	49,152 Bytes 10,240 Bytes	51,200 Bytes 10,240 Bytes	55,296 Bytes 10,240 Bytes	59,392 Bytes 10,240 Bytes	63,488 Bytes 10,240 Bytes
Number of UCWs	Processor Storage Required				
128 next 32 next 32 etc.,	8,192 Bytes +2,048 Bytes +2,048 Bytes etc.,				

up to:

65,536 Bytes

up to:

Channel Configurations

The input/output channel configurations available for the Model Group 12 are:

• One byte-multiplexer channel and five block-multiplexer channels. In *byte mode*, simultaneous operation of several low-speed devices is permitted. Data transfer can be interleaved.

Note: I/O devices that are subject to data overrun (that is, the possibility of data loss), such as magnetic tape units, are not supported in burst mode on the byte-multiplexer channel.

The five block-multiplexer channels permit simultaneous operation of high-speed devices and are either 2-MB or 3-MB maximum data rates. Block-multiplexer channels are for relatively high-speed burst operations. They can multiplex complete blocks of data, and thereby permit a device to disconnect only after channel end, or after a halt instruction has been executed. This facility allows the interleaved execution of several channel programs by one channel.

The actual data rates depend upon the types of devices being serviced and the effect of concurrent processor and channel activity. For channel data rates, refer to the Data Rate chart.

Channel 4 or 5 can be selected as a second byte-multiplexer channel. If channel 4 is selected as a second byte-multiplexer channel, the congifuration is two byte-multiplexer channels, two 3-MB block-multiplexer channels and two 2-MB block-multiplexer channels. If channel 5 is selected as a second byte-multiplexer channel, the configuration is two byte-multiplexer channels, three 3-MB block-multiplexer channels and one 2-MB block-multiplexer channel.

Programming Support

Programming support for the 4341 in ECPS:VSE mode is provided by DOS/VSE. In System/370 mode, programming support is provided by DOS/VSE, OS/VS1, MVS/SP, and VM/370. Note that in System/370 mode, DOS/VSE will not operate on 4341 Models N12 or P12.

Control Storage Expansion

This optional feature adds 16K additional bytes of control store, and provides the capability to run ECPS:VM and ECPS:MVS concurrently.

Basic Channel Set (6 Channels)

Channel

Data Rate

Byte-Multiplexer Channel 0* Block-Multiplexer Channel 1 Block-Multiplexer Channel 2 Block-Multiplexer Channel 3 Block-Multiplexer Channel 4 (or Byte-Multiplexer Channel 5) (or Byte-Multiplexer Channel 5) (See Table Below)
3 Megabytes/Second
3 Megabytes/Second
2 Megabytes/Second
3 Megabytes/Second
(See Table Below)
2 Megabytes/Second)
(See Table Below)

* These devices attached natively to the Support Processor are addressed from the IPU through addresses on the byte-multiplexer channel 0.

Note: When five block-multiplexer channels are operating concurrently, the aggregate data rate is up to 13 megabytes per second.

Byte-Mode Operation (Channels 0 and 4 or 5)

These data rates are with no other channel activity. For data rates with other channel activity, see 4341 Processor Channel Characteristics.

	Byte Channel 0 Data Rate	Byte Channel 4 or 5 Data Rate						
Byte-Mode Type								
Single-Byte Transfer Four-Byte	24 kb/Second	24 kb/Second						
Transfer	96 kb/Second	96 kb/Second						
Burst-Mode Operation	(Channel 0)	(Channel 4 or 5)						
Device Type	Channel Data Rate							
Buffered Input:	2.0 mb/Second	2.0 mb/Second						
Output:	Average kb/Secon	d =						
	(1000 x device rate in kb/sec) divided by							
	(1000 + 0.8 device rate in kb/sec)							
Note: These data ro	ites assume small interi	face-cable and control-unit						

Figure C-1. Channel Configurations and Data Rates

generated delays.

Instruction Timing Information

This section describes the basic Model Group 12 instruction timings and storage cycle times.

Processor Instruction Timings

Listed below are the formulas for determining the instruction execution times in nanoseconds. These formulas do not include any allowance for high-speed buffer storage misses, TLB misses, or processor interference because of I/O operations, processor storage refresh, or interval timer updates. These factors must be considered separately.

Accuracy Codes: Formulas with no identified accuracy code yield instruction execution times accurately.

Formulas with Accuracy Code I may not yield exact instruction times, but represent *average* execution times that can be expected in representative instruction sequences.

Formulas with Accuracy Code II yield instruction timing values given below, within 5 percent.

	For-	Mne-	Op-	Accura	cy
Instruction Name	mat	monic	Code	Code	Formula and Comments
Add	RR	AR	1 A		288
Add	RX	Α	5A		$403 + (DW2 + X) \cdot 115$
Add Decimal	SS	AP	FA		863+RC•403+ZR•805+(2.5•DW1+DW2)•115
					(for $L1 \leq 8$ and $L2 \leq 8$)
					1323+RC•518+ZR•805+(2.5•DW1+DW2)•115
					(for $L1 \leq 8$ and $L2 > 8$)
					1840+RC•518+ZR•805
					+(2.5•DW1+2.5•EDW1+DW2)•115
					(for L1>8 and L2 \leq 8)
					1840+RC•518+ZR•805
					+(2.5•DW1+2.5•EDW1+DW2+EDW2)•115
					(for L1>8 and L2>8)
Add Halfword	RX	AH	4A		460+(DW2+X)•115
Add Logical	RR	ALR	1E		230
Add Logical	RX	AL	5E		345+(DW2+X)•115
Add Normalized (Ext.)	RR	AXR	36	I	2408
Add Normalized (Long)	RR	ADR	2A	I	764
Add Normalized (Long)	RX	AD	6A	I	994+(DW2+X)•115
Add Normalized (Short)	RR	AER	3A	I	892
Add Normalized (Short)	RX	AE	7 A	I	$1028 + (DW2 + X) \cdot 115$
Add Unnormalized (Long)	RR	AWR	2E	I	718
Add Unnormalized (Long)	RX	AW	6E	I	$1063 + (DW2 + X) \cdot 115$
Add Unnormalized (Short)	RR	AUR	3E	I	841
Add Unnormalized (Short)	RX	AU	7E	I	1185+(DW2+X)•115
And	RR	NR	14		230
And	RX	N	54		345+(DW2+X)•115
And (Character)	SS	NC	D4		$805+DW1 \cdot 230+DW2 \cdot 115$ (for L ≤ 8)
					690+L8•518+MA•(floor of L/8)•115+LZ8•58
					(for L>8)
					690+L•518+LZ•58 (for operands overlap)
And (Immediate)	SI	NI	94		518
Branch And Link	RR	BALR	05	I	690+S•(230)
Branch and Link	RX	BAL	45	I	690
Branch on Condition	RR	BCR	07	Ι	115+S•(230)
					633 (for serialization case)
Branch on Condition	RX	BC	47	Ι	115+S•(230)
Branch on Count	RR	BCTR	06	I	288+S•(230) (for R2≠0)
					230 (for R2=0)

4341 Model Group 12 Instruction Timing List

		For-	Mne-	Op-	Accura	асу
	Instruction Name	mat	monic	Code	Code	Formula and Comments
	Branch on Count	DV	DCT	16	т	288 (220)
	Branch on Index High	RA		40	1	288+S•(230)
	Branch on Index High	RS	влн	80	1	575+8•(230)
	Branch on Index Low of Equal	KS	BALE	87	I	575+8•(230)
	Clear Page	S	CLRP	B215	П	89911
	Compare	RR	CR	19		230
	Compare	RX	С	59		$345 + (DW2 + X) \cdot 115$
	Compare (long)	RR	CDR	29		287 (for E1=E2)
						602 (for E1•E2)
	Compare (long)	RX	CD	69		$575+(DW2+X) \cdot 115$ (for E1=E2)
						863+(DW2+X)•115 (for E1•E2)
	Compare (short)	RR	CER	39		287 (for E1=E2)
						842 (for E1•E2)
	Compare (short)	RX	CE	79		$517 + (DW2 + X) \cdot 115$ (for E1=E2)
						$1035 + (DW2 + X) \cdot 115$ (for E1 <e2)< td=""></e2)<>
						920+(DW2+X)•115 (for E1>E2)
	Compare and Swap	RS	CS	BA		690 (for OP1≠OP2)
						748 (for $OP1=OP2$)
	Compare Decimal	SS	СР	F9		$863 + (DW1 + DW2) \cdot 115$
						(for $L1 \le 8, L2 \le 8$)
I						$1208 + (DW1 + EDW2 + DW2) \cdot 115$
						(for L1 < 8, L2 > 8)
						$1323 + (EDW1 + DW1 + DW2) \cdot 115$
I						(for L1>8, L2<8)
						$1323 + (EDW1 + EDW2 + DW1 + DW2) \cdot 115$
						(for L1>8, L2>8)
I	Compare Double and Swap	RS	CDS	BB		1208 (for OP1≠OP2)
I						1380 (for OP1=OP2)
	Compare Halfword	RX	СН	49		$402 + (DW2 + X) \cdot 115$
	Compare Logical	RR	CLR	15		230
	Compare Logical	RX	CL	55		345+(DW2+X)•115
	Compare Logical (character)	SS	CLC	D5		575+DW1•115
						+DW2•115 (for $L \le 8$)
						$402+L8\cdot403+MA\cdot(floor of L/8)\cdot115$ (for L>8)
	Compare Logical (immediate)	SI	CLI	95		345
	Compare Logical Characters					
	under Mask	RS	CLM	BD	I	517+(DW2+CM)•115+M•173
	Compare Logical Long	RR	CLCL	0F	I	-1630+C256•5175
						+F256(5980+(L1>L2)•230)
						+(L8F+L8C)•403+MA•L8C•115
						(for C256>0 and F256>0)
						$518 + F256 \cdot (5865 + (L1 > 0) \cdot 230) + L8F \cdot 403$
						(for C256=0 and F256>0)
						518+C256•5175+L8C•403+MA•L8C•115
						(for C256>0 and F256=0)
	Connect Page	RS	CTP	B0		3795-C•2185
	Convert to Binary	RX	CVB	4F		920+NDD2•575+(DW2+X)•115
	Convert to Decimal	RX	CVD	4E		$978 \cdot (1 + \text{NHD}) + (\text{K} + \text{X}) \cdot 115 + \text{DW2} \cdot 173$
	Deconfigure Page	S	DECP	B21B		10293
	Disconnect Page	S	DCTP	B21C		2243+C•9258
	Divide	RR	DR	1D		5003+N2•115+N1•288
	Divide	RX	D	5D		5635+N1•288+(DW2+N2+X)•115
	Divide (long)	RR	DDR	2D		7705+PN•173

	For-	Mne-	Op-	Accura	icy
Instruction Name	mat	monic	Code	Code	Formula and Comments
Divide (long)	RX	DD	6D		$8050 + (DW2 + X) \cdot 115 + PN \cdot 173$
Divide (short)	RR	DER	3D		4255+PN•173
Divide (short)	RX	DE	7D		$4600 + (DW2 + X) \cdot 115 + PN \cdot 173$
Divide Decimal	SS	DP	FD	I	$5444 + (L1 > 8) \cdot 1093 - (L1 = 16) \cdot 230$
					$+(EDW1+DW1+DW2) \cdot 115$
					$+(EDWQ+DWQ+DWR)\cdot 173$
					-(NDD1=0)•115
					(for case NDD1 <ndd2)< td=""></ndd2)<>
					8926+(NDD1-NDD2)•1037
					+DW1•287+DW2•115
					(for case NDD1 \geq NDD2 and L1 \leq 8)
					9844+(1037+LC•115)•(NDD1-NDD2)
					$+(NDD1>15)\cdot 633-(L1=16)\cdot 230$
					$+(DW1+EDW1+DW2)\cdot 115$
					$+(EDWQ+DWQ+DWR)\cdot 173$
					(for case NDD1 \geq NDD2 and L1>8)
					Note: For cases with NDD1 \geq NDD2 times assume
					equal probability of digits 0 through 9.
Edit	SS	ED	DE	I	748+(Cond.Code•0)•115+Sum(EB+SA)
					(Where the sum is taken over all pattern characters.
					Refer Table 2)
Edit and Mark	SS	EDMK	DF	I	1093+(Cond.Code•0)•115+Sum(EB+SA+EMK)
					(Where the sum is taken over all pattern characters.
					Refer Table 2)
Exclusive OR	RR	XR	17		230
Exclusive OR	RX	Х	57		345+(DW2+X)•115
Exclusive OR (character)	SS	XC	D7		$805 + DW1 \cdot 230 + DW2 \cdot 115$ (for $L \le 8$)
					920+L8•518+LZ8•58+MA•(floor of L/8)•120 (for L>8)
					690+L•518+LZ•58 (for operands overlap)
					$1035+L8 \cdot 173$ (for L>8 and ADDR(OP1)=ADDR(OP2))
Exclusive OR (immediate)	SI	XI	97		518
Execute	RX	EX	44		748 + SI (for R1=0)
					978+SI (for R1≠0)
Halve (long)	RR	HDR	24		863+PN•173
Halve (short)	RR	HER	34		748+PN•173
Insert Character	PY	IC	13		230+¥-115
Insert Characters under Mask	RS	ICM	45 RF		575+(IM+DW2)-115
Insert Page Bits	RS RS	IPR	B4		1610+C•403
Insert PSW Key	KJ S	IPK	B20B		1035
Involidate Page Table Entry	S	IDTE	B20B		1055 4025 (2K Page Size)
Invaluate Lage Lable Litty	5	II IL	D221		4025 (2K Page Size)
Insert Storage Key	PP	ISK	00		1093 (for BC S/370 mode)
Insert Storage Key	КК	151	07		1208 (for EC, S/370 mode)
					1208 (for BC ECPS VSE mode)
					1898 (for EC, ECPS:VSE mode)
			4.5		
Load	RR		18		230
Load	RX		58		230+(DW2+X)•115
Load (long)	RR		28		230
Load (long)	RX		68		230+(DW2+X)•115
Load (short)	RR	LER	38		230

Instruction Name mat monic Code Code Formula and Comments Load (short) RX LA 41 230+(DW2+X)+115 Load and Complement (long) RR LTDR 23 460 Load and Test RR LTDR 23 460 Load and Test RR LTDR 23 345 Load and Test RR LTR 32 345 Load Complement RR LCR 13 288 Load Complement (short) RR LCR 3 460 Load Complement (short) RR LCR 13 288 Load Complement (short) RR LCR 8 1553 Load Halfword RX LH 48 288+(DW2+X)-115 Load Negative (long) RR LNR 11 460 Load Negative (short) RR LNR 1460 140 Load Positive (short) RR LPR 10 288+V2:173 Load Positive (short)		For-	Mne-	Op-	Accura	ю
Load (short) RX LA 41 230+(DW2+X)+115 Load and Complement (long) RR LTR 12 230+(DW2+X)+115 Load and Test RR LTR 12 230 Load and Test (long) RR LTBR 22 345 Load and Test (long) RR LTER 32 345 Load Complement (short) RR LCER 33 460 Load Control RS LCTL B7 17250+PU-4715 (for S/370 mode, load all Regs) Load Control RS LCTL B8 15348 (for ECPS:VSE, load all Regs) Load Haffword RX LH 48 288+(DW2+X)+115 Load Megative (short) RR LFI B8 15348 (for ECPS:VSE, load all Regs) Load Megative (long) RR LNR 11 400-N2*VII5 Load Megative (long) RR LNR 1345 Load Negative (long) Load Positive (long) RR LPR 10 288+V2+173 Load Negative (short) RR LPR 10 288+N2+173 Load Positive (long) RR	Instruction Name	mat	monic	Code	Code	Formula and Comments
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$	Load (short)	RX	LE	78		$230 \pm (DW2 \pm X) = 115$
$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	Load Address	RX	LA	41		$230 \pm X \cdot 115$
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$	Load and Complement (long)	RR	LCDR	23		460
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	Load and Test	RR	LTR	12		230
Load and Test (short) RR LTER 32 345 Load Complement (short) RR LCR 13 288 Load Complement (short) RR LCR 33 460 Load Control RS LCTL B7 17250+PU+4715 (for S/370 mode, load all Regs) 16348 (for ECPS:VSE, load all Regs) 3795+PU@4715 (for S/370 mode, load all Regs) 16348 (for ECPS:VSE, load Regs 0,1) 2880 (for ECPS:VSE, load Regs 0,1) 288 (for NO-OP case) Move (character) RR LRDR 25 518 (for NO-OP case) Move (long to short) RR LRDR 25 518 (for NO-OP case) Move (long to short) RR LSER 35 MVC D2 518 (for LS 8 and operands overlap) 748 + L-288 (for M256-0 and P256-0) -460 + P256 - 502 + P256 - 5894 + L80 + L89 + L80 + L80 + L80 + L80 + L80 + L80 + L	Load and Test (long)	RR	LTDR	22		345
Load ComplementRRLCR1328kLoad Complement (short)RRLCR33460Load ControlRSLCTLB717250+PU+4715 (for S/370 mode, load all Regs) 3795+PU@4715 (for S/370 mode, load Regs 0, 1) 2880 (for ECPS:VSE, load All Regs) 3795+PU@4715 (for S/370 mode, load Regs 0, 1)Load Frame IndexRSLFIB81553Load HalfwordRXLH48288+(DW2+X)+115Load MultipleRSLM98115+N+NWB+115Load NegativeRRLNR1403-N2-115Load Negative (short)RRLNR11406Load Positive (short)RRLPR10288+N2-173Load Positive (long)RRLPDR20460Load Positive (short)RRLPER30460Load Real AddressRXLRAB11495Load Real AddressRXLRA11495Load Rounded (extended to long)RRLERR35690Make AddressableSMUNB21E2128+A-345Move (character)SSMVCD2575+DW1-173+DW2+58 (for L<8)	Load and Test (short)	RR	LTER	32		345
Load Complement (short) RR LCER 33 460 Load Control RS LCTL B7 17250+PU-4715 (for S/370 mode, load all Regs) 16348 (for ECPS:VSE, load all Regs) 3795+PU@4715 (for S/370 mode, load Regs 0,1) Load Frame Index RS LFT B8 1553 Load Hafword RX LH 48 288+(DW2+X)+115 Load Multiple RS LM 98 115+N+NWB-115 Load Negative (RR LNR 11 403-N2-115 Load Negative (short) RR LNR 31 460 Load Pestive (short) RR LPDR 20 460 Load Positive (short) RR LPDR 20 460 Load Positive (short) RR LPDR 20 460 Load Positive (short) RR LPR 30 460 Load Positive (short) RR LFR 31 495 Load Nounded (extended to load Positive (short) RR LFR 35 690 Make Addressable S MAD B1 1495 Load Addressable S MAD B21D 2243-A-115 Make Unaddressable S MAD B21D 2243-A-115 Make Unaddressable S MVC D2 575-DW1-173+DW2-58 (for L ≤8) S18+K2-88+MA-(floor L >8-128) Move (character) SS MVC D2 875-DW1-173+DW2-58 (for L ≤8) Move (character) SS MVC D2 875-DW1-173+DW2-58 (for L ≤8) Move Long RR MVC 0E I Move LONG MALESSADE SS MAVE 0E I MOVE LONG MALESSADE SS MAVE 0E I MVC 0E I MOVE LONG MALESSADE SS MALESSADE ASSON MAEOB-4399+L8M-230 +L8P-173+MA-L8M+115 (for M256-0 and P256-0) MOVE: If operand 15 on a 288 bndy, and the pad length is MOVE: If operand 15 on a 285 bndy, and the pad length is Note: If operand 15 on a 285 bndy, and the pad length is Note: If operand 15 on a 285 bndy, and the pad length is	Load Complement	RR	LCR	13		288
Load Control RS LCTL B7 17250+PU-4715 (for S/370 mode, load all Regs) 3795+PU@4715 (for S/370 mode, load Regs 0,1) Load Frame Index RS LFI B8 1553 Load Halfword RX LH 48 288 (for ECPS:VSE, load Regs 0,1) Load Halfword RX LH 48 288 (for ECPS:VSE, load Regs 0,1) Load Multiple RS LM 98 115+N+NWB-115 Load Negative (long) RR LNR 11 403-N2+115 Load Negative (long) RR LNR 14 460 Load Positive (long) RR LPR 10 288+N2+173 Load Real Address RX LRA B1 1495 Load Rounded (extended to - - - - load Rounded (long to short) RR LRR 25 575 Load Rounded (long to short) RR	Load Complement (short)	RR	LCER	33		460
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	Load Control	RS	LCTL	B7		$17250 \pm PU \pm 4715$ (for \$/370 mode, load all Bass)
$\begin{tabular}{ c c c c c c c c c c c c c c c c c c c$				227		16348 (for ECPS:VSE load all Page)
Load Frame Index RS LFI B8 1553 Load Halfword RX LH 48 288 (107 ECPS:VSE, load Regs 0, 1) Load Halfword RX LH 48 288 (107 ECPS:VSE, load Regs 0, 1) Load Multiple RS LM 98 115 +N +NWB+115 Load Negative (long) RR LNR 11 403-N2+115 Load Negative (long) RR LNR 31 460 Load Positive (short) RR LPR 10 288 + N2+173 Load Positive (short) RR LPR 20 460 Load Positive (short) RR LPR 30 460 Load Positive (short) RR LPR 30 460 Load Rega Address RX LRA B1 1495 Load Rounded (extended to 00g) RR LRDR 25 575 Load Rounded (ong to short) RR LRE 35 690 518 (for N-OP case) Make Addressable S MUN B21E 2128+A.345 518 (for N-OP case) Move (character) SS <td></td> <td></td> <td></td> <td></td> <td></td> <td>3795 + PU@4715 (for S/370 mode load Page 0.1)</td>						3795 + PU@4715 (for S/370 mode load Page 0.1)
Load Frame Index RS LFI B8 1553 Load Halfword RX LH 48 288+(DW2+X)+115 Load Multiple RS LM 98 115+N+NWB+115 Load Negative RR LNR 11 403:N2+115 Load Negative (long) RR LNR 31 460 Load Negative (short) RR LNR 31 460 Load Positive (long) RR LPR 0 288+N2+173 Load Positive (long) RR LPR 0 460 Load Positive (short) RR LPR 0 460 Load Rounded (short) RR LPR 30 460 Load Rounded (long to short) RR LRR 35 690 Make Addressable S MAD B21E 2128+A-\$115 Make Unaddressable S MUN B21E 2128+A-\$15 Move (character) SS MVC D2 575+DW1+173+DW2+\$8 (for L \$8) Move (immediate)						2880 (for ECPS:VSE load Page 0.1)
	Load Frame Index	RS	LFI	B 8		1553
$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	Load Halfword	RX	гн	48		$288 \pm (DW2 \pm V) \pm 115$
	Load Multiple	RS	LM	98		$115 \pm N \pm NWP_{-}115$
	Load Negative	RR	INR	11		403 N2.115
$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	Load Negative (long)	RR	INDR	21		245
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	Load Negative (short)	RR	LNER	31		460
$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	Load Positive	RR	L PR	10		
$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	Load Positive (long)	RR		20		460
$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	Load Positive (short)	RR	LPER	30		460
Load Real Address RX LRA B1 1495 Load Rounded (extended to long) RR LRDR 25 575 Load Rounded (long to short) RR LRER 35 690 Make Addressable S MAD B21D 2243-A.115 Make Unaddressable S MUN B21E 2128+A.345 Monitor Call SI MC AF 10707 Move (character) SS MVC D2 575+DW1.173+DW2.58 (for L \leq 8) 518+L8.288+MA.(floor of L/8).115 (for L>8) 633+L.288 (for L \leq 8 and operands overlap) 748+L.288 (for L \leq 8 and operands overlap) 748+L.288 (for L \geq 8 and operands overlap) Move linverse SS MVCI E8 345+(L.288) Move Long RR MVCL 0E I 288+M256.5262+P256.5894-MEOB.4399+L8M.230 +L8P.173+MA.L8M.115 (for M256>0 and P256>0) 230+M256-5262+MEOB.115+L8M.288+MA.L8M.115 (for M256-50 and P256=0) -460+P256.5894+L8P.173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is	Load PSW	S	LI EK	82	т	2000
Load Rounded (extended to long) RR LRDR 25 575 Load Rounded (long to short) RR LRER 35 690 Make Addressable S MAD B21D 2243-A \cdot 115 Make Unaddressable S MUN B21E 2128+A \cdot 345 Monitor Call SI MC AF 10707 Move (character) SS MVC D2 575+DW1 \cdot 173+DW2 \cdot 58 (for L \leq 8) 518+L8 \cdot 288+MA \cdot (floor of L/8) \cdot 115 (for L>8) 633+L \cdot 288 (for L \leq 8 and operands overlap) 748+L \cdot 288 (for L \geq 8 and operands overlap, OP1 \neq OP2 $+$ 1) 1323+L8 \cdot 173 (for L>8 and operands overlap, OP1 \neq OP2 $+$ 1) Move Inverse SS MVCI E8 345+(L \cdot 288) Move Long RR MVCL 0E I 288+M256 \cdot 5262 $+$ P256 \cdot 5894-MEOB \cdot 4399+L8M \cdot 230 +L8P \cdot 173+MA \cdot L8M \cdot 115 (for M256 \circ 0 and P256 \circ 0) 230+M256 \cdot 5262 $+$ MEOB \cdot 115+L8M \cdot 288+MA \cdot L8M \cdot 115 (for M256 \circ 0 and P256 \circ 0) -460+P256 \cdot 5894+L8P \cdot 173 (for M256 \circ 0 and P256 \circ 0) Note: If operand 1 is on a 2KB bndy, and the pad length is	Load Real Address	RX	IRA	81	1	1495
	Load Rounded (extended to		LINI	ы		1775
Load Rounded (long to short) RR LRER 35 690 Make Addressable S MAD B21D 2243-A+115 Make Unaddressable S MUN B21E 2128+A+345 Monitor Call SI MC AF 10707 Move (character) SS MVC D2 575+DW1+173+DW2+58 (for L \leq 8) 518+L8+288+MA+(floor of L/8)+115 (for L>8) 633+L+288 (for L \leq 8 and operands overlap) 748+L+288 (for L \leq 8 and operands overlap) 748+L+288 (for L \leq 8 and operands overlap) 748+L+288 (for L \leq 8 and operands overlap, OP1 \neq OP2+1] 1323+L8+173 (for L>8 and operands overlap, OP1 $=$ OP2+ Move (immediate) SI MVI 92 288 Move Long RR MVCL 0E I 288+M256+5262+P256+5894-MEOB+4399+L8M+230 +L8P+173+MA+L8M+115 (for M256>0 and P256>0) 230+M256+5262+MEOB+115+L8M+288+MA+L8M+115 (for M256>0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is	long)	RR		25		575
$\begin{array}{llllllllllllllllllllllllllllllllllll$	Load Rounded (long to short)	RR	LRER	35		690
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	Make Addressable	s	MAD	D 21D		2242 4 115
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	Make Unaddressable	3 6	MUN	B21D B21E		2243-A•115
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	Marc Onaddressable	5	MON	BZIE		2128+A•345
Move (character) SS MVC D2 575+DW1•173+DW2•58 (for L≤8) 518+L8•288+MA•(floor of L/8)•115 (for L>8) 633+L•288 (for L≤8 and operands overlap) 633+L•288 (for L≥8 and operands overlap, OP1≠OP2+1) Move (immediate) SI MVI 92 288 Move Inverse SS MVCIN E8 345+(L•288) Move Long RR MVCL 0E 1 288+M256•5262+P256•5894-MEOB•4399+L8M•230 +L8P•173+MA•L8M•115 (for M256>0 and P256>0) 230+M256•5262+MEOB•115+L8M•288+MA•L8M•115 (for M256>0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is Note: If operand 1 is on a 2KB bndy, and the pad length is	Monitor Can	51	MC	Аг		10/07
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	Move (character)	66	MVC	Da		518 (Ior NO-UP case)
$S18+L8+288+MA \cdot (1000 \text{ of } L/8) \cdot 115 \text{ (for } L>8)$ $633+L-288 \text{ (for } L \le 8 \text{ and operands overlap)}$ $748+L \cdot 288 \text{ (for } L \ge 8 \text{ and operands overlap, } OP1 \neq OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } L>8 \text{ and operands overlap, } OP1 = OP2 + 1)$ $1323+L8 \cdot 173 \text{ (for } M256 \cdot 5262 + P256 \cdot 5894 - MEOB \cdot 4399 + L8M \cdot 230 + 1288 + MA \cdot L8M \cdot 115 \text{ (for } M256>0 \text{ and } P256>0)$ $230+M256 \cdot 5262+MEOB \cdot 115+L8M \cdot 288+MA \cdot L8M \cdot 115 \text{ (for } M256>0 \text{ and } P256=0)$ $-460+P256 \cdot 5894+L8P \cdot 173 \text{ (for } M256=0 \text{ and } P256=0)$ $-460+P256 \cdot 5894+L8P \cdot 173 \text{ (for } M256=0 \text{ and } P256>0)$ Note: If operand 1 is on a 2KB bndy, and the pad length is	move (character)	33	M V C	DZ		$5/5+Dw1 \cdot 1/3+Dw2 \cdot 58$ (for $L \le 8$)
633+L-288 (for L≤8 and operands overlap) 748+L-288 (for L>8 and operands overlap, OP1≠OP2+1) 1323+L8+173 (for L>8 and operands overlap, OP1=OP2+ Move (immediate) SI MVI 92 288 Move Long RR MVCL 0E 1 288+M256+5262+P256+5894-MEOB+4399+L8M+230 +L8P+173+MA+L8M+115 (for M256>0 and P256>0) 230+M256+5262+MEOB+115+L8M+288+MA+L8M+115 (for M256>0 and P256=0) -460+P256+5894+L8P+173 (for M256=0 and P256=0) -460+P256=0 Note: If operand 1 is on a 2KB bndy, and the pad length is						$518+L8 \cdot 288+MA \cdot (tloor of L/8) \cdot 115$ (for L>8)
748+L-288 (for L>8 and operands overlap, OP1≠OP2+1) 1323+L8.173 (for L>8 and operands overlap, OP1=OP2+ Move (immediate) SI MVI 92 288 Move Inverse SS MVCIN E8 345+(L.288) Move Long RR MVCL 0E 1 288+M256.5262+P256.5894-MEOB.4399+L8M.230 +L8P.173+MA.L8M.115 (for M256>0 and P256>0) 230+M256.5262+MEOB.115+L8M.288+MA.L8M.115 (for M256>0 and P256=0) -460+P256.5894+L8P.173 (for M256=0 and P256=0) -460+P256.5894+L8P.173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is Note: If operand 1 is on a 2KB bndy, and the pad length is						$633+L \cdot 288$ (for $L \le 8$ and operands overlap)
Move (immediate) SI MVI 92 288 Move Inverse SS MVCIN E8 345+(L•288) Move Long RR MVCL 0E 1 288+M256•5262+P256•5894-MEOB•4399+L8M•230 +L8P•173+MA•L8M•115 (for M256>0 and P256>0) 230+M256•5262+MEOB•115+L8M•288+MA•L8M•115 (for M256>0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256>0) -460+P256•5894+L8P•173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is Note: If operand 1 is on a 2KB bndy, and the pad length is Note: If operand 1 is on a 2KB bndy, and the pad length is						$748+L \cdot 288$ (for L>8 and operands overlap, $OP1 \neq OP2+1$)
More (miniculate) 31 MVT 92 288 Move Inverse SS MVCIN E8 345+(L•288) Move Long RR MVCL 0E 1 288+M256•5262+P256•5894-MEOB•4399+L8M•230 +L8P•173+MA•L8M•115 (for M256>0 and P256>0) 230+M256•5262+MEOB•115+L8M•288+MA•L8M•115 (for M256>0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is Note: If operand 1 is on a 2KB bndy, and the pad length is	Move (immediate)	C1	MAX/I	0.2		$1323+L8 \cdot 1/3$ (for L>8 and operands overlap, OP1=OP2+1)
Move Long 33 MVCIN E8 343+(1.2288) Move Long RR MVCL 0E 1 288+M256.5262+P256.5894-MEOB.4399+L8M.230 +L8P.173+MA.L8M.115 (for M256>0 and P256>0) 230+M256.5262+MEOB.115+L8M.288+MA.L8M.111 (for M256>0 and P256=0) -460+P256.5894+L8P.173 (for M256=0 and P256>0) .460+P256.5894+L8P.173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is	Move (minediate)	51	MVCIN	92 E9		288
Model Long KK Model OE 1 288+M256+5262+P256+5894-MEOB+4399+L8M+230 +L8P+173+MA+L8M+115 (for M256>0 and P256>0) 230+M256+5262+MEOB+115+L8M+288+MA+L8M+115 (for M256>0 and P256=0) -460+P256+5894+L8P+173 (for M256=0 and P256=0) -460+P256+5894+L8P+173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is	Move Long	33 DD	MVCIN	E8 OF		345+(L+288)
+L8P•173+MA•L8M•115 (for M256>0 and P256>0) 230+M256•5262+MEOB•115+L8M•288+MA•L8M•115 (for M256>0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is	Move Long	ĸĸ	MVCL	UE	1	288+M256•5262+P256•5894-MEOB•4399+L8M•230
(for M256>0 and P256>0) 230+M256•5262+MEOB•115+L8M•288+MA•L8M•11: (for M256>0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is						$+L\delta P \cdot 1/3 + MA \cdot L\delta M \cdot 115$
(for M256=0 and P256=0) -460+P256•5894+L8P•173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is						(101 M250>0 and P250>0)
-460+P256•5894+L8P•173 (for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is						230 + M250 + 3202 + MEOB + 113 + L8M + 288 + MA + L8M + 115 (for M256 > 0 and P256 = 0)
(for M256=0 and P256>0) Note: If operand 1 is on a 2KB bndy, and the pad length is						(101 M250>0 and P250=0)
Note: If operand 1 is on a 2KB bndy, and the pad length is						(for M256 = 0 and B256 > 0)
Note. If operating 1 is on a 2KB birdy, and the pad length is						(for M250=0 and P250>0) Note: If opproved 1 is one $2KB$ had a subtract the subtract is
dependent them are equal to 2K, this instance in the						store in operand 1 is on a 2KB bindy, and the pad length is
than the formula shows. Altheurs all shows at the state of the state o						than the formula shows. Although alcourse the set
reduction in cache missee and as a result there is a						reduction in cache misses and as a result there is a
not performance improvement						net performance improvement

	For-	Mne-	Op-	Accura	cy
Instruction Name	mat	monic	Code	Code	Formula and Comments
Move Numerics	SS	MVN	D1		1265+DW1•230+DW2•115 (for L≤8)
					863+L8•748+MA•(floor of L/8)•115 (for L>8)
					633+L•748+DW1•58
					(for $L \leq 8$ and operands overlap)
					575+L•748
					(for L>8 and operands overlap)
Move with Offset	SS	MVO	F1		1725+(2.5•DW1+DW2)•115+(A2 <a1+l1)•288< td=""></a1+l1)•288<>
					(for $L1 \leq 8, L2 \leq 8$)
					1840+(2.5•DW1+DW2+EDW2) •115+(A2 <a1+l1)•288< td=""></a1+l1)•288<>
					;(for $L1 \le 8, L2 > 8$)
					2243+(2.5•EDW1+2.5•DW1+DW2)•115+(A2 <a1+l1)•288< td=""></a1+l1)•288<>
					$(for L1 > 8, L2 \le 8)$
					2243+(2.5•EDW1+2.5•DW1+DW2+EDW2)
					•115+(A2 <a1+l1)•288< td=""></a1+l1)•288<>
					(for L1>8,L2>8)
					$2300 + [6.5 \cdot min(L1,L2) + 3.5(L1 > L2) + 1.5$
					(L1>L2+8)+(L1>8)+(L2>8)] 115
					(for overlapping case)
Move Zones	SS	MVZ	D3		$1265 + DW1 \cdot 230 + DW2 \cdot 115$ (for L ≤ 8)
					863+L8•748+MA•(floor of L/8)•120 (for L>8)
					633+L•748+DW1•58
					$(L \leq 8 \text{ and operands overlap})$
					633+L•748 (L>8 and operands overlap)
Multiply	RR	MR	1C		2760+N2•345
					805 (for VP=0)
Multiply	RX	М	5C		2933+(3•N2+DW2+X)•115
Multiply (extended)	RR	MXR	26		11922+PN•58
Multiply (long to extended)	RR	MXDR	27		4485+PN•58
Multiply (long to extended)	RX	MXD	67		4830+PN•58+(DW2+X)•115
Multiply (short to long)	RR	MER	3C		2933+PN•173
Multiply (short to long)	RX	ME	7C		3278+PN•173+(DW2+X)•115
Multiply (long)	RR	MDR	2C		3738+PN•173
Multiply (long)	RX	MD	6C		4083+PN•173+(DW2+X)•115

	For-	Mne-	Op-	Accuracy
Instruction Name	mat	monic	Code	Code Formula and Comments
Multiply Decimal	SS	МР	FC	$3565 + (3.5 \cdot DW1 + DW2) \cdot 115 + (L1 > 8)$ $\cdot (2 + FDW1 - DW1 + 2.5 \cdot EDW1) \cdot 115 - (NDD2 = 0) \cdot 113$ (for NDD1 = 0) $4140 + (FDW1 + DW2 + 2.5 \cdot EDW1 + 2.5 \cdot DW1) \cdot 115$ (for NDD1 > 15 and NDD2 = 0) $3738 + (3.5 \cdot DW1 + DW2) \cdot 115 + (L1 > 8)$ $\cdot (5.5 + FDW1 - DW1 + 2.5 \cdot EDW1) \cdot 115 - CLZ \cdot 403$ (for NDD1 \leq 15 and NDD2 = 0) $6095 + (NDD2 + 3.5 \cdot DW1 + DW2) \cdot 115$ $+ (L1 > 8) \cdot (4.5 + FDW1 - DW1 + 2.5 EDW1)$ $\cdot 115 - CLZ \cdot 403$ (for NDD1 \leq 15 and NDD2 > 0) $8683 + (3.5 \cdot DW1 + DW2) \cdot 115 + (L1 > 8)$ $\cdot (4.5 + FDW1 - DW1 + 2.5 \cdot EDW1) \cdot 115 - CLZ \cdot 403$
				(for NDD1=15 and NDD2<15) 13915+(3.5•DW1+DW2)•115+(L1>8) •(4.5+FDW1-DW1+2.5•EDW1)•115-CLZ•403 (for NDD1=15 and NDD2=15) 10465+(2•NDD2+FDW1+DW2
				+2.5•EDW1+2.5•DW1)•115 (for NDD1>15 and NDD2>0) 4524+(4•VP+FDW1+DW2 +2.5•EDW1+2.5DW1)•115 (for NDD1=29 and NDD2=1)
Multiply Halfword	RX	MH	4C	2358+(DW2+N2+X)•115
OR	RR	OR	16	230
OR	RX	0	56	345+(DW2+X)•115
OR (character)	SS	OC	D6	$805+DW1\cdot230+DW2\cdot115$ (for L ≤ 8) $690+L8\cdot518+MA\cdot$ (floor of L/8) $\cdot115+LZ8\cdot58$ (for L>8) $690+L\cdot518+LZ\cdot58$ (for even the state)
OR (immediate)	SI	OI	96	518
Pack	SS	PACK	F2	$1265 + DW2 \cdot 115 + DW1 \cdot 173$ (for L1 ≤ 8 , L2 ≤ 8) $1553 + DW2 \cdot 115 + (DW1 + EDW1) \cdot 173$ (for L1 > 8 , L2 ≤ 8) $1725 + DW1 \cdot 173 + (DW2 + EDW2) \cdot 115$ (for L1 ≤ 8 , L2 > 8) $2013 + (DW1 + EDW1) \cdot 173 + (DW2 + EDW2) \cdot 115$ (for L1 > 8 , L2 > 8)
Purge TLB	S	PTLB	B20D	6268
Reset Reference Bit	S	RRB	B213	2300+C•115 (for ECPS:VSE mode) 1325 (for S/370 mode)
Retrieve Status and Page	SS	RSP	D8	805
Set Clock	S	SCK	B204	5520 (for Enable ON) 1322 (for Enable OFF)
Set Clock Comparator	S	SCKC	B206	3910
Set CPU Timer	S	SPT	B208	4140
Set Page Bits	RS	SPB	B5	2473+(C•230)
Set Program Mask	RR	SPM	04	575
et PSW Key from Address	c	SDV A	B20 A	1660

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	For-	Mne-	Op-	Accura	cy
Instruction Name	mat	monic	Code	Code	Formula and Comments
	DD	CCV	08		7130 (for $\frac{5}{370}$ mode)
Set Storage Key	ĸĸ	221	00		3284 (for ECPS:VSE mode)
	0	CCM	90		1150 (for BC: \$/370 mode or ECPS:VSE mode)
Set System Mask	8	55M	80		1725 (for EC) ECPS:VSE mode: turn I/E bits off or on)
					1/25 (for EC, ECF3, V3E mode, turn 1/E bits off or on)
					1610 (for EC; $S/370$ mode, turn T/E bits off or on)
					1840 (for EC; S/ 370 mode; turn 1 bit off of off)
Shift and Round Decimal	SS	SRP	F0		2645-ZR•1/3+DW1•288
					(for $L \leq 8$, shift left and amount < 16)
					2415-ZR•115+DW1•288
					(for $L \leq 8$, shift right and amount < 16)
					1725+DW1•288
r					(for $L \leq 8$, shift left and amount ≥ 16)
					1552+DW1•288
					(for $L \leq 8$, shift right and amount>16)
					1725+DW1•288
					(for $L \leq 8$, shift right and amount=16)
					4255+DW1•690-ZR•115
					(for L>8, shift left)
					3737+DW1•690-ZR•115
					(for L>8, shift right and amount<16)
					3507+DW1•690-ZR•115
					(for $L>8$, shift right and amount=16)
					3622+DW1•690-ZR•115
					(for L>8, shift right and 16 <amount<32)< td=""></amount<32)<>
					3392+DW1•690-ZR•115
					(for $L>8$, shift right and amount=32)
Shift Laft Double	RS	SLDA	8F		1092+N1•575
Shift Laft Double Logical	RS	SLDL	8D		920
Shift Laft Single	RS	SLA	8B		747+N1•115
Shift Left Single Logical	RS PS	SUI	89		460
Shift Disht Darkla	R5 DC	SPDA	8E		920
Shift Right Double	DC	SRDA	80		920
Shift Right Double Logical	R3 DC	SEDL	84		578
Shift Right Single	R3 DC	CDI	0A 00		460
Shift Right Single Logical	K5 DV	SKL	00 50		$288 \pm DW_{2} = 173 \pm X = 115$
Store	KA DV	STD	60		$288 \pm DW^{2} \cdot 173 \pm X \cdot 115$
Store (long)	KA DV	STD	70		$288 \pm DW2 \cdot 173 \pm X \cdot 115$
Store (short)	KA C	SIE	70 D010	7	1/28
Store Capacity Counts	3	SICAP	42		1450 200 - V - 115
Store Character	RX	SIC	42 DE		$200 \pm A \cdot 115$
Store Characters Under Mask	RS	SICM	BE	. I	518+CM•115+Dw2•175
Store Clock	S	STCK	B205	,	1610
Store Clock Comparator	S	STCKC	B207		
Store Control	RS	STCTL	B 6		690+[ORS+ceiling of (N-OR)/2]0288
					+(floor of N-ORS/2)•RSA•1/3
Store CPU ID	S	STIDP	B202	2	1093
Store CPU Timer	S	STPT	B209)	1898
Store Halfword	RX	STH			288+DW2•173+X•115
Store Multiple	RS	STM	90		115+N•173+NWB•173
Store then AND System Mask	SI	STNSM	AC		1380 (for BC; S/370 mode or ECPS:VSE mode
					1495 (for EC; ECPS:VSE mode; turn I/E bits off)
					1495 (for EC; S/370 mode; turn I/E bits off)
					1783 (for EC; S/370 mode; turn T bit off)

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	For-	Mne-	Op-	Accura	су
Instruction Name	mat	monic	Code	Code	Formula and Comments
Store then OR System Mask	SI	STOSM	AD		1380 (for BC; S/370 mode or ECPS:VSE mode
					1495 (for EC; ECPS:VSE mode: turn I/E bits on
					1495 (for EC; S/370 mode; turn I/E bits on)
					1783 (for EC; $S/370$ mode; turn T bit on)
Subtract	RR	SR	1B		288
Subtract	RX	S	5B		403+(DW2+X)•115
Subtract Decimal	SS	SP	FB		863+RC•403+ZR•805+(2.5•DW1+DW2)•115
					(for $L1 \leq 8$ and $L2 \leq 8$)
					1265+RC•518+ZR•805
					+(2.5•DW1+DW2+EDW2)•115
					(for $L1 \leq 8$ and $L2 > 8$)
					1840+RC•518+ZR•805+(2.5•DW1
					+2.5•EDW1+DW2)•115
					(for L1>8 and L2 \leq 8)
					1840+RC•460+ZR•805+(2.5•DW1
					+2.5•EDW1+DW2+EDW2)•115
					(for L1>8 and L2>8)
Subtract Halfword	RX	SH	4 B		460+(DW2+X)•115
Subtract Logical	RR	SLR	1F		230
Subtract Logical /	RX	SL	5F		345+(DW2+X)•115
Subtract Normalized					
(extended)	RR	SXR	37	I	2408
Subtract Normalized (long)	RR	SDR	2B	I	765
Subtract Normalized (long)	RX	SD	6B	Ι	1053+(DW2+X)•115
Subtract Normalized (short)	RR	SER	3B	I	826
Subtract Normalized (short)	RX	SE	7 B	I	$1028 + (DW2 + X) \cdot 115$
Subtract Unnormalized (long)	RR	SWR	2F	I	720
Subtract Unnormalized (long)	RX	SW	6F	I	$1064 + (DW2 + X) \cdot 115$
Subtract Unnormalized (short)	RR	SUR	3F	I	841
Subtract Unnormalized (short)	RX	SU	7F	I	$1243 + (DW2 + X) \cdot 115$
Supervisor Call	RR	SVC	0A	I	4140 (BC mode)
					3680 (EC mode)
Test and Set	S	TS	93		747
Test under Mask	SI	ТМ	91		288
Test Protect	SS	TPRT	E501		3853 Condition Code 0
					4428 Condition Code 1
					4428 Condition Code 2
					3278 Condition 3
Translate	SS	TR	DC		460+L•633
Translate and Test	SS	TRT	DD		460+NP•633+CCV•115
Unpack	SS	UNPK	F3		1610+DW2•115+DW1•173
					(for $L1 \le 8$, $L2 \le 8$)
					2013+DW2•115+(DW1+EDW1)•173
					(for L1>8, L2≤8)
					1840+DW1•173+(DW2+EDW2)•115
					(for L1≤8, L2>8)
					2358+(DW2+EDW2)•115+(DW1+EDW1)•173

(for L1>8, L2>8)

Instruction Name mat mat mat Code Code Fermitik and Commons Zero and Add SS ZAP F8 1840+RC-403+ZR+403+(15-DW1+DW2) 115+(A2 <a1+l1)-288< td=""> (for L158 and L258) 200=RC-318+ZR-518+(15-DW1+LDW2+DW2) 115+(A2<a1+l1)-288< td=""> (for L158 and L258) 200=RC-318+ZR-518 (for L158 and L258) 205=RC-518+ZR-518 (for L158 and L258) 2155-RC-518+ZR-518 (for L158 and L258) 205=RC-518+ZR-518 (for L158 and L258) Legend for Timing Formulas (for L158 and L258) (for L158 and L258) (for L258 and L258) Legend for Timing Formulas (for L158 and L258) (for L258 and L258) (for L258 and L258) Legend for Timing Formulas (for C 152 and L256) (for L258 and L258) (for L258 and L258) Legend for Timing Formulas (for C 152 and L256) (for L258 and L256) (for L258 and L256) A1 address of operand 1. (for C 152 and C 258) (for L258 and L254) (for L258 and L254) C26 1 if multiplicand has 8 bytes or more of leading zeros: operand 1. (for C 152 and L254) (for L258 and L254) C17. 1 if multip</a1+l1)-288<></a1+l1)-288<>				For-	Mne-	Op-	Accura	сy				
 Zero and Add SS ZAP FR (840+RC-403+(1.5.DW1+DW2) -115+(1.25 and 1.25) 2300+RC-515+ZR-518+(1.5.DW1+DW2+DW2) -115+(1.25 and 1.25) 2300+RC-515+ZR-518+(1.5.DW1+DW2+DW2) -115+(1.25 and 1.25) 238+RC-515+ZR-518+(1.5.DW1+DW2+DW2) -115+(1.25) -115+(Instructio	n Na	me	mat	monic	Code	Code	For	rmula a	nd C	Comments	
Zero and Add SS ZAP FS 1800+RC-0042 (X+005+1(3-100 + 100*2) 1 State State State State 2000 RC-05142 (X+015+11-288 (for L158 and L258) 2358+RC-0518+2(R-018) 1 State State State State 1 State State State State State 2 address of operand 1. State State State State State 1 State State <td></td> <td></td> <td></td> <td>66</td> <td>7.4.0</td> <td>Ea</td> <td></td> <td>1040 - 1</td> <td>D.C. 40</td> <td>a</td> <td>ZD 402 - (1.5. DW(1 DW(2)</td> <td></td>				66	7.4.0	Ea		1040 - 1	D.C. 40	a	ZD 402 - (1.5. DW(1 DW(2)	
 (for L1 \$ and L2 \$8) 200+RC-518+2RA-518+(1.5-DW1+EDW2+DW2) (15+128-874-12-288 (for L158 and L2>8) 2258-RC-518+2RA-518 (for L158 and L2>8) (for L158 and L159 (for L158 and L159 (for L158 and L159 (for L158 and L150 file fint (L1-8) bytes of operand 1. (for L158 and L150 file fint (L1-8) bytes of operand 1. (for L158 and L2) file file fint (L1-8) bytes of operand 1. (for L158 and L150 file fint (L1-8) bytes of operand 1. (for L158 and L150 file fint (L1-8) bytes of operand 1.	Zero and	Add	1	55	ZAP	F8		1840+1 •115+(RC•40 (A2 <a< td=""><td>(1+)</td><td>$LR \cdot 403 + (1.3 \cdot D W 1 + D W 2)$</td><td></td></a<>	(1+)	$LR \cdot 403 + (1.3 \cdot D W 1 + D W 2)$	
 2000 RC-134 / ZR-134 / L-30W1+EDW2+DW2) 115+42C-41-11)-288 (for L153 and L2>8) 2358-RC-134 / ZR-138 41.5-EDW1+1.5-DW1+DW2)+115 41.42C-41-EL1)-288 (for L153 and L2>8) 2358-RC-314 / ZR-318 41.5-EDW1+1.5								(for	$r L1 \leq 8$	3 and	d L2≤8)	
 113+(A2CA1+L1)-288 (for L158 and L258) 2358+RC-318+ZR-318 (13+EDW1+5+DW1+DW2)+115 +(A2CA1+L1)-288 (for L158 and L258) 2358+RC-318+ZR-318 (15+EDW1+5+DW1+EDW2) +115+(A2CA1+L1)-288 (for L158 and L258) 2358+RC-318+ZR-318 (15+EDW1+DW2+1.5+EDW1+EDW2) +115+(A2CA1+L1)-288 (for L158 and L258) (15+EQM1+DW2+1.5+EDW1+EDW2) +115+(A2CA1+L1)-288 (for L158 and L258) (15+EQM1+DW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1.5+EDW1+EDW2+1								2300+I	- RC•51	8+2	$ZR \cdot 518 + (1.5 \cdot DW1 + EDW2 + DW2)$	
 (for L1 sk and L2-8) 2358+RC-518+RC-518 (for L1-5k and L2-8) 2158+RC-518+RC-518 2169 2160+RC-518+RC-518 2161+RC-518+RC-518 2161+RC-528+RC-518 2161+RC-528+RC-518 2161+RC-528+RC-518 2161+RC-528+RC-518 2161+RC-5								•115+((A2 <a< td=""><td>1+1</td><td>L1)•288</td><td></td></a<>	1+1	L1)•288	
 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+1 - 5 - DW 1+ DW 2): 115 +(A2 < A1 + L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 2358 + RC - 518 + ZR - 518 +(1.5 - DW 1+L): - 288 (or L1 > 8 and L2 < 8) 240 = 116 block is connected: les 0. 250 = number of 256 - byte blocks, on 256 - byte boundaries, spanned by that portion of the longer operand this compared to the pad character, sexinding through the first doubleword, if any, which compared to the pad character, if any the pad character, if any which compared to the pad character, if any a								(for	r L1≤8	3 and	d L2>8)	
 H(1.5 EDW1+15-2W1+DW2)+115 H(2<2A1+L))-288 ((or L)>8 and L2<8) 2358+RC-518+RC+518+RC+518 H(3-5DW1+DW2+1.5 -EDW1+EDW2) +115+(A2<a1+l))-288 ((or="" l)="">8 and L2<8)</a1+l))-288> EDW1-DW2+1.5 -EDW1+EDW2) +115+(A2<a1+l))-288< li=""> ((or L)>8 and L2<8) EDW2 = 1 for L2.5 kind L2<16 if the first (1.2-8) bytes of operand 1. EDW2 = 1 for L2.5 kind L2<16 if the first (1.2-8) bytes of operand 2. EDW3 = 1 if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = length of the operand 1. EZ = if mutiplicand has 8 bytes or more of leading zeros; else 0. EL = length of the operand in bytes. LC = 0 if NDD1 > 15 and NDD1-NDD2 > 14 LC = 10 for L2.5 if operand 2 cross a doubleword boundary, else 0. EDW3 = 1 if first 8 bytes or operand 1 cross a doubleword boundary, else 0. EDW4 = 1 if the steight bytes of operand 1 cross a doubleword boundary, else 0. EDW4 = 1 if the steight bytes of operand 1 cross a doubleword boundary, else 0. EDW4 = 1 if the result a crosses a doubleword boundary, else 0. EDW4 = 1 if the result else obubeword boundary, else 0. EDW4 = 1 if the result else obubeword boundary, else 0. EDW5 = 1 if the goutient cross a doubleword boundary, else 0. EDW5 = 1 if the goutient cross a doubleword boundary, else 0. EDW5 = 1 if the goutient cross a doubleword boundary, else 0. EDW5 = 1 if the goutient cross a doubleword boundary, else 0. EDW5 = 1 if the gouti</a1+l))-288<>								2358+I	RC•51	8+2	ZR•518	
 H(AZA1+L1)-288 (for L1>8 and L28) 2358+RC-518+ZR-518 H(1-DV2-15 + EDW1+EDW2) +15+(AZA1+L1)-288 (for L1>8 and L28) Edgend for Timing Formulas A = 1 if block is addressable; else 0. AB = 1 if block is addressable; else 0. AB = 1 if block is address of operand 2. C = 1 if block is connected; else 0. CC = 1 if block is connected; else 0. CC = 1 if block is connected; else 0. CC = 1 if block is connected; else 0. CC = 1 if block is connected; else 0. CC = 1 if block is connected; else 0. CC = 1 if block is connected; else 0. CC = 1 if block is connected; else 0. CC = 0 if condition code (CC) is 0 5 if CC is 1. CC = 1 if block is connected; else 0. CC = 0 if condition code (CC) is 0. 5 if CC is 2. C1.7 = 1 if multiplicand has 8 bytes or more of leading zeros: else 0. C256 = number of 256-byte blocks, on 256-byte								+(1.5•)	EDW1	+1.	5•DW1+DW2)•115	
Legend for Timing Formulas FUND 1258 and L258 of the first (L2-8) bytes of operad 12 cores a doubleword boundary, else 0. A = 1 if block is addressible; else 0. AB = 1 if block is addressible; else 0. AI = address of operand 1. A2 = address of operand 2. CV = 0 if continion code (CC) is 0 S if CC is 1. 6 if CC is 2. CUZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. CLZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. CLZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. CLZ = 1 if for L1>8 if looperand 1 cores a doubleword boundary, else 0. CLZ = 1 if cores 1 alignment. CLZ = 1 if for L1>8 if looperand 1 cores a doubleword, if any, which compares unequal. MW1 = 1 for L1>8 if looperand 1 coress a doubleword boundary, else 0. DW1 = 1 for L1>8 if borgerand 1 coress a doubleword boundary, else 0. DW2 = 1 for L1>8 if the bast eight bytes of operand 1 coress a doubleword boundary, else 0. DW2 = 1 for L1>8 if the last eight bytes of operand 2 coress a doubleword boundary, else 0. DW2 = 1 for L1>8 if the last eight bytes of operand 2 coressa a doubleword boundary, else 0. <td< td=""><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td>+(A2<</td><td>(AI+L</td><td>.1)•: </td><td>288 d 1 2 < 8)</td><td></td></td<>								+(A2<	(AI+L	.1)•: 	288 d 1 2 < 8)	
 H13-DW1+DW2+15-EDW1+EDW2) H13-K42<a1+l1)-288< li=""> H14-K42<a1+l1)-288< li=""> H14-K42 </a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<></a1+l1)-288<>								2358+1	$RC \cdot 51$	8 ± 7	ZE • 518	
115+(A2 <a1+l1)-288< th="">Legend for Timing FormulasA= 1 if black is addressable; else 0.AB= 1 if branch target is on doubleword boundary, else 0.AI= address of operand 1.A2= address of operand 2.C= 1 if black is addressable; else 0.C= 1 if black is addressable; else 0.CCV= 0 if condition code (CC) is 0Si if CC is 1.EDW2CLZ= 1 if multipleand has 8 bytes or more of leading zeros; else 0.CK2= number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand that is compared to the shorter operand, extending through the first doubleword, if any, which compares unequal. If operands have equal length, this refers to operand 1 alignment.CM= Ste Figure C-2.CM= Ste Figure C-2.DW1= 1 for L258 if doerand 2 crosses a doubleword boundary, = 1 for L258 if doerand 2 crosses a doubleword boundary, = 1 for L258 if doerand 2 crosses a doubleword boundary, = 1 for loces a doubleword boundary, else 0.DW2= 1 if first 8 bytes of operand 1 crosse a doubleword boundary, else 0.DW3= 1 if first 8 bytes of operand 2 crosses a doubleword boundary, = 1 if or loces a doubleword boundary, else 0.DW4= 1 if first 8 bytes cross a doubleword boundary, else 0.DW4= 1 if first 8 bytes cross a doubleword boundary, else 0.DW4= 1 if first 8 bytes cross a doubleword boundary, else 0.DW4= 1 if first 8 bytes cross a doubleword boundary, else 0.DW4= 1 if first 8 bytes cross a doubleword</a1+l1)-288<>								+(1.5•]	DW1+	DW	/2+1.5•EDW1+EDW2)	
(for L1>8 and L2>8)Legend for Timing FormulasA= 1 if black is addressable; ide 0.A= 1 if black is addressable; ide 0.AB= 1 if brack target is on doubleword boundary, else 0.AI= address of operand 1.A2= address of operand 2.CV= 0 if condition code (CC) is 0S if CC is 1.= caponent corresponding to operand 1.C2= 1 if multiplicand has 8 bytes or more of leading zeros;else 0.= 1 if multiplicand has 8 bytes or more of leading zeros;else 0.= 1 if multiplicand has 8 bytes or more of leading zeros;else 0.= 1 if multiplicand has 8 bytes or more of leading zeros;else 0.= 1 if multiplicand has 8 bytes or more of leading zeros;else 0.= 1 if multiplicand has 8 bytes or operand 1 cores a doubleword, if any, which compared to the shorter operand, which is actually compared to the shorter operand, which compares unequal. If operands have equal length, this refers to operand 1 coreses a doubleword boundary, else 0.DW1= 1 for L128 if idperand 1 crosses a doubleword boundary, else 0.DW2= 1 for L228 if it he last eight bytes of operand 1 crosses a doubleword boundary, else 0.DW2= 1 for (length of quotient) \$ if the quotient crosses a doubleword boundary, else 0.DW3= 1 for (length of quotient) \$ if the last eight bytes of operand 2 crosses a doubleword boundary, else 0.DW4= 1 for L28 if the last eight bytes of operand 2 crosses a doubleword boundary, else 0.DW2= 1 for (length of quotient) \$ if the quotient crosses a doubleword boundary, else 0.DW4 <td< td=""><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td>•115+(</td><td>(A2<!--</td--><td>1+</td><td>L1)•288</td><td></td></td></td<>								•115+((A2 </td <td>1+</td> <td>L1)•288</td> <td></td>	1+	L1)•288	
Legend for Timing FormulasA= 1 if block is addressable; else 0.AB= 1 if block is addressable; else 0.AB= 1 if block is addressable; else 0.A1= address of operand 1.A2= address of operand 2.C= 1 if block is connected; else 0.CC= 0 if condition code (CC) is 0S if CC is 1CLZ= 1 if multiplicand has 8 bytes or more of leading zeros; else 0.CZ4= 0 if condition code (CC) is 0.S if CC is 1.CLZ= 1 if multiplicand has 8 bytes or more of leading zeros; else 0.CZ56= number of 256-byte blocks, on 256-byte boundaris, satually compared to the shorter operand, extending through the first doubleword, if any, which compared use unequal. If operands have equal length, this refers to operand 1 alignment.CM= See Figure C-2.DW1= 1 for L2>8 if operand 1 crosses a doubleword boundary, = 1 for L2>8 if the last eight bytes of operand 1 cross a doubleword boundary, else 0.DW2= 1 for lass fit the last eight bytes of operand 2 cross a doubleword boundary, else 0.DW2= 1 for lass if the last eight bytes of operand 2 cross a doubleword boundary, else 0.DW3= 1 if frist 8 bytes ross a doubleword boundary, else 0.DW4= 1 if frist 8 bytes ross a doubleword boundary, else 0.DW4= 1 if frist 6 quotient cross a doubleword boundary, else 0.DW4= 1 if frist 6 quotient cross a doubleword boundary, else 0.DW4= 1 if frist 8 bytes ross a doubleword boundary, else 0.DW7= 1 for lass a dubleword boundary, else								(fo	r L1>8	3 and	d L2>8)	
A = 1 if block is addressable; else 0.EMK = EDMK adjustment. Refer to Figure C-3.AB = 1 if back is cometed; else 0.EMK = EDMK adjustment. Refer to Figure C-3.A1 = address of operand 1.E1 = exponent corresponding to operand 2.C = 1 if block is connected; else 0.FDW1 = 1 if first 8 bytes or operand 1.C = 1 if block is connected; else 0.FDW1 = 1 if first 8 bytes or more of leading zeros; else 0.C = 1 if multiplicand has 8 bytes or more of leading zeros; else 0.F256 = number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand which is actually compared to the shorter operand, extending through the first doubleword, any, which compares unequal. If operand 1 eross a doubleword boundary, else 0.CM = See Figure C-2.M = See Figure C-2.CM = See Figure C-2.L = length of the operand 1 eross a doubleword boundary, else 0.DW1 = 1 for L12 \$ if operand 1 crosses a doubleword boundary, else 0.DW2 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW3 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes cross a doubleword boundary, else 0.DW4 = 1 if first 8 bytes	Legen	d f	or Timing Forn	nulas				E	DW2	=	1 for L2>8 and L2 \leq 16 if the first (L2-8) bytes of	
AB = 1 if branch target is on doubleword boundary,else 0. EMK = EDMK adjustment. Refer to Figure C-3. A1 = address of operand 1. EI = exponent corresponding to operand 1. A2 = address of operand 2. EI = exponent corresponding to operand 1. CC = 1 if block is connected; else 0. EI = exponent corresponding to operand 1. CCV = 0 if condition code (CC) is 0 5 if CC is 2. = number of 256-byte boundaries, spanned by that portion of the longer operand which is refers to operand 1 dingment. F256 = number of 256-byte boundaries, spanned by that portion of the longer operand which is refers to operand 1 alignment. IM = See Figure C-2. CM = See Figure C-2. IM muber of hat digits with value greater than 9. L = length of the operand in bytes. CM = See Figure C-2. IM muber of hat digits with value greater than 9. L = length of the operand in bytes. DW1 = I for L1 2 is 1 forerand 1 crosses a doubleword boundary, else 0. LZ = 0 if NDD1 : S 15 an NDD1-NDD2 > 14 DW2 = I for L1 2 is 1 forerand 1 crosses a doubleword boundary, else 0. LZ = length of operand 1 in bytes. DW3 = I for L1 2 is 1 forerand 2 crosses a doubleword boundary, else 0. LZ <	A	=	1 if block is addressa	ble; else	0.						operand 2 cross a doubleword boundary, else 0.	
A1 = address of operand 1. A2 = address of operand 2. C = 1 if block is connected; lele 0. CV = 0 if condition code (CC) is 0 5 if CC is 1. EDW1 6 if CC is 2. = number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand the is compared to the pad character, extending through the first doubleword, if any, which compares unequal. If operand, have equal length, this refers to operand 1 alignment. CM = See Figure C-2. CM = See Figure C-2. DW1 = I for L1 ≤ 8 if operand 1 crosses a doubleword boundary, else 0. DW2 = I for L2 ≤ 8 if operand 1 crosses a doubleword boundary, else 0. DW2 = I for L2 ≤ 8 if operand 2 crosses a doubleword boundary, else 0. DW3 = I if first 8 bytes cross a doubleword boundary, else 0. DW4 = I if first 8 bytes cross a doubleword boundary, else 0. DW4 = I if first 8 bytes cross a doubleword boundary, else 0. DW4 = I if first 8 bytes cross a doubleword boundary, else 0. DW4 = I if first 8 bytes cross a doubleword boundary, else 0. DW4 = I if first 8 bytes cross a doubleword boundary, else 0. DW4 = I if first 8 bytes cross a doubleword boundary, else 0. DW8	AB	=	1 if branch target is o	on doub	eword bo	undary,e	lse 0.	EN	MK	=	EDMK adjustment. Refer to Figure C-3.	
A2 = address of operand 2. C = 1 if block is connected; else 0. CCV = 0 if condition code (CC) is 0 5 if CC is 1. CLZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. C256 = number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand which actually compared to the horter operand, extending through the first doubleword, if any, which compared to the parade 1 disparant. CM = See Figure C-2. DW1 = 1 for L15 & if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DW2 = 1 for L25 & if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DW3 = 1 if for (L25 & if operand 1 crosses a doubleword boundary. = 1 for L25 & if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DW4 = 1 if for (L25 & if operand 2 crosse a doubleword boundary, else 0. DW4 = 1 if for (length of quotient) > 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DW4 = 1 if the remainder crosse a doubleword boundary, else 0. DW4 = 1 if the remainder crosse a doubleword boundary, else 0. DW7 = 1 if the remainder crosse a doubleword boundary, else 0. DW8 = 1 if the remainder cross a doubleword boundary, else 0. DW8 = 1	A1	=	address of operand 1	•				E1	1	=	exponent corresponding to operand 1.	
 FDW1 = 1 if block is connected; else 0. CCV = 0 if condition code (CC) is 0 5 if CC is 1 6 if CC is 2. CLZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. C256 = number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand this is compared to the pad character, extending through the first doubleword, if any, which compared to the shorter operand which is actually compared to the shorter operand which actes to operand 1 arcosse a doubleword boundary. I for L12 & if operand 1 crosses a doubleword boundary. I for L12 & if operand 1 crosses a doubleword boundary. I for L12 & if operand 2 crosses a doubleword boundary. I for L12 & if the last eight bytes of operand 2 crosse a doubleword boundary, else 0. DWB = 1 if first 8 bytes cross a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary, else 0. EDW = 1 if or (length of quotient) > 8 if the last eight bytes of operand 1 cross a doubleword boundary, else 0. EDW = 1 if or (l	A2	=	address of operand 2	2.				E2	2	=	exponent corresponding to operand 2.	
 CCV = 0 if codition code (CC) is 0 5 if CC is 1 6 if CC is 2. CLZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. C256 = number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand which is actually compared to the shorter operand, if any, which compares unequal. If operands have equal length, this refers to operand 1 alignment. CM = See Figure C-2. DW1 = 1 for L1 ≤ 8 if operand 1 crosses a doubleword boundary. 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. 1 for L2 ≤ 8 if the last eight bytes of operand 1 cross a doubleword boundary, else 0. DW2 = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. 1 for L2 > 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DW3 = 1 if or trosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW5 = 1 if move cod oublewords spanned by pad in operand 1. L8F = number of doublewords spanned by pad in operand 1. <li< td=""><td>C</td><td>=</td><td>1 if block is connected</td><td>ed; else (</td><td>Э.</td><td></td><td></td><td>FI</td><td>DW1</td><td>=</td><td>1 if the first 8 bytes of operand 1 cross a doubleword boundary, else 0.</td><td></td></li<>	C	=	1 if block is connected	ed; else (Э.			FI	DW1	=	1 if the first 8 bytes of operand 1 cross a doubleword boundary, else 0.	
 S if CC is 1 6 if CC is 2. CLZ a 1 if multiplicand has 8 bytes or more of leading zeros; else 0. C256 a number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand which is actually compared to the shorter operand, extending through the first doubleword, if any, which compare unequal. IM See Figure C-2. CM Sift operand 1 crosses a doubleword boundary, else 0. DW2 If or L12 & Si ft operand 2 crosses a doubleword boundary, else 0. DW3 If first 8 bytes of operand 2 crosses a doubleword boundary, else 0. DW4 If or (L12 & Si ft he last eight bytes of operand 2 crosses a doubleword boundary, else 0. DW4 If or (L12 & Si ft he last eight bytes of operand 2 crosses a doubleword boundary, else 0. DW4 If or (L12 & Si ft he last eight bytes of operand 2 crosses a doubleword boundary, else 0. DW4 If or (L12 & Si ft he last eight bytes of operand 2 crosses a doubleword boundary, else 0. DW5 If or (Length of quotient) > 8 if t	CCV	=	0 if condition code (CC) is 0				F2	256	-	number of 256-byte blocks, on 256-byte boundaries,	
 CLZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. CLZ = 1 if multiplicand has 8 bytes or more of leading zeros; else 0. C256 = number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword, if any, which compared to the pad character, extending through the first doubleword boundary, else 0. DW2 = 1 for L1>8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DW3 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. EDW			5 if CC is 1								spanned by that portion of the longer operand that is	
 C256 = number of 256-byte blocks, on 256-byte boundaries, spanned by that portion of the longer operand, extending through the first doubleword, if any, which compares unequal. If operands have equal length, this refers to operand 1 alignment. CM = See Figure C-2. DW1 = 1 for L1 ≤ 8 if operand 1 crosses a doubleword boundary, else 0. DW2 = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary, else 0. DW2 = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary, else 0. DW2 = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary, else 0. DW3 = 1 if first 8 bytes cross a doubleword boundary, else 0. DW4 = 1 if for L3 ≤ 8 if the last eight bytes of operand 1 crosses a doubleword boundary, else 0. DW4 = 1 if for L3 ≤ 8 if the last eight bytes of operand 2 crosses a doubleword boundary, else 0. DW8 = 1 if first 8 bytes cross a doubleword boundary, else 0. DW8 = 1 if first 8 bytes cross a doubleword boundary, else 0. DW8 = 1 if for the autient crosses a doubleword boundary, else 0. DW8 = 1 if for the datient crosses a doubleword boundary, else 0. DW8 = 1 if for length of quotient) > 8 if the last eight bytes of the quotient crosses a doubleword boundary, else 0. DW8 = 1 if for length of quotient) > 8 if the last eight bytes of the quotient crosse a doubleword boundary, else 0. DW8 = 1 if for (length of quotient) > 8 if the first (L1-8) bytes of operand 1. L8 = number of doublewords spanned by pad in operand 1. L8P = number of doublewords spanned by pad in operand 1. ME0W = 1 for (length of quotient) > 8 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for (length of quotient) > 8 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1 ≤ 16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L	CL7	_	6 if UC is 2.	8 hvtes	or more o	f leading	zeros:				compared to the pad character, extending through the first doubleword, if any, which compare unequal.	
C256= number of 256-byte blocks, on 256-byte bloundaries, spanned by that portion of the longer operand which is actually compared to the shorter operand, extending through the first doubleword, if any, which compares unequal. If operands have equal length, this refers to operand 1 alignment.K= Number of hex digits with value greater than 9. LCM= See Figure C-2.UC0 if NDD1 > 15 and NDD1-NDD2 > 14 2 if NDD1 > 15 and ND1-NDD2 > 14 2 if NDD1 > 15 and ND1-ND2 > 14 2 if NDD1 > 15 and ND1-ND2 > 14 2 if NDD1 > 15 and ND1-ND2 > 14 2 if ND1 > 15 and ND1-ND2 > 14 		-	else 0.	8 Dytes	or more o	r icading	, 20103,	IM	Л	=	See Figure C-2.	
 spanned by that portion of the longer operand, which is actually compared to the shorter operand, extending through the first doubleword, if any, which compares unequal. If operands have equal length, this refers to operand 1 alignment. CM = See Figure C-2. DW1 = 1 for L1 ≤ 8 if operand 1 crosses a doubleword boundary, = 1 for L1 ≤ 8 if operand 2 crosses a doubleword boundary, = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary, else 0. DW2 = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary, else 0. DW8 = 1 if first 8 bytes crosses a doubleword boundary, else 0. DWQ = 1 for (length of quotient) > 8 if the last eight bytes of operand 2 crosses a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary, else 0. DWR = 1 if or (length of quotient) > 8 if the first (L1-8) bytes of of the quotient cross a doubleword boundary, else 0. DWR = 1 for (length of quotient) > 8 if the first (L1-8) bytes of of the quotient cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the	C256	=	number of 256-byte	blocks,	on 256-by	te bound	faries,	к		=	Number of hex digits with value greater than 9.	
Is actually compared to the shorter operand, extending through the first doubleword, if any, which compares unequal. If operands have equal length, this refers to operand 1 alignment.LC=0 if NDD1 ≤ 15 1 if NDD1 > 15 and NDD1-NDD2 ≤ 14 2 if NDD1 > 15 and NDD1-NDD2 > 14CM=See Figure C-2.LZ=number of contiguous zero result bytes starting from the beginning of the result.DW1=1 for L1 ≤ 8 if operand 1 crosses a doubleword boundary. =1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. =1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. =1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. =1 for L2 ≤ 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0.LZ=number of contiguous zero result doublewords starting from the beginning of the result.DW2=1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. =1 for L2 ≤ 8 if the last eight bytes of operand 2 crosses a doubleword boundary, else 0.LSC=number of doublewords spanned by processing.DW8=1 for (length of quotient) ≤ 8 if the quotient crosses a double word boundary. =1 for (length of quotient) ≥ 8 if the last eight bytes of the quotient cross a doubleword boundary; else 0.L8F=number of doublewords spanned by move in operand 1.DWR=1 if the remainder crosses a doubleword boundary; else 0.=1 for (length of quotient) > 8 if the first (L1-8) bytes of the quotient cross a doubleword boundary, else 0.L8F=number of doublewords spanned by move in operand 1.DWQ=1 for (length of quo			spanned by that port	ion of th	ne longer o	operand	which	L		=	length of the operand in bytes.	
 compares unequal. If operands have equal length, this refers to operand 1 alignment. CM = See Figure C-2. DW1 = 1 for L1>8 if operand 1 crosses a doubleword boundary. a for L1>8 if the last eight bytes of operand 1 cross a doubleword boundary, else 0. DW2 = 1 for L2>8 if operand 2 crosses a doubleword boundary, else 0. DW3 = 1 if first 8 bytes cross a doubleword boundary, else 0. DW4 = 1 if or (length of quotient) >8 if the last eight bytes of the quotient crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW4 = 1 if the remainder crosses a doubleword boundary, else 0. DW5 = 1 if the remainder crosses a doubleword boundary, else 0. DW6 = 1 if the remainder crosses a doubleword boundary, else 0. DW6 = 1 if the remainder crosses a doubleword boundary, else 0. EDW4 = 1 for (length of quotient) >8 if the first (L1-8) bytes of of the quotient cross a doubleword boundary, else 0. EDW4 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of of the quotient cross a doubleword boundary, else 0. EDW4 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW5 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cr			extending through th	to the s he first d	norter ope	erano, 1. if anv.	which	LC	С	=	$0 \text{ if NDD1} \leq 15$	
CM = See Figure C-2. LZ = number of contiguous zero result bytes starting from the beginning of the result. DW1 = 1 for L1 ≤ 8 if operand 1 crosse a doubleword boundary, else 0. LZ = number of contiguous zero result bytes starting from the beginning of the result. DW2 = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary, else 0. L1 = length of operand 1 in bytes. DW2 = 1 for L2 > 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0. L2 = length of operand 1 doublewords spanned by processing. DW8 = 1 if first 8 bytes cross a doubleword boundary, else 0. L8C = number of doublewords in the longer operand Nich are actually compared to the shorter operand. If operand 1 doublewords in the longer operand Mich are actually compared to the pad character. DW4 = 1 for (length of quotient) > 8 if the last eight bytes of the quotient cross a doubleword boundary, else 0. L8F = number of doublewords in the longer operand Nich are actually compared to the pad character. DW4 = 1 for (length of quotient) > 8 if the last eight bytes of the quotient cross a doubleword boundary, else 0. L8F = number of doublewords spanned by move in operand 1. DW6 = 1 for (length of quotient) > 8 if the last eigh			compares unequal. I	f operan	ds have e	qual leng	th, this				1 if NDD1 > 15 and NDD1-NDD2 \leq 14	
CM= See Figure C-2.LZ= number of contiguous zero result bytes starting from the beginning of the result.DW1= 1 for L1 ≤ 8 if operand 1 crosses a doubleword boundary. = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. = 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. = 1 for L2 ≤ 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0.LZ= number of contiguous zero result doublewords starting from the beginning of the result. L1DW2= 1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary. = 1 for L1 ≥ 8 if the last eight bytes of operand 2 cross a double word boundary, else 0.L2= length of operand 1 doublewords spanned by processing. L8CDW3= 1 if first 8 bytes cross a doubleword boundary, else 0.L8C= number of doublewords in the longer operand which are actually compared to the shorter operand 1. operands have equal length, this refers to operand 1. L8FDW4= 1 if the remainder crosses a doubleword boundary, else 0.L8F= number of doublewords spanned by move in operand 1.DW7= 1 for (length of quotient) > 8 if the first (L1-8) bytes of of the quotient cross a doubleword boundary, else 0.L8P= number of doublewords spanned by pad in operand 1.ED= £D, EDMK base value. Refer to Figure C-3.M= 0 if all test mask bits selected are 0; else 1.EDW0= 1 for (length of quotient) > 8 if the first (L1-8) bytes of of the quotient cross a doubleword boundary, else 0.MA <td< td=""><td></td><td></td><td>refers to operand 1 a</td><td>lignmen</td><td>t.</td><td></td><td></td><td></td><td>_</td><td></td><td>2 if NDD1 > 15 and NDD1-NDD2 > 14</td><td></td></td<>			refers to operand 1 a	lignmen	t.				_		2 if NDD1 > 15 and NDD1-NDD2 > 14	
DW1=1 for L1 \$ 8 if operand 1 crosses a doubleword boundary.LZ8=number of contiguous zero result doublewords starting from the beginning of the result.DW2=1 for L1>8 if the last eight bytes of operand 1 cross a doubleword boundary.=1 lor L2 \$ 8 if operand 2 crosses a doubleword boundary.=1 lor L2 \$ 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0.L2=length of operand 1 in bytes.DW2=1 for L2 \$ 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0.L2=length of operand 1 doublewords spanned by processing.DW8=1 if first 8 bytes cross a doubleword boundary, else 0.L8C=number of doublewords in the longer operand .If operands have equal length, this refers to operand 1.DWQ=1 for (length of quotient) \$ 8 if the last eight bytes of the quotient cross a doubleword boundary, else 0.L8F=number of doublewords in the longer operand which are actually compared to the pad character.DWR=1 if the remainder crosses a doubleword boundary; else 0.L8M=number of doublewords spanned by pad in operand 1.EDWQ=1 for (length of quotient) > 8 if the first (L1-8) bytes of the quotient cross a doubleword boundary, else 0.MEOB=1 if move ends on other than a 256 byte boundary in operand 1; else 0.	CM	=	See Figure C-2.					LZ	Z	=	number of contiguous zero result bytes starting from the beginning of the result.	
 a for L1>8 if the last eight bytes of operand 1 cross a doubleword boundary, else 0. DW2 = 1 for L2≤8 if operand 2 crosses a doubleword boundary. a 1 for L2>8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DWB = 1 if first 8 bytes cross a doubleword boundary, else 0. DWQ = 1 for (length of quotient) ≤8 if the last eight bytes of the quotient cross a doubleword boundary, else 0. DWR = 1 if the remainder cross a doubleword boundary, else 0. DWR = 1 if the remainder cross a doubleword boundary, else 0. DWR = 1 if the remainder cross a doubleword boundary, else 0. DWR = 1 if the remainder cross a doubleword boundary, else 0. DWR = 1 if the remainder cross a doubleword boundary, else 0. ED = £D, EDMK base value. Refer to Figure C-3. ED = 1 for (length of quotient) >8 if the first (L1-8) bytes of the quotient cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. MEOB = 1 if move ends on other than a 256 byte boundary in operand 1 cross a doubleword boundary, else 0. 	Dwi		I for $L1 \leq 8$ if operation boundary.	nd I cros	sses a dou	bleword		LZ	Z8	=	number of contiguous zero result doublewords	
doubleword boundary, else 0.L1=length of operand 1 in bytes.DW2=1 for L2 ≤ 8 if operand 2 crosses a doubleword boundary.=1 L2=length of operand 2 in bytes.=1 for L2 > 8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0.L8=number of operand 1 doublewords spanned by processing.DW8=1 if first 8 bytes cross a doubleword boundary, else 0.L8C=number of doublewords in the longer operand. If operands have equal length, this refers to operand 1.DWQ=1 for (length of quotient) ≤ 8 if the quotient crosses a double word boundary.=1 lif the remainder crosses a doubleword boundary, else 0.DWR=1 if the remainder crosses a doubleword boundary, else 0.L8F=number of doublewords in the longer operand hich are actually compared to the pad character.DWR=1 if the remainder crosses a doubleword boundary, else 0.L8F=number of doublewords spanned by move in operand 1.DWR=1 if the remainder crosses a doubleword boundary, else 0.L8F=number of doublewords spanned by move in operand 1.ED=ED, EDMK base value. Refer to Figure C-3.M=0 if all test mask bits selected are 0; else 1.EDW1=1 for (L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0.MEOB=1 if move ends on other than a 256 byte boundary in operand 1; else 0.		=	1 for $L1 > 8$ if the las	t eight b	ytes of op	erand 1	cross a				starting from the beginning of the result.	
 DW2 = 1 for L2≤8 if operand 2 crosses a doubleword boundary. = 1 for L2>8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DWB = 1 if first 8 bytes cross a doubleword boundary, else 0. DWQ = 1 for (length of quotient) ≤8 if the quotient crosses a double word boundary. = 1 for (length of quotient) >8 if the last eight bytes of the quotient crosses a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary; else 0. DWR = 1 if the remainder crosses a doubleword boundary; else 0. DWR = 1 if the remainder crosses a doubleword boundary; else 0. ED = ED, EDMK base value. Refer to Figure C-3. EDWQ = 1 for (length of quotient) >8 if the first (L1-8) bytes of the quotient cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. 			doubleword boundar	ry, else ().			LI	1	=	length of operand 1 in bytes.	
boundary.L8=number of operand 1 doublewords spanned by processing.DWB=1 for L2>8 if the last eight bytes of operand 2 cross a doubleword boundary, else 0.L8=number of doublewords in the longer operand which are actually compared to the shorter operand. If operands have equal length, this refers to operand 1.DWQ=1 for (length of quotient) ≤8 if the quotient crosses a double word boundary. =1 for (length of quotient) >8 if the last eight bytes of 	DW2	=	1 for $L2 \le 8$ if operation	nd 2 cros	sses a dou	bleword		L2	2	=	length of operand 2 in bytes.	
 a Troi 122% if the last eight bytes of operand 2 cross a doubleword boundary, else 0. DWB = 1 if first 8 bytes cross a doubleword boundary, else 0. DWQ = 1 for (length of quotient) ≤8 if the quotient crosses a double word boundary. = 1 for (length of quotient) >8 if the last eight bytes of the quotient cross a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary; else 0. DWR = 1 if the remainder crosses a doubleword boundary; else 0. ED = €D, EDMK base value. Refer to Figure C-3. EDWQ = 1 for (length of quotient) >8 if the first (L1-8) bytes of the quotient cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. 		_	boundary.	t aight h	wter of or	arand 2	oross a	L8	8		number of operand 1 doublewords spanned by	
 DWB = 1 if first 8 bytes cross a doubleword boundary, else 0. DWQ = 1 for (length of quotient) ≤8 if the quotient crosses a double word boundary. = 1 for (length of quotient) >8 if the last eight bytes of the quotient cross a doubleword boundary, else 0. DWR = 1 if the remainder crosses a doubleword boundary; else 0. ED = ED, EDMK base value. Refer to Figure C-3. EDWQ = 1 for (length of quotient) >8 if the first (L1-8) bytes of the quotient cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. EDW1 = 1 for L1>8 and L1≤16 if the first (L1-8) bytes of operand 1 cross a doubleword boundary, else 0. 		=	doubleword boundar	rv. else ().		cross a				processing.	
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$EDW1 = 1 \text{ for } L1>8 \text{ and } L1 \le 16 \text{ if the first (L1-8) bytes of} \\ operand 1 \text{ cross a doubleword boundary, else 0.} MEOB = 1 \text{ if move ends on other than a 256 byte boundary in} \\ operand 1; else 0.$	EDWQ	-	I for (length of quot of the quotient cross	a doubl	s if the first eword bo	st (L1-8) undary, e) bytes else 0.	М	IA	=	0 if operands are mutually aligned on doubleword boundaries, else 1.	,
	EDW1	=	1 for L1>8 and L1 \leq operand 1 cross a do	≤16 if th oublewor	e first (L1 d bounda	-8) byte ry, else (s of).	М	IEOB	=	1 if move ends on other than a 256 byte boundary in operand 1; else 0.	

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M256	=	number of 256 byte blocks, on 256 byte boundaries, spanned by move in operand 1.
N	=	number of registers in LM, STM.
N1	=	1 if operand 1 is negative; else 0.
N2	=	1 if operand 2 is negative; else 0.
NDD1	=	number of significant decimal digits in operand 1.
NDD2	=	number of significant decimal digits in operand 2.
NHD	=	number of significant hex digits.
NP	=	number of bytes processed.
NWB	=	number of source or destination words that cross doubleword boundaries.
ORS	=	1 if starting register number is odd, else 0.
Р	-	0 if processing halts before the last doubleword is reached; 1 if the last doubleword is processed.
PN	=	1 if post normalization is required; else 0.
PU	=	1 if Translation Lookaside Buffer purge is required; else 0.
P256	=	number of 256 byte blocks, on 256 byte boundaries, spanned by the pad in operand 1.
R	=	1 if returned to the instruction following the execute instruction after completing the subject instruction; else 0.
RC	=	1 if recomplementation is required; else 0.
RMN	=	1 if remainder is negative; else 0.
RSA	=	1 if starting register number is even and op2 is off doubleword boundary or if starting register number is odd and op2 is on a doubleword boundary; else 0.
S	=	1 if successful branch; else 0.
SA	=	sign adjustment. Refer to Figure C-3.
SI	=	time to execute the subject instruction.
VP	=	value of the operand 2.
х	-	1 if index register number is not zero; else 0.
ZR	-	1 if result is zero; else 0.
•	=	multiply.
()	=	1 if the logical condition within () is satisfied; else 0.

Figures C-2 and C-3 are used to determine the processing time for each pattern character occurring in ED and EDMK instructions. For each such pattern character, determine the EB value and add it to the instruction time. If the pattern character is a digit selector or a significance starter, and if the source digit is the lowest-order digit in its field, add in also the sign adjustment (SA). If the instruction is EDMK and a nonzero source digit is encountered with the significance indicator off, add the EMK adjustment to the instruction time.

Mask	СМ	IM	Mask	СМ	IM
0000	0	0	1000	1	1
0001	1	0	1001	2	2
0010	1	1	1010	2	2
0011	1	0	1011	3	4
0100	1	1	1100	1	1
0101	2	2	1101	3	4
0110	1	1	1110	2	3
0111	2	2	1111	1	0

Figure C-2. Mask Character

High-Speed Buffer Storage and TLB Miss Service Times

These timings (in microseconds) are added to the instruction execution time for each miss in the high-speed buffer storage.

High-speed buffer storage miss where page to be replaced has not been altered.
High-speed buffer storage miss where page to be replaced has not been altered.
High-speed buffer storage miss where page to be replaced has been altered (castout case).

For typical workloads, the average high-speed buffer storage miss time is approximately 1.34 microseconds. The frequency depends upon the addressing pattern of the program being executed.

Pattern Character Type	Significance Indicator	Source Digit	EB Value	SA - Sign Adj Add if digit is last before:	ustment the		EMK Adjustment add if
and the second second				Sign B	Sign D	+Sign	EDMK
Digit Selector	Off Off On On	0 1-9 0 1-9	1323 1553 1323 1438	345 518 518 518	345 690 690 690	345 690 805 805	0 230 0 0
Signifi- cance Starter	Off Off On On	0 1-9 0 1-9	1840 1955 1668 1783	518 518 460 460	690 690 633 633	575 690 748 748	0 230 0 0
		First Byte of Pattern					.
Field Separator		Yes No	575 920				
Message Char.	Off Off On	Yes No 	575 920 690				

Figure C-3. ED and EDMK Pattern Character Timings

For each miss in the Translation Lookaside Buffer (TLB), these times are added:

- 1.04 μs TLB miss for S/370 mode, BC mode or EC DAT off.
 0.98 μs TLB miss for ECPS:VSE mode.
- 1.38 μ s TLB miss for S/370 mode, EC mode DAT on.

These times assume that the necessary translation table entries are found in the high-speed buffer storage. The TLB miss frequency depends on the addressing pattern of the program being executed. Values outside the ranges given are possible.

Performance Degradation of PER Mask Setting

With bit 1 of the PSW, EC mode turned on, and the PER mask bits all zero, there is no performance degradation. With PER mask bits enabled for various categories of events shown below, and with typical instruction mixes, PER degradation (time in addition to normal time), is approximately as follows:

PER Enabled For	Degradation
Register alteration	620 percent
Instruction Fetch	730 percent
Storage alteration	570 percent
Successful branch	520 percent

The above figures are for cases in which no PER events are raised. The accumulative degradation with more than one event selected is not an addition of the above figures.

Effect of Hardware-Assist Features on Performance

The ECPS:VM Assist and ECPS:VS1 Assist facilities simulate certain frequently used functions in hardware. The effect that these have on performance depends on the workload being executed and the frequency with which it requests services that are assisted.

ECPS:VM Assist and ECPS:VS1 Assist are mutually-exclusive facilities. Also, ECPS:VM Assist and ECPS:MVS are mutually-exclusive facilities. If the Control Store Expansion feature is installed, ECPS:VM and ECPS:MVS can be run concurrently and provide a relative batch throughput improvement of 1.8 to 4.0 times the VM environment with only the ECPS:MVS assist present.

Engineering Scientific Assist improves the native assembler code of equal function up to 30 percent.

4341 Processor Model Group 12 Complex Configurator



Figure C-4. 4341 Processor Model Group 12 Complex Configurator

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